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Advanced Computer Program Models : A Talking Textbook Based on Three Languages

John Thomas Gardner
Walden University

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ADVANCED COMPUTER PROGRAM MODELS
A TALKING TEXTBOOK
BASED ON THREE LANGUAGES

by

John Thomas Gardner

B.M.E. Kansas State College, 1951

M.S. Kansas State University, 1957



Robert W. Swanson, Ph.D., Advisor
Director of Computer Center
California State University
San Diego

A Project Submitted in Partial Fulfillment of
The Requirements for the Degree of
Doctor of Philosophy

Walden University
July, 1972

AN ABSTRACT OF A DISSERTATION

By

John Thomas Gardner

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The purpose of this dissertation was to develop a learning instrument, to be used by programmers preparing for the Data Processing Management Association Test as a self study book, or by college business programming and computer science students who have completed a course in data processing and a course in programming a higher level language.

The mathematical ability requirement was minimized by developing the algorithms in parallel with the programs.

The learner should experience emphasis in the following areas:

1. The type of activities required to pass the DPMA test
(the programming part)
2. Data Structures
3. Fortran (at the level of the DPMA test)
4. RPG (at the level of the DPMA test)
5. Flow chart reading and writing

Fortran and RPG (Report Program Generator) languages were used, since their proficiency is required for the DPMA test; however a subset of IBM Basic Assembler language was used, because the author believed that a person who is more than superficially interested in computers should demonstrate a proficiency with a machine language.

An important part of this method of presentation are the cassette recordings which allow the learner to progress outside the classroom. The recordings plus the hard copy of the actual programs, diminished in size, give the learner material which he can move to any location and study without the presence of the instructor.

Organization

The book is divided into four main parts:

1. Introduction (containing the Swanson Systems Study)
2. Data Structures
3. Fortran
4. RPG

The Introduction contains the Swanson Study which provides a model to answer the question "WHY" for much of the rest of the book.

Section II, Data Structures, includes many of the topics included in the recommendations of the Curriculum Committee on Computer Science of the Association of Computing Machinery.

Section III, Fortran, is designed to prepare the advanced programmer for the DPMA test, but can be used to develop programs in other languages by using the fortran programs presented here as models. This section does contain some BAL programs.

Section IV, RPG, provides models which will aid in the preparation for the DPMA test. The BAL models used here were selected to match the same topic, Invoicing, as the RPG program.

Experimental Results

The experimental objective was to determine to what extent, if any, Advanced Computer Program Models affected the attitudes towards and proficiencies in programming of college level students in San Diego County.

The original design of the experiment was a completely crossed three by four factorial experiment: three level (novice, beginner and expert) versus four types of computer instructions: (1) standard classroom, (2) no classroom, only the talking textbook material, (3) classroom

and material, (4) open study with no guidance from instructor or experimental material.

The observations are individual students' different scores with respect to a standard examination given in January at the start of the school semester, and the same examination given in May at the end of the semester. Classes were held for one night a week for a total of eighteen meetings. One section consisted of students learning computer programming. The other sections consisted, with one exception, of programmers who had some experience with programming. In this section the objective was to develop techniques for data structures.

There were several factors which prevented us from fulfilling the requirements of the originally completed cross factorial experiment:

1. We were dependent on the availability of computer equipment.
2. We could not control the composition of the programming classes so that one class was a mixed grade level and the other class was not mixed.
3. We were dependent upon volunteers and had no control over the number of programming sections assigned to these volunteers.

To measure the affect on proficiency, we used the CTSS test data bank maintained in the California State University, San Diego. To measure the affect on attitude, we used the Aiken Devised Programming Attitude Scale, which is a Likert-type scale, developed by Louis Aiken, professor of psychology, at Guilford College. They were administered once within a week period between January 15th and January 21st, and

the second time between May 15th to May 22, 1972. The tests were administered by the instructors to their own classes. The grading and scoring was performed by the staff at San Diego University.

A statistical analysis was performed with the assistance of the computer facilities of San Diego State University. Biomed Statistical Program, BMD-V, was used in one analysis. The remaining analyses and data summaries were made using the statistical program library (S.T.L.) available at the University.

Conclusions

The conclusions derived from the testing and experiments were: Novices scored significantly higher with classroom instruction with little difference in text used. The absence of significance in the effects of the classroom experience and text on student proficiency and attitude in the other groups was contrary to expectations. The curriculum in question had been developed for a different approach for the subject. At the end of the experiment, the instructors were of the opinion that for a more effective utilization of the material, it would be necessary to adapt the text to the curriculum.

Whether or not a curriculum especially adapted to this experiment would show any changes in student proficiency or attitude is another question for objective evaluation.

In any case, it is clear that the function of this text, as it was coordinated for this experiment, did not provide any reasonable improvement in the overall student proficiency or attitude beyond the variations associated with students, teachers and experience level.

PREFACE

This book is a learning instrument, to be used by programmers preparing for the Data Processing Management Association Test as a self study book, or by college business programming and computer science students who have completed a course in data processing and a course in programming a higher level language.

The mathematical ability requirement of the student has been minimized, since the algorithms are developed in parallel with the programs.

The learner will experience emphasis in the following areas:

1. The type of activities required to pass the DPMA test,
(the programming part)
2. Data Structures
3. RPG (at the level of the DPMA test)
4. Fortran (at the level of the DPMA test)
5. Flow chart reading and writing

Fortran and RPG (Report Program Generator) languages were used since their proficiency is required for the DPMA test; however a subset of IBM BAL language was used because I believe that the person who is more than superficially interested in computers should be able to demonstrate a proficiency with a machine language.

The cassette recordings are an integral part of the presentation, since they allow the reader to progress outside the classroom. The recordings, plus the hard copy of actual programs, diminished in size, give the learner material which he can move to any location and study without the presence of the instructor.

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Section III, Fortran, is designed to prepare the advanced programmer for the DPMA test, but can be used to develop programs in other languages by using the Fortran programs presented here as models. This section does contain some Basic Assembler programs.

Section IV, RPG, provides models and information which will aid in the preparation for the DPMA test. The Basic Assembler program presented here is used as a comparison for the invoicing program, written in RPG.

HOW TO USE THIS BOOK

The programmer reviewing for the DPMA test can study the Swanson Study, use the cassettes, review the Fortran Section (paying particular attention to the first half), and review the RPG Section. This can be done in six weeks without an instructor.

The college student can make a semester, or a year study, with or without an instructor. The type of computer is incidental, since the tape cassettes analyze each BAL program. In fact a computer isn't needed (however it does work as a motivating force). This allows several different languages to be programmed at the same time during a semester.

An instructor who uses this book need not know the Basic Assembler Language, because the tapes will assist him.

The long programs were diminished in size intentionally so that a learner could manipulate them with ease. It is sometimes better for analysis, if the learner copies certain sections manually, if he wants to carry on a concentrated study.

The book starts with a real systems study, which uses a fictitious name. The study was made by the firm of Swanson Associates, (real name), and the corporation in the study is the Marine Division of the Oceano-Graphic Corporation (fictitious name).

Why use a systems study to start a programming text? The answer to this question is that the study gives meaning to the whole text, and serves as a vehicle for the text. The programmer reviewing for the DPMA test might have lost view of the whole corporation, he might have forgotten some of his systems work he had a few years ago,

and might not have a model of a systems study to review before his test.

The college student might have a good background in accounting, but it will add to his ability if he can work with a model that develops a management information system, which includes the whole or at least a part of the corporation.

ACKNOWLEDGEMENTS

Information and illustrative materials have been used from publications copyrighted by the International Business Machines Corporation.

Direct quotations appear as itemized below:

1. C-26-3600 RPG Programs Invoicing and Pre-Billing Calculation with Inventory Control
2. C-33-6000 RPG Disk Examples

The Systems Study was adapted from a real systems study by Dr. R. L. Swanson and Associates.

Help from many sources has contributed to this book. To my students and teachers, named and unnamed, I extend my deepest gratitude. I also express my thanks to Mr. Wayne Langley, Mr. Richard Schenck, Mr. Lawrence Huster, Mr. Herman Dailey, Mr. Roy Towne and Mrs. Ethel Castle, who taught, shared knowledge and provided aid in working various programs.

Special thanks go to Dr. Robert Swanson, who patiently guided me through the dissertation and provided the systems study upon which the book is based.

Most of all, I wish to express my deepest appreciation to my wife, Patty, and seven children, Clorinda, Patrice, Deborah, Curtis, Melisa, Michele and Melanie, who were, I am sure, very pleased when "the computer book" was finished.

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SECTION I

INTRODUCTION

THE SWANSON SYSTEMS STUDY

The objective of this section is to present a complete model of a systems study, and the presentation of the files serve as models to be programmed at the end of the book.

Since so many elementary programming books start with explanations of computers and use only short examples, it was felt that an advanced book should contain a system study and complete models of programs when possible.

It is suggested that the reader return again and again to the charts of this study so he can develop a feeling for the functional groupings of this real corporation.

PRODUCTION SYSTEM PROPOSAL
for the
MARINE DIVISION
of
OCEANO-GRAPHIC CORPORATION

by

Dr. Daryl G. Mitton

Dr. Robert W. Swanson

1.0 ABSTRACT

This is a proposal of a system design for the Marine division of the Oceano-Graphic Corporation. We believe its unique features will provide a multitude of advantages for the firm including:

- a. A planning and control capability which cannot even be approached by your current system and which, we feel, has not been attained for any similar sized company in the past.
- b. A capability of response through information accessibility for all tactical and strategic situations.
- c. A reduction in both direct and indirect production costs-- at present or increased production levels.
- d. A competitive advantage in the ability to obtain proper proposal information from a technical, cost, and delivery standpoint.

- e. A reduction in inventory and inventory costs (we feel that under existing output levels, the present inventory level of \$85,000 can be cut to \$20,000).
- f. The availability of analyses of consequences of alternative decisions.
- g. The many sundry advantages brought about by good systems design, such as more complete and improved cost information, the signaling of impending operational difficulties, the reduction of operation cycle time, the improvement of customer service and feedback, the determination of future equipment and manpower needs, and the preparation of most essential paperwork.

The strength of the system is the ability to obtain needed information easily. This strength provides a strong motivation to keep the system operative. No operational system, regardless of its detailed design, will function properly unless it is used as its design intended it to be used. Most systems are designed in such a way that their intended function is obstructed in some way to satisfy organizationally local needs. Therefore, most systems fall short of their intended goals.

At present, the Marine Division of the Oceano-Graphic Corporation has a very detailed and elaborate system for operational planning and control. It does not function. Its literal atrophy is caused by an informal realization at all levels within the division that the struggle to supply input and the difficulty of receiving output are not worth the rewards obtainable by the functioning system.

This is no direct criticism of the Marine Division system, for what is symptomatic here is found in most organizations: the super-

imposition of an informal "make-do" system which operates on "get by" and exists to accommodate certain paper work demands that may or may not have any real relevance to the effective running of an organization. In fact, it is not uncommon for the resulting working system to be actively dysfunctional, particularly in the long run.

Our study first confirmed the existence of a non-working system. We next submitted preliminary designs of two possible alternative systems which would provide a greater degree of automaticity in speeding input and spilling output. The first was a somewhat conventional approach, an integrated data processing system, which would mechanize the various subsystems (a subsystem at a time) that impacted on the production function--production, purchasing, accounting, engineering, etc. Each system, being separate to orient its particular output demands, would, nonetheless, be updated by information fed into any of the subsystems. Output would have a specialist orientation--not suited to generalship. The data files would be redundant and the system would never be totally current. We admit to editorializing against this system in favor of a second approach, which the Marine Division decided to pursue.

The system devised is built around the concept of a single flow of information. Essential to this system is a common data base. Information input is, therefore, immediately triggered to update purchasing, production, engineering, etc. Information is always current. No output is made unless called for. Information release is on a "need to know" basis rather than a "nice to know" basis, and it is in an interpretable form to serve as a decision base. The ability to obtain needed information very easily serves as the strongest motivation to supply essential input. In spite of the universality of the system, we have

devised a way to introduce it in modules, which will enable faster realization from the system as it is programmed and introduced, and which provides for current debugging of the programming.

We recommend an IBM 000/000 computer as the equipment best suited for immediate and future needs of the Marine Division. We estimate its first-year costs at \$00,000 (\$0,000 per month for ten months) to program and debug the system recommended. Current Marine Division personnel would be able to man the program, so no additional cost would be incurred here. It is unfortunate, in a report such as this, that we cannot include as a cost the price of not making the change, for we feel, if the Marine Division is to be a part of the growth of oceanographic explosion that is sure to come, these changes must be made. What is the cost of losing a position of leadership in oceanographic instrumentation? What is the cost of losing a contract because of poor performance? What is the cost of slipping ever so slightly behind the competitor? There are costs that do not show in the profit and loss statement directly, but their costs are just as real as those that are displayed.

1.1 PLAN OF REPORT

The report is presented in five parts, starting with a discussion of the objectives and limitations. This is followed by an explanation of the methods used and alternatives available to us in conducting the study and design. The next section of the report is devoted to a summary analysis of the existing system and what we considered to be the major limitations for continued successful use of the system as it currently exists. Beginning in Section 5.0, we present the proposed single

information flow system, complete with the details of the proposed filing system. In Section 5.5, we further specify the benefits that can accrue from the employment of such an approach to the production information system. Section 6.0 of this report deals with our estimated cost and schedule of implementation.

2.0 OBJECTIVES

The objective of this study is to analyze existing procedures and conditions of operations prevalent in the Marine Division of the Oceano-Graphic Corporation to determine the feasibility of use and choice between available methods of operations of automatic data processing systems.

A part of this objective has been to develop an operating system that would provide the management of this division with insight and control of the production function, insight being defined as access to the facts in an orderly manner with an insured degree of accuracy.

A second function of the objective was to develop the state of automation to provide individual response to the system.

The functional criteria for measuring how well these objectives can be obtained are:

- a. Reduction of the number of manual checking functions by one-half (i.e., routine decisions and paperwork).
- b. The increase of the division's profit to thirty percent prior to taxes.
- c. Reduction of the total direct and indirect cost of production by twenty percent.
- d. Reduction of the response time to a customer's order to thirty days.

2.1 SCOPE

There are many limitations to a report of this nature. Time only permits the minimum of attention to the existing disciplines and methods of management information systems. The chief center of interest here is to determine the prospective utilization or applicability of present systems. Further, since many management decision factors are prefaced upon future events not totally definable now, no effort has been made to define the future needs of Oceano-Graphic other than what is needed now. These decisions must be reserved for the future as the technology evolves; however, the provisions for growth were considered and incorporated as a function of the study. When we consider the peculiarly sensitive nature of a company's investment in computers--with its potential for major impact on its position in competition--judgments on how, where and when--the investment should be made to assume special importance. As a result of advances in technology, the computer can now play a more central role in corporate planning and operations. Since present criteria for allocating resources to ADP are based on outdated data processing operations or on rules of thumb, they do not take into account the ability of computers to contribute to profits and/or cut operating costs outside of ADP. Clearly, it would be more appropriate to evaluate the computer investment in terms of its contribution to the entire management process.

We recognize how difficult a task this will be. There can be no simple rules of thumb for analyzing Oceano-Graphic's ADP expenditures. These expenditures must be geared not to the company's size or competitor's spending, but to the benefits to be derived in each specific case. Since these benefits often accrue from the value of the information to

be supplied by the system, as well as from the added efficiencies and direct savings in data processing operations, they are elusive and extremely difficult to measure.

It should further be expressed as a function of the scope of this study that we, as consultants, are not in the business of selling computer hardware. Our task is to evaluate the present system and recommend changes as they relate to the objectives of the study. If the results of this study are implemented, the choice of equipment would be the decision of the Oceano-Graphic Corporation.

3.0 BACKGROUND

Our approach to this particular study was to review the present system as it is currently operating. This review included those documents and procedures that are defined in the systems and procedure manuals and a complete survey of the actual operations and procedures as interpreted by those performing the functions. There always exists a certain dichotomy in all enterprises. There are the procedures set by the workers and lower levels of management and the procedures formally defined by top management. Rarely are the interpretations identical. Suborganizations within organizations adopt symptoms to satisfy the needs of their organizational area. They add to and ignore the system as they see it and make it tolerable and workable from their point of view. Therefore, there exists a system working the way the lower echelons believe it has to work. It is necessary to review both systems.

Each of the activities was reviewed from the perspective of how well it contributed to the production of the end product and how well it fit into the function of management control. Instead of starting

with the costs of processing data, we started with the positive value of the information produced by the processing. The benefits do not stem directly from the fact that given data are processed; they stem from the results of data processing. The value of information from data processing is not what it costs to obtain but, rather, what it can do for management. Costs have their place, but in a different part of the equation.

This, the first effort becomes one of analysis of the intended system. It is an effort to establish a system based upon the total objectives of the company. Recognizing the profitability of rapid growth in the oceanographic field and the limitations of time and cost in designing and implementing a total system, we feel there are two feasible alternatives to choose from--integrated data processing or single information flow.

The integrated data processing approach is an attractive alternative from several aspects. First of all, it is the conventional approach. Most commitments in the field of data processing take this approach. It is a step-by-step approach. You can merchandise a subsystem at a time, thereby getting a system on the computer, gaining some utilization and economic value out of the system without the total system being designed. In this case, production could be put into operation first, followed by a purchasing system, then accounting, etc. All of the subsystems would be designed so that information from one system would be used to update other subsystems, thus integrating all the systems into a single integrated whole. Each of the subsystems would produce the desired reports. In the technical jargon, we refer to it as an "output oriented system."

The single information flow concept developed from some of the

apparent limitations of the integrated data processing (IDP) approach. First, consider the IDP method each subsystem being linked to the other subsystems. This means each subsystem must be updated without respect to the volume of transactions affecting it. Secondly, each subsystem must, by its nature of independence, support separate filing systems, much of which is redundant data. This requirement for processing and redundant files necessitates considerably more machine capacity than would be required if they could be eliminated. A second deficiency arises in that the system is never totally current by the nature of the interlacing of subsystems. Thus, management in each of the separate functions always sees different data creating a continuous point of contention.

This system is designed whereby the source information in any department updates a single information file. For example, an engineering change would update the purchasing files and production files, as well as the engineering files at the time of entry. Conceptually, they would be a single file. In this mode, we could eliminate a considerable amount of processing. In the technical sense, the system would become input oriented. The results would then be a system containing the same information bank, but requiring less equipment and more current data.

When presented with these alternatives, the management of the Marine Division chose to pursue development in the direction of the single information flow concept. The design of the system was then undertaken with the intent of developing this concept on a limited scope, primarily directed to production, inventory control, material control, and purchasing functions.

3.1 PROCEDURAL ANALYSIS

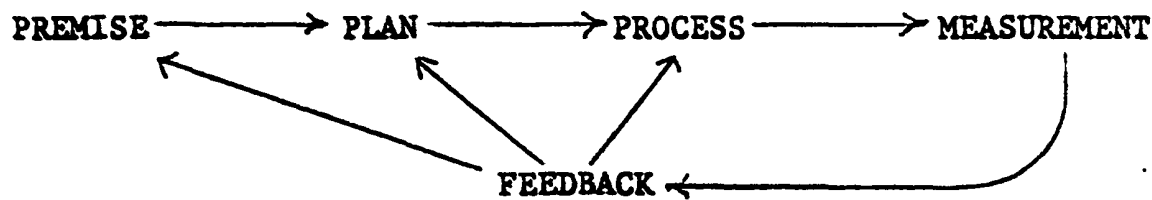
The benefits accruing from any management information system, whether computerized or not, may be seen to fall into three categories:

- a. Cost savings--i.e., savings in data processing cost because of reductions in clerical work force and other changes.
- b. Operational gains--i.e., efficiencies in corporate operations resulting from the application by managers of information received through the system (for instance, data on inventory reductions and faster production).
- c. Intangible benefits--i.e., improvement in customer service, corporate planning and forecasting, the ability to sustain growth, and other advantages which may not be present without the system but which depend upon management's astuteness in using it.

Thus, the functions for measuring the value of the present system were categorized into these three areas in an attempt to establish a baseline.

It is also necessary at this point to differentiate between the business system and the information system of the enterprise. For the most part, the business system has its main function of converting the raw materials, through a process of manufacturing, testing, and assembling, into oceanographic measuring instruments. The management information system, on the other hand, translates from the environment and from within its own components as its inputs. It stores this information and associates it with previously stored information in order to provide a frame of reference for the next vital step, that of making a decision.

Defined in this manner, the analysis of the present system and design of the proposed system were evaluated against the following premise. Management control is a function of measuring how well a task is being accomplished as compared to what was planned. This is pictorially described as:



4.0 THE EXISTING SYSTEM

As previously stated in the procedure of analysis section of this report, the existing system was evaluated primarily from the following points of view:

- a. What are the objectives or premise for operations?
- b. How was the premise put into operation, i.e., how did the planning take place?
- c. In what manner was the plan implemented?
- d. What was the measurement for determining how well the plan was being accomplished?
- e. What was the medium for feedback and how did it affect the process?

The first step was to define the processing system as it currently exists. EXHIBIT "E1" describes the system by means of a PERT type network as it was described by the employees within the system.

As we appraise the system from an overall standpoint, we find

Detail of Production

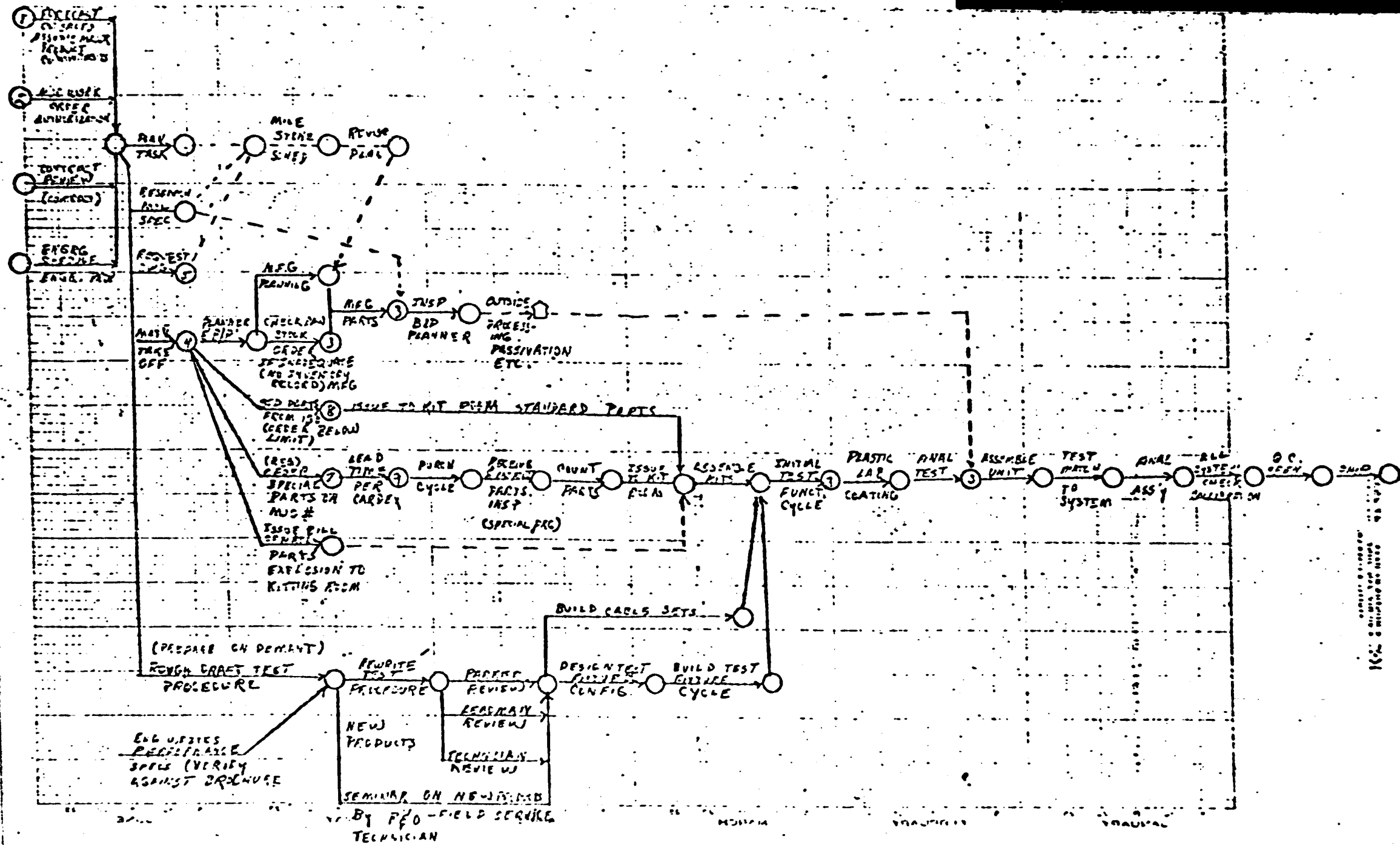


Figure SWAN 1

the information is usually defined in two documents, the MJO and contract. This information is then combined with the forecast of sales to plan the task. The planning task for inventory MJO's is completed and issued in the form of a milestone schedule, which is filed and never distributed. As a matter of interest, during the period of this investigation, much dissatisfaction was expressed by those making this schedule and they had discontinued its process and were in the process of designing a new one. No where in the organization was any other reference made of the milestone schedule nor did anyone seem to miss it. This was somewhat overcome by the authorization for work document. This document does give the information as to what is to be done and some limiting criteria, such as budget and when due.

In pursuing the authorization for work as the planning and control mechanism, we found its primary use was to provide the major task definition and a charge number by which to accumulate costs. It contains no provision for stating how the job is to be done and if there is a change in the plan. For instance, there is no revising of this document to reflect the change. Hence, the document doesn't really provide the planning function nor the feedback medium for control.

To further pursue the process of planning, we found the foreman in each of the departments was familiar with what had to be accomplished, based upon past experience, and assigned the work accordingly. This work assignment was related to their area only and reflected none of the interface requirements of the other subsystems. Thus, the product test plans were in no way coordinated with the machine shop. Again, there was a complete lack of any formal feedback mechanism. It would appear to an outsider that the system primarily reacted to crises and was not

Procedure Flow Chart

R. W. SWANSON, associates

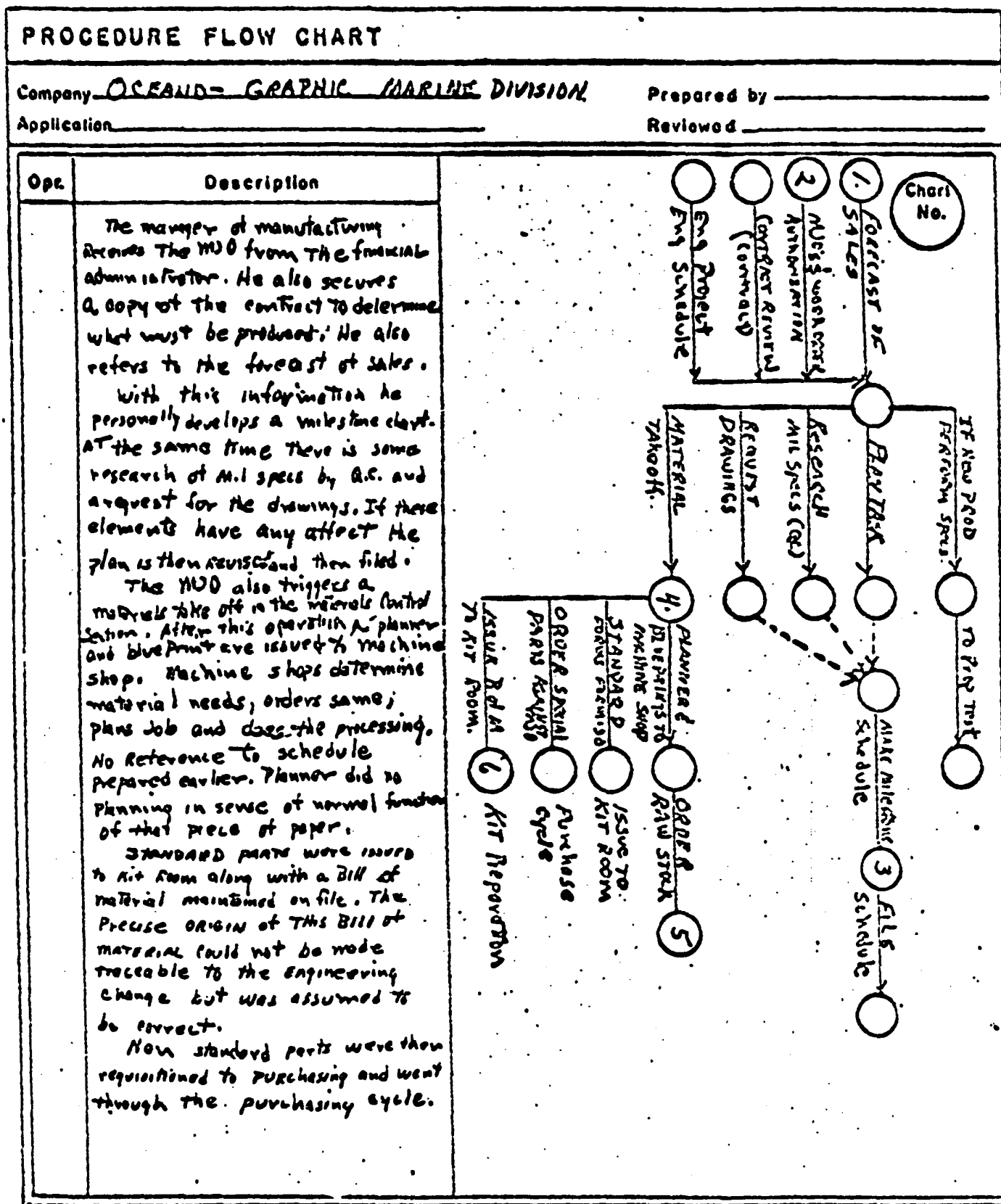


Figure SWAN 2

REPORT DISTRIBUTION and USE

Customer: OCEANO - GRAPHIC MARINE DIVISION

Date: 00/00/00

Use Value Comment Sheet

Process Frequency	Report Title	Unit	Report Distribution (Dept. Function, Person, Etc.)																										
			1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24			
			BUDGET	PURCHASING	FILE	SHIP	MKT.	DAD	ACST	ENG	G.C.	DEG. CONTRO	REC. CONT	R.F.D.	M.E.G.	INV. CONT	PROD. TEST	ELECT. ASSEMB	MACH. SHOP	OTHER	CONTRACTS	PH. Division	Rel. Div.	Memorial Delivery	Burst	Bind	Decollate	Total No. Copies	
D	WORK ORDER AUTHORIZATION	o																											
		b		✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓									16 MIN
		c																											20 MIN
		d																											
D	MJO	e		✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓										
		f																											
		g																											
		h																											

Notes

- WORK ORDERS CUT ON AN AS NEEDED BASIS. APPROXIMATELY 450 PER YEAR WITH AN AVERAGE OF 75 BEING OPEN AT ONE TIME. COST/VALUE STUDY NOT MADE AT PRESENT TIME. NOT CONSIDERED AS AN ACCOUNTABLE COST OR A DIRECT COST.
- MJO TRIGGERS WORK ORDER AUTHORIZATION - ONE IS REDUNDANT.

Processing Frequency

D - Daily M - Monthly
 W - Weekly O - Quarterly
 B - Bi-Weekly SA - Semi-Annually
 S - Semi-Monthly A - Annual

KEY

Burst -
 Bind
 Reg Left - L
 Top - T
 Right - R
 Decollate - /

Figure SWAN 4

R. W. SWANSON, associates

VALUE COMMENT			
Ref. No.	Rep. Use Lic.	Report Title	Value Narrative
F-2	1	FORECAST OF SALES	<p>THE forecast is of value to some executives at management, but is of little value to Production-planning. No commitment to production can be made on the basis of a forecast. MDR is the only correlation factor between the issue of the WORK AUTHORIZATION AND THE SALES FORECAST. THERE ARE NO STATISTICAL CORRELATIONS COMPLETED to establish any measure of validity. In my discussions with respect to production planning under present system, Report has only an IMMEDIATE value, if any.</p>
F-2	2	MDR & WORK ORDER AUTHORIZATION	<p>BASICALLY three documents. THE MDR IS THE ONLY documented evidence of planning that appears to tie all operations together. Document contains Customer identification, funding, Budget data and approvals. The AUTHORIZATIONS PART DOES identify the major functions to be accomplished. However it is department oriented and NOT process oriented. Budget is not related to time nor work station. May or may not reflect estimate.</p>

*Consider at least: (1) Earning on freed capital. (2) Yield on inventory reductions. (3) Benefits. (4) Equipment. (5) Displaceable costs. (6) Profit on increased sales.

Figure SWAN 5

Request For Master Job Order

OCEANO - GRAPHIC AIRLINE DIVISION

1. REQUEST FOR MASTER JOB ORDER

Initiator _____ Security _____ Contract # _____
 Project Mgr _____ Priority _____ P. O. # _____
 Project Engr. _____ RFP # _____

Customer Name Address/Job Title	Type of Project:	
	Contract/P. O.	<input type="checkbox"/>
	Mfg. for Inventory	<input type="checkbox"/>
	Construction of Asset	<input type="checkbox"/>
	Bidding	<input type="checkbox"/>
	Basic or Applied Research	<input type="checkbox"/>
	B. B. C. Repair	<input type="checkbox"/>

Job Description

See attached

Contract Funding		Target or Budget	
Cost	_____	Labor	_____
Fee or Profit	_____	Overhead	_____
Total	_____	Pur. Dir. Mat'l	_____
Contractual	_____	Co. Const. Mat'l	_____
Completion Date	_____	Other Direct	_____
Purchase Orders or Subcontracts may be placed against this MJO on or after _____		Total	_____
Contracts Approval _____		G & A	_____
		Total Costs	_____
		Profit	_____
		Profit %	_____
		Engr. O/H _____ %	Mfg. O/H _____ %
		G&A _____ %	

Approved:

Head of Budgets _____ Plant/Project Manager _____ Vice President _____ Pres/Exec. Vice President _____

2. MASTER JOB ORDER

MJO No. _____ Change No. _____ Cross Reference _____
 Effective Date _____ Expiration Date _____

Ship to: _____ Billing Instructions: _____
 Taxable Non-Taxable
 FOB Point _____
 Mark For: _____ Terms _____ % Net _____ Days _____
 Via: _____

Reports and Other Requirements including QC

See attached

Scheduled Completion Date	Scheduled Shipping Date
Distribution <input type="checkbox"/> Budgets <input type="checkbox"/> File <input type="checkbox"/> Mkt. <input type="checkbox"/> Acct. <input type="checkbox"/> Eng. <input type="checkbox"/> Mfg. Control <input type="checkbox"/> R&D	
<input type="checkbox"/> Pur. <input type="checkbox"/> Ship <input type="checkbox"/> D&D <input type="checkbox"/> S. A. T. <input type="checkbox"/> QC <input type="checkbox"/> Doc. Control <input type="checkbox"/> Mfg.	

Approved:

Vice President, Administration

(2)

Figure SWAN 6

R. W. SWANSON, associates

VALUE COMMENT*		
Est. No.	Rep. Title	Value Narrative
F-2 4	PRODUCTION PLANNING OPERATION	<p>DOCUMENT CONTAINS THE SIGNIFICANT IDENTIFICATION DATA. Content in order as a planning document is worthless. In executed form, manufacturing operations bear no resemblance to process of manufacture. TOOL & TIME STANDARDS are not maintained nor relevant. Clerk in inventory control stated she made planner. Her interpretation of making planner consists of reproducing a copy from a notice planner previously reproduced. Operation bears no correlation to present production volume or Requirements document could disappear and have no effect upon the system. IF this operation were to be redesigned to function correctly. Inventory costs alone could be reduced from current \$85,000 level to a \$20,000 level easily. This could be accomplished with no impact to production requirements. Recommend a Buy whats needed type of system.</p>

* Consider at least (1) Earning on freed capital. (2) Yield on inventory reductions. (3) Benefits. (4) Equipment. (5) Displaceable costs. (6) Profit on increased sales.

Figure SWAN 8

a function of discerning which of many alternatives should be taken. Unfortunately, it is the empty wagon that makes the most noise.

Because of the lack of a formal flow of the planning function, it became impossible to measure economically the cost of that system and we, therefore, abandoned attempting to establish that of baseline. We also recognize that the Marine Division is basically an intermittent type of operation or job shop--organized about multipurpose machines to perform specific functions. This type of environment normally relies on heavy paperwork in order to control closely its cost and production capacities.

We believe the present management has designed systems that would accomplish the intent of control. However with the dynamics, particularly present in the oceano-graphic field, any effort to accomplish this manually with a sizable production volume would be futile. First, the cost of manpower, primarily clerical and management, would be economically infeasible. Second, the response time would be of such length as to negate most of the competitive advantages of the product and seriously affect the function of sales.

For the purposes of this report, we have included some of the analytical data used to arrive at the above conclusion. This, of course, is only a sample of the effort and volume of data collected. It is not our intent to burden this report with data that is meaningful only to fellow colleagues in this field. However, some insight should be given the reader as to the logic of the conclusions.

5.0 PROPOSED SYSTEM

Experience has shown that the gathering and dissemination of

information is the manufacturing company's most difficult problem. Information is voluminous, scattered, and often difficult to obtain. In one of the early stages of this study, Mr. Fictitious emphatically stated he wanted a system that would give him "access to the facts." In addition, if we are to maintain the function of management control, we must provide for a means of planning and measurement that is currently lacking authenticity in the current system. Thus, our system must have a central information system and a framework that will facilitate mechanization.

We are proposing a single flow information system where the information will be accumulated in a production control center and where, literally, one set of books is maintained. This takes the form of records stored on computer disk files, readily accessible to all interested parties at a moment's inquiry. These records are designed to contain as much data as is deemed important to management. The accuracy of the records, also an essential element, will be easier to achieve and maintain, for only one set of archives will now be put to use.

The essential element of this system is the development of a data base that generally covers all the operational information needed to handle this company's business. It will be stored on disk files and, therefore, directly on-line with a computer. Because of this, summary and detail information can be accessed, updated, and retrieved from multiple entry points. Data is stored one way through reference to symbolic record field labels and can be printed in various output formats.

Each system of records is linked in a particular manner. For example, a part number, accessed through the basic record, may lead to

what we call the product structure--a where-used going-to file, or standard routing record or the manufacturing sequence, or an open order status, or an open job summary, or a detail record. In the latter instance, the specific work center in which the job is being performed may even be pinpointed.

5.1 THE SYSTEM OVERVIEW

Drawing No. 1, System Overview, shows the interaction on data flow within the Marine Division organization. The interaction of these events is grouped within seven major areas:

- a. Sales analysis
- b. Engineering
- c. Inventory control and production scheduling
- d. Manufacturing facilities
- e. Finance
- f. Purchasing
- g. Sales and distribution

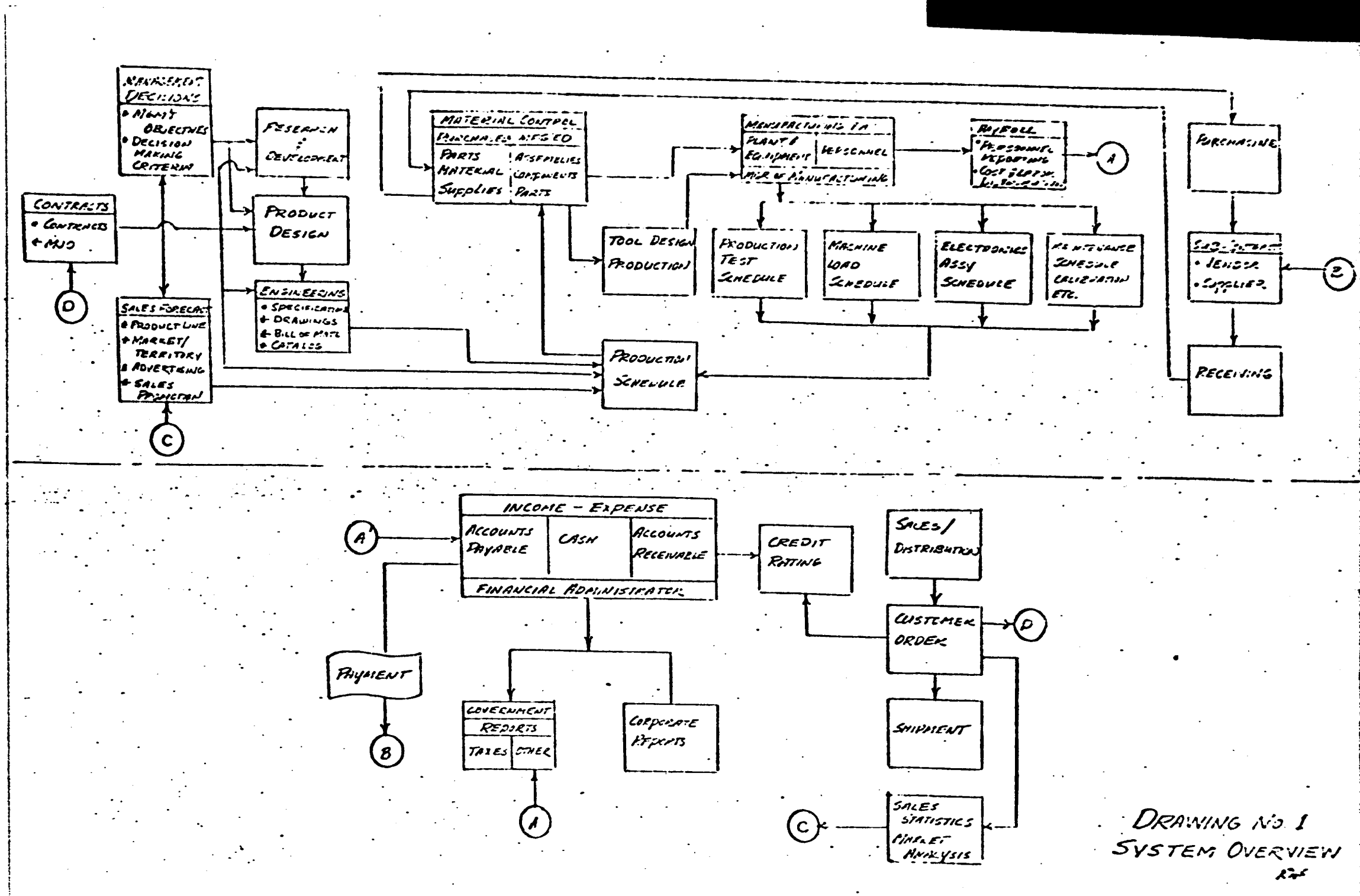
While the data base is designed to process segments of all these areas, it does not contain the additional detail records needed to handle sales analysis, finance, and sales/distribution.

5.2 PRIMARY FLOW OF INFORMATION

The input of information leads from an initial input of customer orders and statistical sales background data to the final shipment of an order. A generalized statement of the system flow is divided into a planning phase and an execution phase.

Planning begins with the preparation and projection of order forecasts. Stock availability and on-order status are screened across

System Overview



DRAWING NO. 1
SYSTEM OVERVIEW
RWF

Figure SWAN 10

product inventory records, but family component characteristics of the product line must also be recognized. Product structure on bills of material enter into these decisions.

After a determination of net requirements, an order quantity analysis takes place to ascertain lot sizes and lead times for both purchased and manufactured items.

To-buy items are routed to where items are placed on a purchase requisition. At this point a selection of vendor is made, price and delivery are negotiated, and purchase order is released. Receipt cards and/or a scheduled receipt document may be prepared simultaneously with the purchase order and forwarded to the inspection-receiving area of the plant. An open purchase order record is now initiated for follow-up.

To-make items are routed to production planning for assembly and fabrication. Some similarity exists within these two units. An assembly order is generated for the assembly area, a shop order for the fabrication area. Material requisitions and job tickets accompany both documents. Three basic types of records (standard routing, work center load, and open job order) permit assembly and fabrication to schedule, to load, and to level the line or shop and to release the order paperwork.

Execution begins at the purchasing level with the need for order follow-up and vendor expediting. The vendor ships material, accompanying his shipment with packing lists and an invoice.

Varied execution functions are performed at the assembly and fabrication levels. Orders are dispatched, rescheduled, and expedited between work centers. In the meantime, current production reporting

updates work center and open job order records.

5.3 SYSTEM FLOW

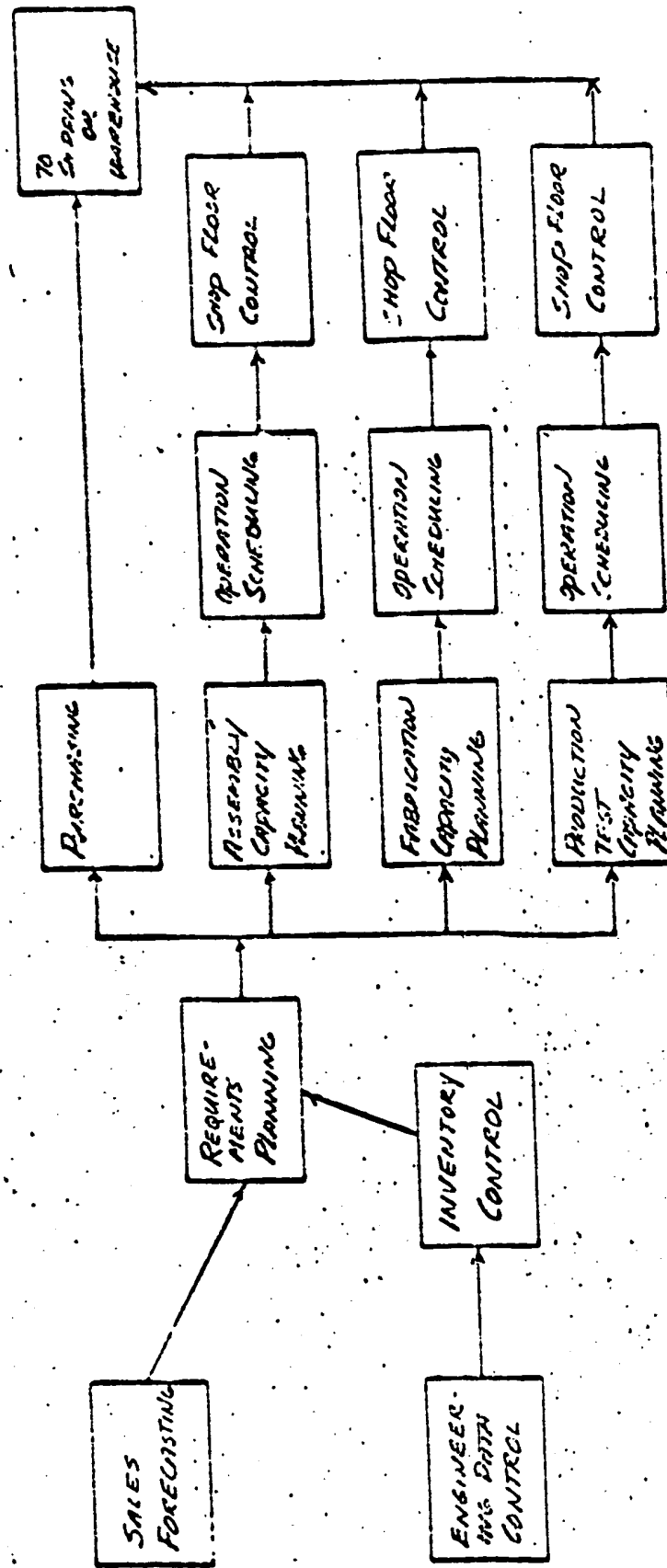
Our proposed system has been designed to fit this basic production model. For ease of implementation, eight subsystems have been developed. This allows for modular programming and should substantially reduce the cost of implementation and speed of accomplishment. The information flow begins from two directions. (See Drawing No. 2, Production Model.) The first path moves from engineering data control, to inventory control, to requirements planning. The mission of the engineering data control subsystem is to organize and maintain basic records. These basic records are what we call the item master file, product structure file, standard routing file, and work center master. The subsystem has the added capacity or capability of retrieving information from the data base. Six retrieval functions are presently developed, three in assembly sequence and three in parts usage sequence.

An inventory control subsystem follows organization of basic records. On-hand inventory, usage history, and on-order fields are used in the item master file so that stock status reports can be generated. Thus, a major objective of this application area is record maintenance and updating inventory. With accessibility to such data, "when to order" and "how much to order" decisions are made.

A second flow line moves from sales forecasting to requirements planning. The sales forecasting subsystem analyzes historical demand data, which may be stored on the item master, to provide requirements planning with a gross finished product forecast plan.

The merger of requirements planning with inventory control now

PRODUCTION MODEL



DATA BASE RECORDS

- ENG DATA CONTROL
 - ITEM MASTER
 - PRODUCT STRUCTURE
 - STANDARD ROUTING
 - WORK CONTROL MASTER
 - OPEN JOB ORDER
 - ORDER SUMMARY
 - SALES FORECASTING
 - ITEM MASTER
-
- INV. CONTROL
 - ITEM MASTER
 - REQUIREMENTS PLANNING
 - ITEM MASTER
 - PRODUCT STRUCTURE
-
- PURCHASING
 - ITEM MASTER
 - PURCHASE MASTER
 - VENDOR MASTER
 - OPEN PURCHASE
 - CAPACITY PLANNING
 - ITEM MASTER
 - STANDARD ROUTING
 - WORK CENTER MASTER
-
- OPERATION SCHEDULING
 - STANDARD ROUTING
 - WORK CENTER MASTER
 - OPEN JOB ORDER
 - TOOL MASTER
-
- SHOP FLOOR CONTROL
 - ITEM MASTER
 - STANDARD ROUTING
 - PRODUCT STRUCTURE
 - OPEN JOB ORDER

DRAWING NO 2
AMS

Figure SWAN 11

makes it possible to determine net requirements, projected into time periods and schedule due dates. Product structure records are used at this point to allow breakdown of finished product items into individual components. These are similarly netted and projected into time periods. All of this results in planned orders, destined to each of the four links: purchasing, electronic assembly, fabrication, and product test.

Planned orders to purchasing result in material requisitions being prepared. Through the use of purchase master and vendor master records, a vendor may be selected and a purchase order with receiving documents created so that purchase follow-up can be initiated in the next sequence of events.

Planned orders to the three production areas go to the capacity planning sub-systems, or long-range scheduler. Its purpose is to identify overloads far enough in the future for both facility and manpower planning. After order start date calculations are performed (utilizing standard routing records), consideration is given to plant capacity. The work center master is used for this purpose. Available techniques are then used to level the loads. A work center load report, projected by time period, is one of the key output documents. The operation scheduling accepts orders which have gone through a releasing cycle from capacity planning and schedules the work center within its short-range time span. Dispatching sequences are prepared and analyses made of the loads. Priority rules are set and order completion dates determined. To the short-range scheduling phase of this subsystem, we have added the control of tools. A tool master record is designed for this function.

Shop floor control is the final subsystem in the flow line. It prepares the shop packets and other factory documentation. It also

constructs the open job summary and operation detail records so that the progress of the work can be reported. Feedback is one of its more important functions so that the system can respond to change.

5.4 STANDARDIZED RECORDS

A set of standardized record layouts is included in this report. These records are designed as the data base to mechanize the application areas we have discussed and lead toward developing a single information flow concept.

These records contain the fields we consider necessary to enable Oceano-Graphic Division to utilize and control their control production output requirements. Each record is described in detail in Drawings numbered three through eight. Their respective lengths determine the size of the storage requirements of any data processing equipment employed. Chart No. 9 is included to show how the chaining sequence of the records is interfaced. This interaction of the records achieve for us two major functions, first to provide access to the facts and secondly, through this chaining sequence a single data base is developed to allow for implementation of the single information flow showing how a change, such as an engineering change, updates the whole data base, which is common to all. Thus, we accomplish one of our main objectives; one set of books is maintained.

5.5 BENEFITS OF THE PROPOSED SYSTEM

a. A Plan for Growth

A plan can be developed to begin implementation for each of the application subsystems leading to the single information flow system. The system can grow as Oceano-

File Layout

RECORDS AND WORK AREAS

110.3

RECORD NAME	ITEM MASTER RECORD																			
HOW SIGNED																				
FIELD DEFINITION	ITEM			PRODUCT STRUCTURE										STANDARD FIXTURE				FILLER		
CHARACTER POSITION	TYPE	NUMBER	DESCRIPTION	FORECAST	UNIT OF MEAS	INV. VALUE CLASSIFICATION	1 ST ASSEMBLY ADDRESS	2 ND ASSEMBLY ADDRESS	LOW LEVEL CODE	NEXT ITEM IN ACT. CHAIN	OVERFLOW CHAIN ADDRESS	1 ST OPERATION ADDRESS	LAST OPERATION ADDRESS							
FIELD LENGTH	2	5	8	2	4	9	8	8	3	8	8	8	8							
SN	5	8	N	4	9	5	8	8												

RECORD NAME	ORDER POLICY																			
HOW SIGNED																				
FIELD DEFINITION	ORDER POLICY										FORECASTING									
CHARACTER POSITION	ORDER CODE	ORDER POINT	ORDER QTY/UP-TO	SAFETY STOCK	MINIMUM	MAXIMUM	MULTIPLE	MODEL TYPE	FIRST AVERAGE	SECOND AVERAGE	TREND	SAFETY FACTOR	AVERAGE DEMAND	MEAN ABSOLUTE DEVIATION	SUM OF DEVIATIONS	BASE PERIOD	START PER FACTOR			
FIELD LENGTH	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2

RECORD NAME	LEAD TIME																			
HOW SIGNED																				
FIELD DEFINITION	LEAD TIME			RAW MATL		UNIT COST						UNIT PRICE			PARTS WARE HISTORY					
CHARACTER POSITION	PURCH ASSESS	PRODUCTION	SET UP	NO	UNIT QTY PER PART	STANDARD COSTS			ACTUAL COSTS			LIST	NET	TAX CODE	DEMAND		ISSUES			
FIELD LENGTH	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	

RECORD NAME	CURRENT PERIOD																			
HOW SIGNED																				
FIELD DEFINITION	CURRENT PERIOD				INVENTORY ON HAND						ALLOCATED	BACK ORDER	PHYSICAL INVENTORY							
CHARACTER POSITION	BEHAVING INVENT	TRANSFER & ACQUIRED	REQUIS	ISSUE	DEMAND	TOTAL QTY	NO LOCATION	PRIMARY LOCATION		ADDRESS TO PHYSICAL LOCATION	QUANTITY	QUANTITY	TYPE OF INV. (FINANCIAL-RELATION) ETC		QUANTITY COUNT	CHECKED NO	DATE OF LAST COUNT	DATE OF NEXT COUNT		
FIELD LENGTH	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	

RECORD NAME	PROJECTED SALES																			
HOW SIGNED																				
FIELD DEFINITION	PROJECTED SALES			ON ORDER		ON HAND		ADDRESS TO PURCHASER		ADDRESS TO VENDOR		ENGINEERING CHANGE CONTROL								
CHARACTER POSITION	DATE	ADDRESS	QUANT	DATE	QUANT	DATE	QUANT	ADDRESS	ADDRESS	ADDRESS	ADDRESS	LAST FACILITY CHANGE		CURRENT ENGINEERING CHANGES						
FIELD LENGTH	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	2	

Figure SWAN 12

Chaining Concept

NO. 4

OTHER FILES TO BE DESIGNED

- PRODUCT STRUCTURE FILE
- PEGGED REQUIREMENTS
- OPEN PURCHASE REQUISITIONS
- OPEN PURCHASE ORDERS
- VENDOR MASTER
- PURCHASE MASTER
- OPEN JOB ORDER SUMMARY
- STANDARD ROUTING
- WORK CENTER MASTER
- OPEN JOB OPERATION DETAIL
- TOOL MASTER

CHAINING CONCEPT

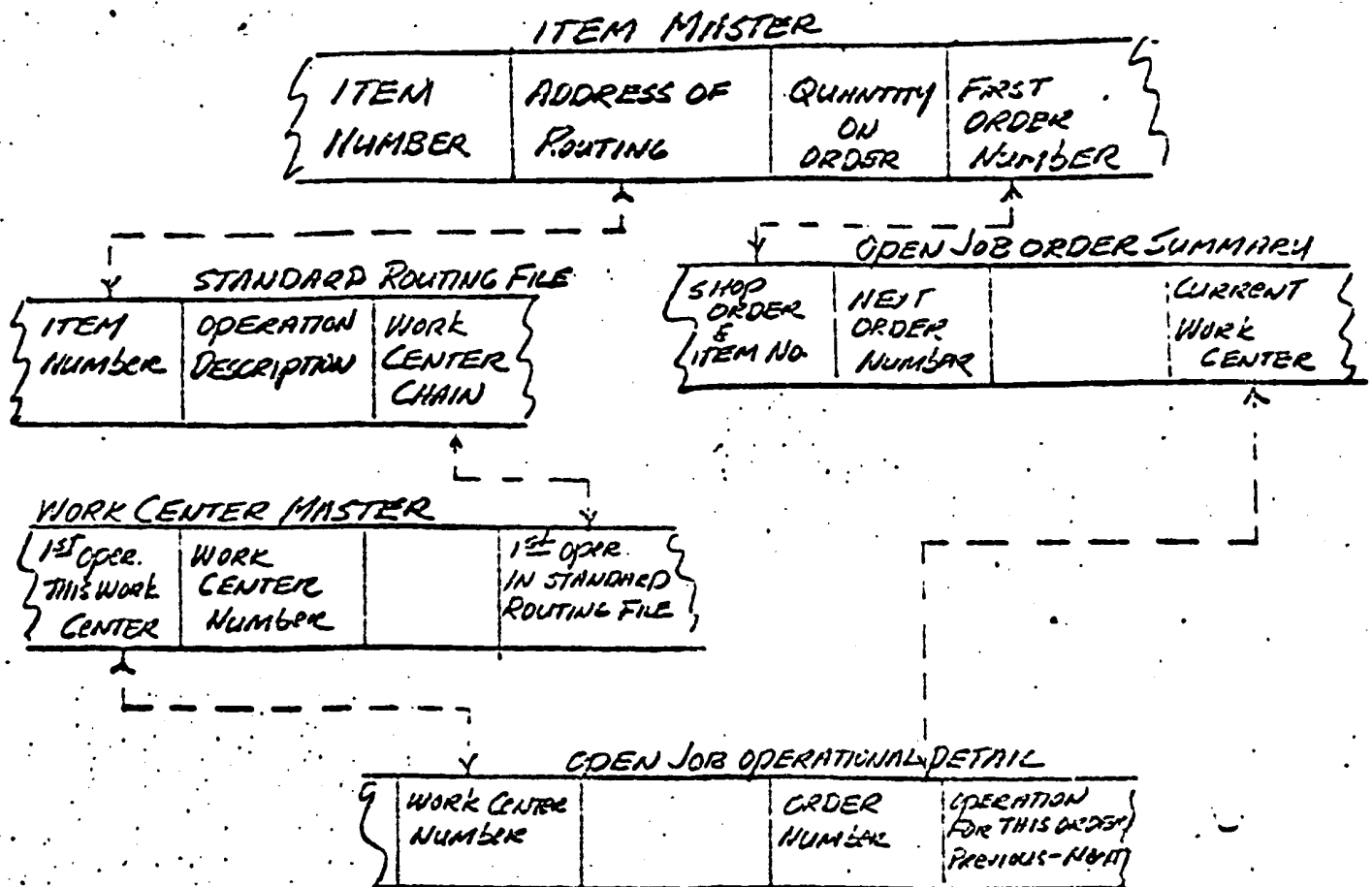


Figure SWAN 13

Procedure Flow Chart

R. W. SWANSON, associates

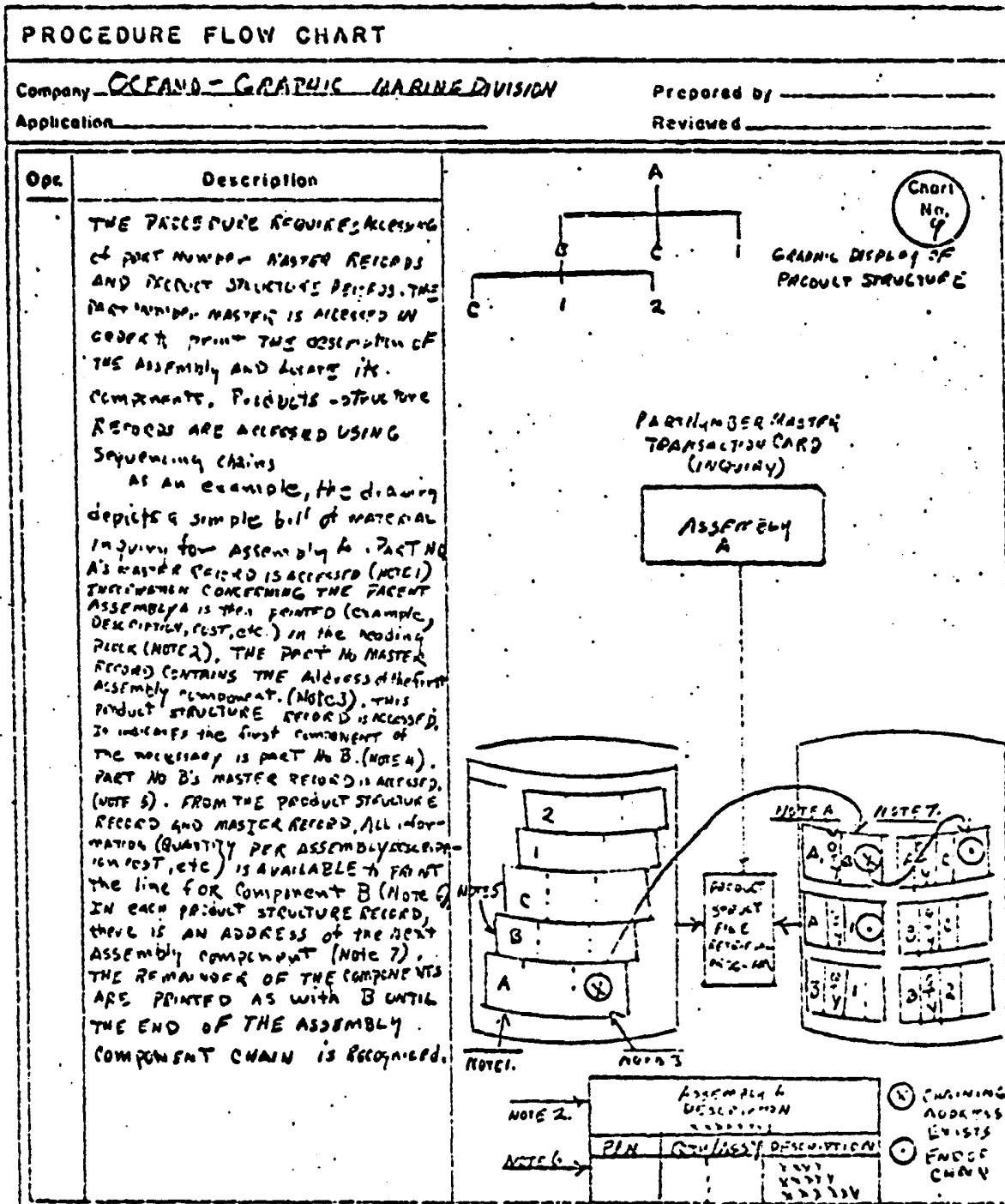


Figure SWAN 14

Graphic grows. And Oceano-Graphic will obtain tangible results long before the total system is installed.

b. Standardization

Information is common to all--it is known and called by one description. One set of records makes this possible. Thus, it makes planning easier with standardized consistency in estimating and pricing.

c. Extensive Data Base

The record base has two important features: accessibility and accuracy. Information is accessible through inquiry to multiple points; detail is available through chaining to all related records. No longer will it be necessary to spend hours or days searching file drawers or ledger cards. Also, information is more accurate; it is updated in only one place. Standard transactions processed within each subsystem assure complete record maintenance.

d. Modular Program Design

Chart No. 10 shows the modular programming concepts. We can obtain tangible results before the total system is installed. This has been the major drawback of the single information flow concept. This will also allow a systematic plan of implementation with easily defined milestones for time and budget control.

e. All production information is now directed into a single channel. Levels of operating and management personnel are made more aware. With this assurance, the following should happen:

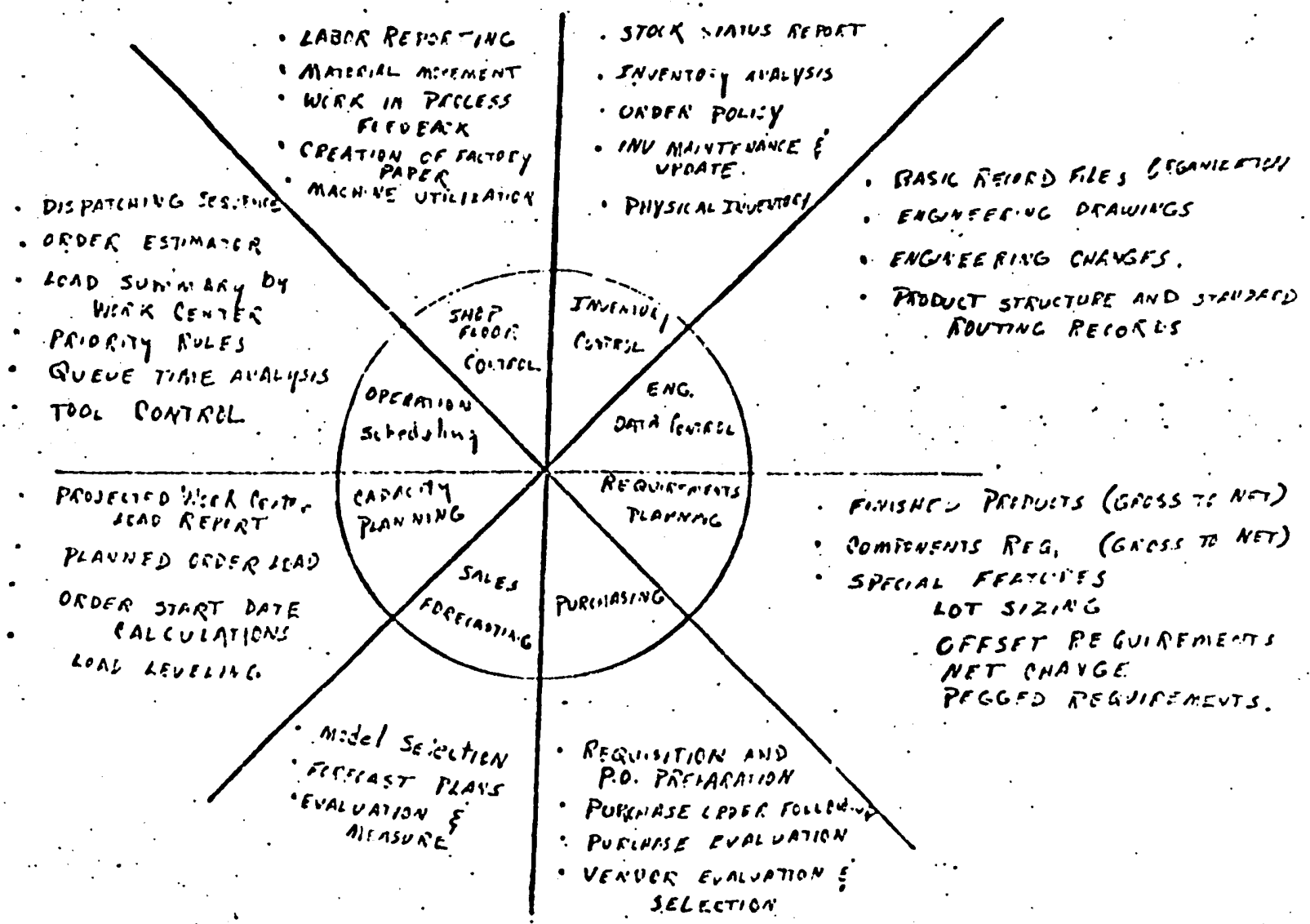


Figure SWAN 15

- (1) Cost can be closely controlled, better surveillance over overtime hours, inventory, and machines--the key to cost reduction.
- (2) More efficient planning.
- (3) More time available to react to changes.
- (4) Less waste, reduced information costs, more profits.

6.0 COST OF IMPLEMENTATION

The cost of implementation is a function of two variables. Machine costs of equipment capable of the level of mechanization required for this system represent one variable. Software cost, the cost of programming, represents a one-time cost necessary to the operation of the system.

6.1 HARDWARE COSTS

The hardware costs have not, for the purposes of this report, been precisely defined by the soliciting of quotes or bid proposals. However, with our experience directly in this field, a realistic estimate can be submitted. For instance, an IBM System 000/000 with a configuration to handle this application would cost \$0,000 per month through the first years of operation. The addition of terminals as the need arises will increase the hardware costs to approximately \$0,000 per month at later stages of development.

Operating and programming costs after the system is fully in operation will be approximately \$000 per month. It is imagined this capability would be developed from within the current organization and no additional people added to the payroll.

6.2 SOFTWARE COSTS

The normal rule of thumb for predicting software cost for a new system such as this is to take the first year's cost for the hardware; in this case, \$00,000. However, the sophistication of this system, which is considerably more elaborate than a payroll system, etc. (which is usually the makeup of software costs), can be accomplished for considerably less investment, since developed, canned programs, programs already written, can be used to minimize this expenditure. However, the canned programs still must be modified and patched to fit the Oceano-Graphic operation. We, therefore, estimate the cost of programming, debugging, and implementation of the system to be \$00,000, approximately \$0,000 per month for a period of ten months (Reference Drawing No. 11, Implementation Plan).

ACTIVITIES

Question 1. The existing system was evaluated from what points of view?

Answer:

- A. What are the objectives or premise for operation?
- B. How was the premise put into operation, i.e., how did the planning take place?
- C. In what manner was the plan implemented?
- D. What was the measurement for determining how well the plan was being accomplished?
- E. What was the medium for feedback and how did it affect the process?

Question 2. What was the object of the engineering data control?

Answer: To organize and maintain basic records.

Question 3. What are the basic records?

Answer: Item Masters, Product Structure, Standard Routing and Work Center Master.

Question 4. What other record is included under Engineering Data Control?

Answer: Open Job Order and Order Summary.

Question 5. What is the major objective of the Inventory Control System?

Answer: Record maintenance and updating of the inventory (it indicates when to order and how to order).

Question 6. What else is of importance as far as inventory control is concerned?

Answer: On-order fields are used in the Item Master so that Stock Status reports can be generated.

Question 7. What functions do the Sales Forecasting and Requirements Planning Connection perform?

Answer: The analization of historical demand data, which may be stored on the item master to provide requirements planning with a gross finished product plan.

Question 8. What is the result of the merger of Requirements Planning with Inventory Control?

Answer: This makes it possible to determine the net time requirements, projected into time periods and schedule due dates.

Question 9. What other records are used?

Answer: Product Structure records are used to allow the breakdown of finished product items into individual components.

Question 10. What are the names of the four links that the Planned Orders are distributed to?

Answer: Purchasing, Electronic Assembly, Fabrication and Production Test.

Question 11. What are the four files involved when the Planned Orders reach the purchasing link?

Answer: Item Master, Purchase Master, Vendor Master and Open Purchase Order.

Question 12. How may a vendor be selected?

Answer: Through the use of a Purchase Master and a Vendor Master record.

Question 13. What is the purpose of the Capacity Planning Subsystem?

Answer: It identifies overloads far enough in the future to allow manpower planning and facility planning.

Question 14. What records are used to calculate the Order Start Date?

Answer: Standard Routing Records.

Question 15. Give the name of one of the key documents developed in this area according to the Production Model.

Answer: Work Center Load Report, projected by time period.

Question 16. What is the purpose of the Operation Scheduling Section?

Answer: It accepts orders which have gone through a releasing cycle from Capacity Planning, and schedules the work center within its short range time span.

Question 17. Shop Floor Control is the final subsystem in the flow. Give two of its functions.

Answer: A. It prepares the shop packets and other factory documentation.

B. It constructs the Open Job Summary and Operating

Detain Records so that the progress of the work can be reported.

Question 18. If we examine the set of Standardized Layout Records, we note that certain fields were used. Why?

Answer: The fields were used to allow Oceano-Graphic to utilize and control their production output requirements.

Question 19. What can the lengths of the records help determine?

Answer: The lengths can determine the size of storage requirements of any data processing equipment used.

Question 20. Chart 9 shows how chaining sequence of the record is interfaced. What two functions does this achieve?

Answer:

- A. This allows the development of the necessary retrieval functions to provide access to the facts.
- B. The chaining (or linking) allows the development of a single data base which means a single information concept.

(Make a check mark for your answer to Data Management Topics--included, not included, or irrelevant in the Swanson Study.)

FEASIBILITY STUDY

1. Basic Steps in Development of Business Systems

1. Problem definition
2. Application Research
3. Scope of study
 - a. Organizational boundaries
 - b. Objective of study
 - c. Resources available for study
4. Objectives
5. Target dates
6. Study phase responsibility
7. Education of several departments
8. Management role

2. Factors Included in Systems Report to Management

1. Introduction
2. Description of study
3. Costs
4. Organizational changes

3. Systems and Procedures Covered

4. Functions of Systems and Procedures Department

1. Systems analysis and design
2. Forms design and control
3. Records management and retention
4. Report analysis
5. Preparation of written procedures
6. Work distribution
7. Process flow chart
8. Procedure flow chart
9. Work measurement
10. Time standards
11. Time and motion study
12. Forms control
 - a. Retain only necessary forms
 - b. Renew and reuse
 - c. Proper design and manufacture
 - d. One person assigned to forms control

5. Forms Design Basic Consideration

1. Purpose
2. Decision for specific information

3. Use and purpose
4. Logical placement of items
5. Physical placement of items
6. Analyzation of number of copies

6. Records Management

1. Creation and use of records
2. Distribution
3. Simplification of paperwork
4. Data retained
5. Filing procedures
6. Security

7. Company Manuals

1. Organizational
2. Policy
3. Operation
4. General information

8. Installation of Electronic Data Processing

1. Systems design
2. Personnel selection and training
3. File conversion
4. Forms design
5. Programming
6. Testing of programs
7. Facilities design and preparation
8. Scheduling and testing the system
9. Changeover
10. Feedback evaluation

9. Elements of Good Reports

1. Accurate
2. Clear
3. Relevant
4. Current

10. File Design

1. Volume consideration
2. Timing
3. Sorting

11. Fact Gathering Techniques of a System Study

1. Interview
2. Questionnaire
3. Review Records

12. Type of Information Gathered

1. Historical
2. Cost
3. Current procedure
4. Effectiveness
5. Relationship between departments

CONSTRUCTION

The Swanson Study uses a PERT type plan with Chart 11. The construction program can be used with this type of plan, and the cost program, which follows can be used to account for costs that occur daily.

Use the tape cassette to analyze this program. It will develop an introduction to the Basic Assembly Language used in this book.

The object of this program is to find the shortest time in which a construction project can be completed and to find the activities that limit the time for completion.

First a network for the project, such as Figure C-1 is drawn. Each arrow represents an activity. The numbered circles represent events. The numbers alongside the arrows are the work time, in days, for the activity. The diagram is laid out such that all activities leading to an event must be completed before any activities leading away from that event can begin. The ending number must be higher than the beginning event number.

The data is read into the computer, one activity per IBM card. The first sixteen spaces are reserved for the activity description, the next three spaces for the starting event number, then three spaces for the end event number, finally three spaces for the work time.

The data is stored in core storage, first the sixteen bytes of description, then two bytes for start event (packed format is used for all number data in the storage tables), two bytes for end event number, two bytes for normal work time. In addition, space is reserved in memory for start time, end time, slack time and an index. Two bytes of core storage are reserved for each activity in the table.

The program proceeds:

1. Initialize
2. Print header
3. Read all data cards and store information in table
4. Print out table
5. Sort the table on basis of lowest starting event number
6. Locate activity with the lowest starting event number
7. Locate and place in table TTWO all activities with end event numbers that match the start event numbers of (6)
8. Sort table TTWO on the basis of ending times
9. Subtract ending time of each activity from the ending time of the largest activity in table TTWO
10. Enter difference in slack time position
11. Enter the longest event time from table TTWO into start time for activity (6)
12. Compute end time for activity (6)
13. Index activity (6)
14. Transfer table TTWO and activity (6) back into main table
15. Test for last event
16. NO-back to (6)
17. Print out main table for second time
18. Sort main table on basis of lowest slack time first
19. Cut table off after a last event with no slack time
20. Sort new short table on basis of start event number, highest first
21. Print headers
22. Print out first activity in table (this activity is actually the last activity to be accomplished on the project)

23. Index first activity with a %%
24. Locate and print all activities with end event numbers the same start event as 21
25. Index each activity printed with a caret
26. Return to beginning of table, find first activity not indexed with %% , but indexed with a caret
27. Locate and print all activities with end event numbers the same as start event number 25
28. At last event end program

Construction--Flow Chart

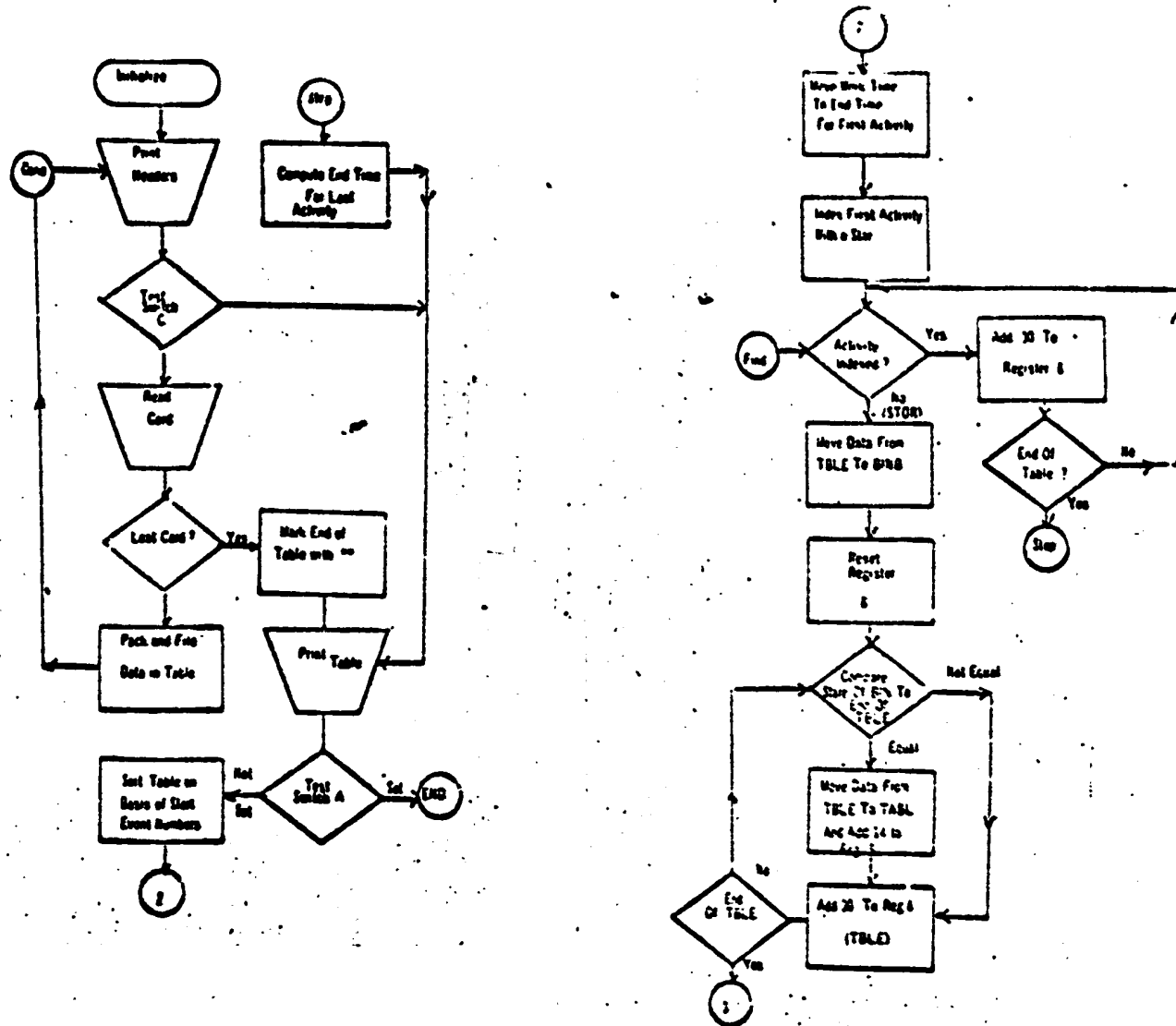


Figure FC-02

Construction--Flow Chart

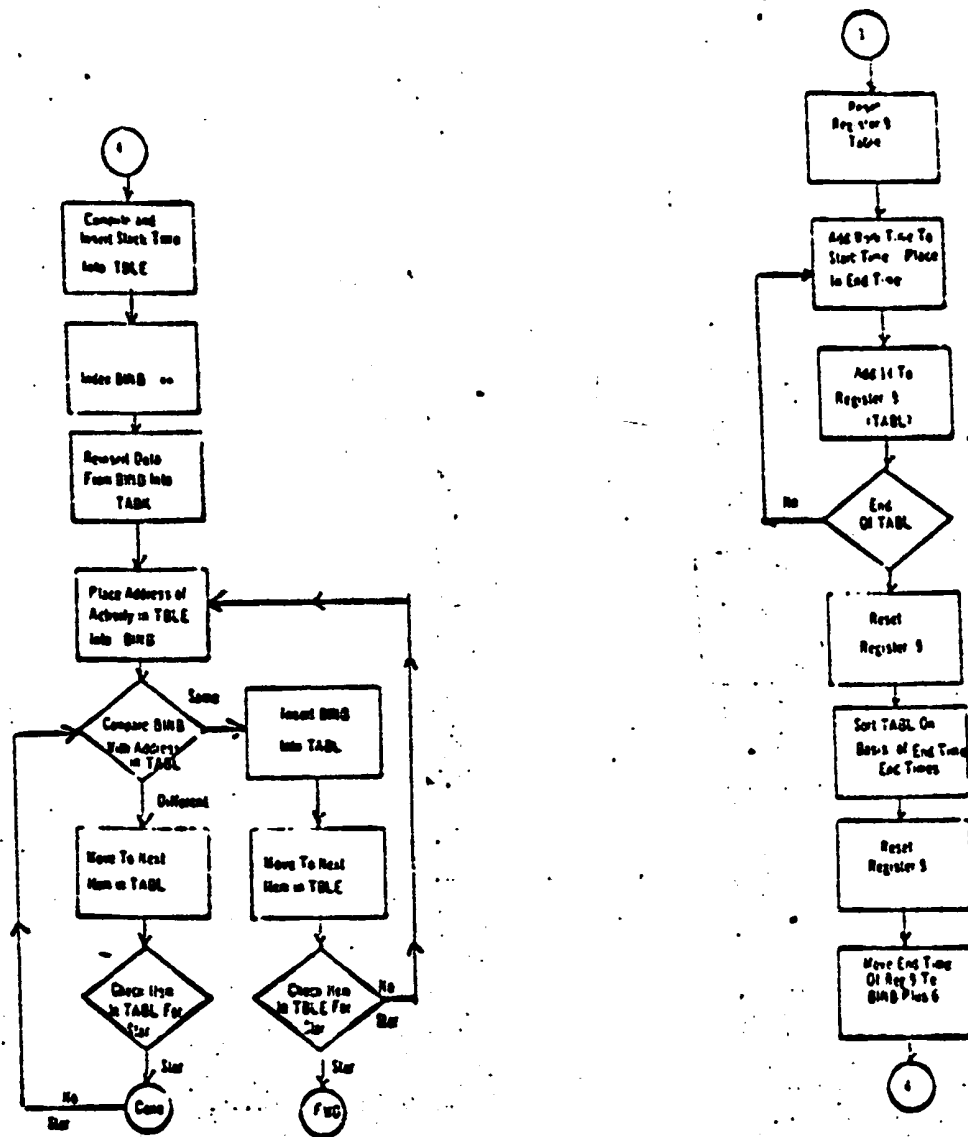


Figure FC-03

Construction--Flow Chart

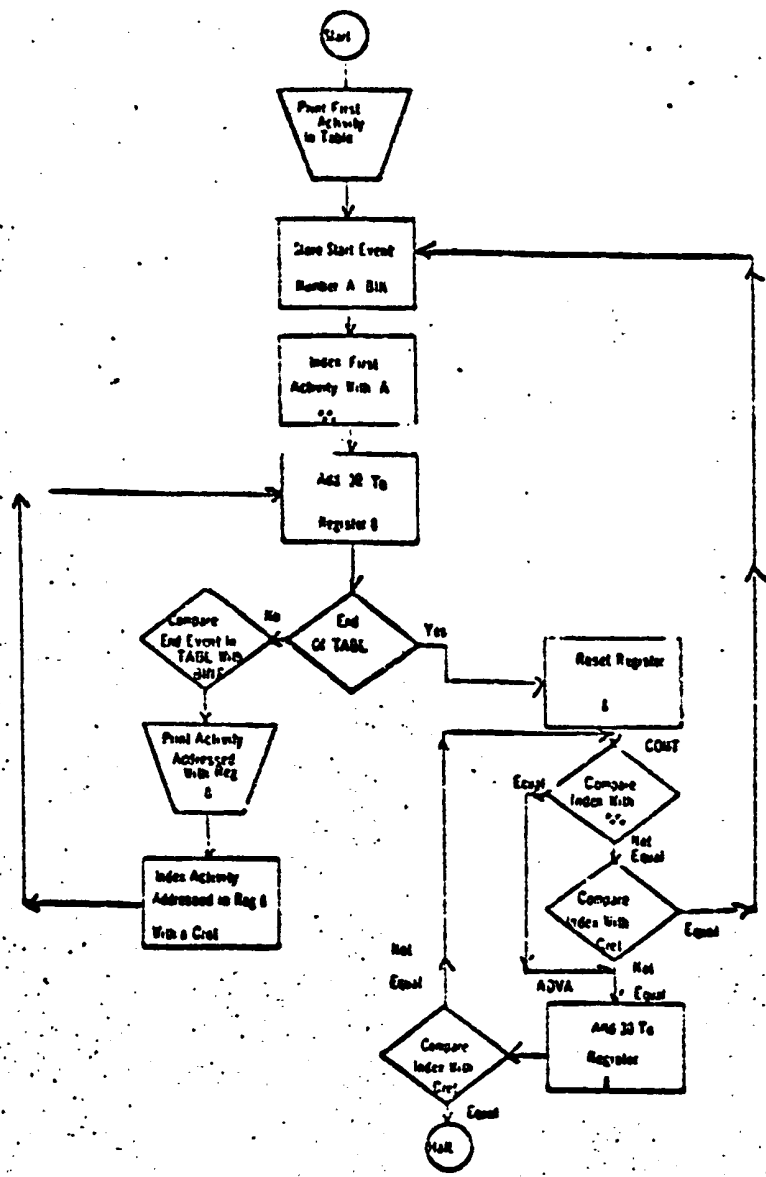


Figure FC-01

Construction Program--Network Layout

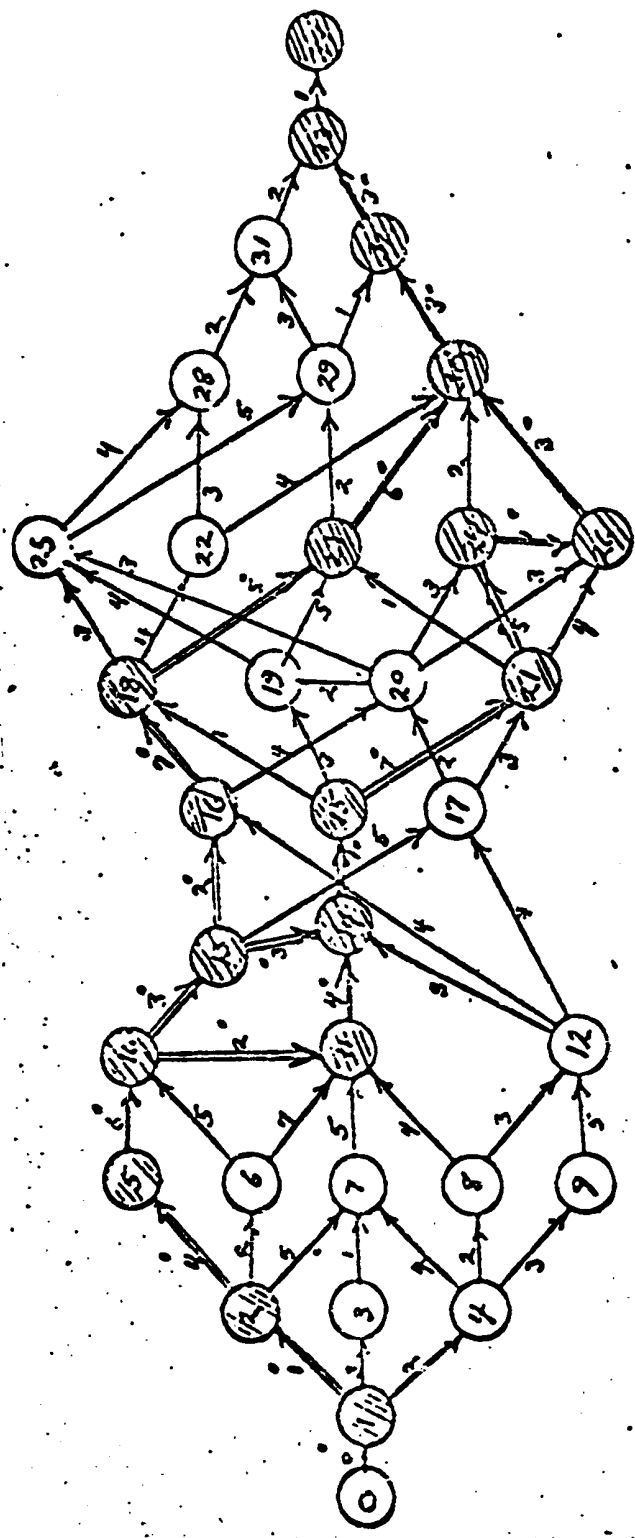


Figure NET-01

Construction Program

Sort Routine

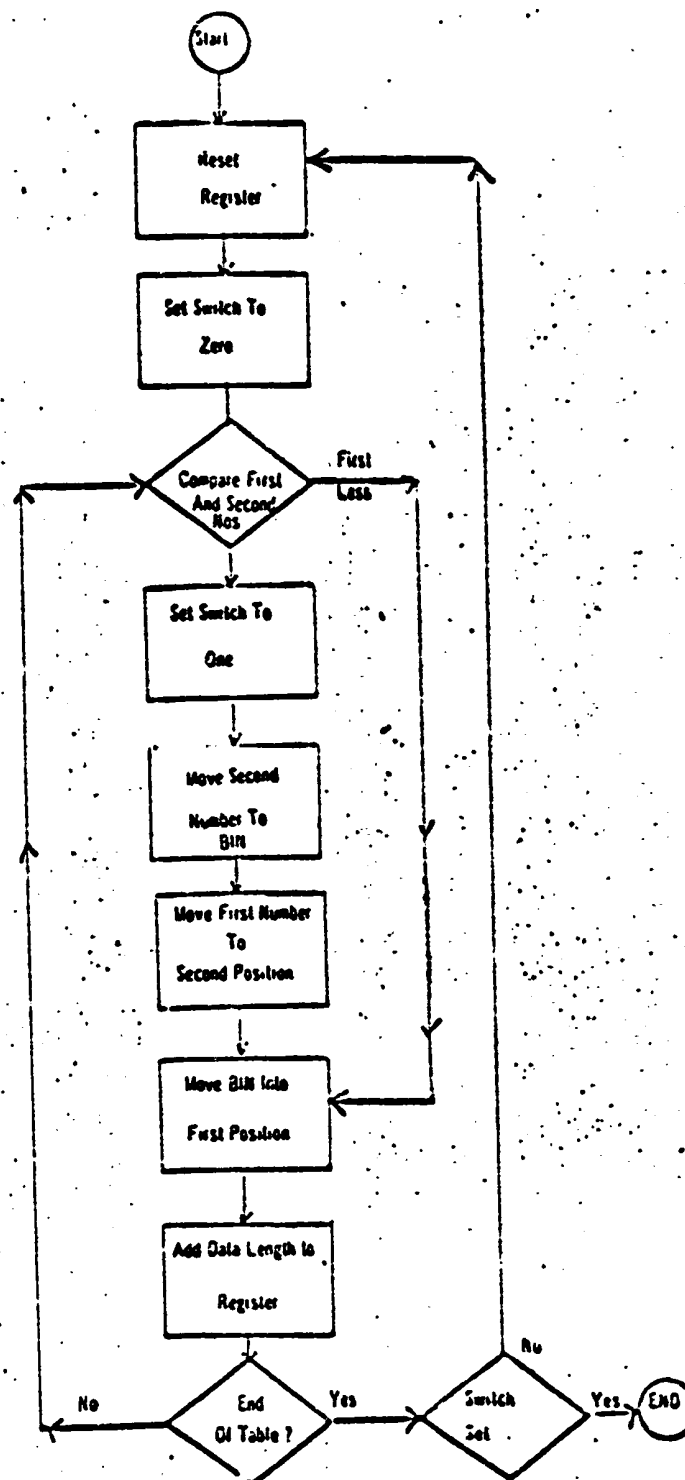


Figure SRT-01

Construction Program

0244	0510 0010 001E		MVC	0130,01,301A)	006
0246	0510 001E 0003		MVC	30170,01,010	006
0248	0400 1500		AM	0,0101	006
024A	0501 0011 E50A		CLC	3017,01,STAR	006
024C	0770 000E		RC	7,0101	006
024E	0101 050E		TM	SW,X'01'	006
0248	0770 0000		RC	1,0101	006
024A	1000		SR	0,0	006
024C	0400 0000		LM	0,0101	007
024E	1000		SA	0,0	007
0272	0400 0000		LM	0,0101	007
0274	0411 001A 0010		ZAP	24(2,0),20(2,0)	007
0276	0411 001A 000A		ZAP	26(2,0),ZERO	007
0278	0501 001C E50A		MVC	20(2,0),STAR	007
027A	0501 001C E50A		CLC	20(2,0),STAR	007
027C	0770 010E		RC	7,0101	007
027E	0400 1500		AM	0,0101	007
0280	0501 0000 E50A		CLC	0(2,0),STAR	007
0282	0770 010E		RC	0,0101	007
0284	0770 0110		RC	15,0101	007
0286	0701 0001 0010		MVC	0100(1),10(0)	009
0288	1000		SR	0,0	009
028A	0400 0000		LM	0,0101	009
028C	0411 0012 0001		CP	10(2,0),010(2)	009
028E	0770 0110		RC	7,0101	009
0290	0200 0000 0010		MVC	01(4,0),010(0)	009
0292	0400 0000 0000		MVC	01(4,0),010(0)	009
0294	0400 0000 0000		AM	0,0101	009
0296	0400 0000 0000		AM	0,0101	009
0298	0501 0000 E50A		CLC	0(2,0),STAR	009
029A	0770 0110		RC	7,0101	009
029C	0201 0000 E50A		MVC	0(2,0),STAR	009
029E	1000		SR	0,0	009
02A0	0400 0000		LM	0,0101	009
02A2	0411 0000 0000		ZAP	ACCA,0(2,0)	009
02A4	0411 0000 0000		ZAP	ACCA,0(2,0)	009
02A6	0400 0000 0000		AM	0(2,0),ACCA	009
02A8	0400 0000 0000		AM	0,0101	009
02AA	0501 0000 E50A		CLC	0(2,0),STAR	009
02AC	0770 0110		RC	7,0101	009
02AE	0201 0000 E50A		MVC	0(2,0),STAR	009
02B0	1000		SR	0,0	009
02B2	0400 0000		LM	0,0101	009
02B4	0411 0000 0000		ZAP	ACCA,0(2,0)	009
02B6	0411 0000 0000		ZAP	ACCA,0(2,0)	009
02B8	0400 0000 0000		AM	0(2,0),ACCA	009
02BA	0400 0000 0000		AM	0,0101	009
02BC	0501 0000 E50A		CLC	0(2,0),STAR	009
02BE	0770 0110		RC	7,0101	009
02C0	0201 0000 E50A		MVC	0(2,0),STAR	009
02C2	1000		SR	0,0	009
02C4	0400 0000		LM	0,0101	009
02C6	0411 0000 0010		ZAP	ACCA,0(2,0)	009
02C8	0411 0000 0010		ZAP	ACCA,0(2,0)	009
02CA	0400 0000 0000		AM	0(2,0),ACCA	009
02CC	0400 0000 0000		AM	0,0101	009
02CE	0501 0000 E50A		CLC	0(2,0),STAR	009
02D0	0770 0110		RC	7,0101	009
02D2	0201 0000 E50A		MVC	0(2,0),STAR	009
02D4	1000		SR	0,0	009
02D6	0400 0000		LM	0,0101	009
02D8	0411 0000 0000		ZAP	ACCA,0(2,0)	009
02DA	0411 0000 0000		ZAP	ACCA,0(2,0)	009
02DC	0400 0000 0000		AM	0(2,0),ACCA	009
02DE	0400 0000 0000		AM	0,0101	009
02E0	0501 0000 E50A		CLC	0(2,0),STAR	009
02E2	0770 0110		RC	7,0101	009
02E4	0201 0000 E50A		MVC	0(2,0),STAR	009
02E6	1000		SR	0,0	009
02E8	0400 0000		LM	0,0101	009
02EA	0411 0000 0010		ZAP	ACCA,0(2,0)	009
02EC	0411 0000 0010		ZAP	ACCA,0(2,0)	009
02EE	0400 0000 0000		AM	0(2,0),ACCA	009
02F0	0400 0000 0000		AM	0,0101	009
02F2	0501 0000 E50A		CLC	0(2,0),STAR	009
02F4	0770 0110		RC	7,0101	009
02F6	0201 0000 E50A		MVC	0(2,0),STAR	009
02F8	1000		SR	0,0	009
02FA	0400 0000		LM	0,0101	009
02FC	0411 0000 0010		ZAP	ACCA,0(2,0)	009
02FE	0411 0000 0010		ZAP	ACCA,0(2,0)	009
0300	0400 0000 0000		AM	0(2,0),ACCA	009
0302	0400 0000 0000		AM	0,0101	009
0304	0501 0000 E50A		CLC	0(2,0),STAR	009
0306	0770 0110		RC	7,0101	009
0308	0201 0000 E50A		MVC	0(2,0),STAR	009
030A	1000		SR	0,0	009
030C	0400 0000		LM	0,0101	009
030E	0411 0000 0010		ZAP	ACCA,0(2,0)	009
0310	0411 0000 0010		ZAP	ACCA,0(2,0)	009
0312	0400 0000 0000		AM	0(2,0),ACCA	009
0314	0400 0000 0000		AM	0,0101	009
0316	0501 0000 E50A		CLC	0(2,0),STAR	009
0318	0770 0110		RC	7,0101	009
031A	0201 0000 E50A		MVC	0(2,0),STAR	009
031C	1000		SR	0,0	009
031E	0400 0000		LM	0,0101	009
0320	0411 0000 0010		ZAP	ACCA,0(2,0)	009
0322	0411 0000 0010		ZAP	ACCA,0(2,0)	009
0324	0400 0000 0000		AM	0(2,0),ACCA	009
0326	0400 0000 0000		AM	0,0101	009
0328	0501 0000 E50A		CLC	0(2,0),STAR	009
032A	0770 0110		RC	7,0101	009
032C	0201 0000 E50A		MVC	0(2,0),STAR	009
032E	1000		SR	0,0	009
0330	0400 0000		LM	0,0101	009
0332	0411 0000 0010		ZAP	ACCA,0(2,0)	009
0334	0411 0000 0010		ZAP	ACCA,0(2,0)	009
0336	0400 0000 0000		AM	0(2,0),ACCA	009
0338	0400 0000 0000		AM	0,0101	009
033A	0501 0000 E50A		CLC	0(2,0),STAR	009
033C	0770 0110		RC	7,0101	009
033E	0201 0000 E50A		MVC	0(2,0),STAR	009
0340	1000		SR	0,0	009
0342	0400 0000		LM	0,0101	009
0344	0411 0000 0010		ZAP	ACCA,0(2,0)	009
0346	0411 0000 0010		ZAP	ACCA,0(2,0)	009
0348	0400 0000 0000		AM	0(2,0),ACCA	009
034A	0400 0000 0000		AM	0,0101	009
034C	0501 0000 E50A		CLC	0(2,0),STAR	009
034E	0770 0110		RC	7,0101	009
0350	0201 0000 E50A		MVC	0(2,0),STAR	009
0352	1000		SR	0,0	009
0354	0400 0000		LM	0,0101	009
0356	0411 0000 0010		ZAP	ACCA,0(2,0)	009
0358	0411 0000 0010		ZAP	ACCA,0(2,0)	009
035A	0400 0000 0000		AM	0(2,0),ACCA	009
035C	0400 0000 0000		AM	0,0101	009
035E	0501 0000 E50A		CLC	0(2,0),STAR	009
0360	0770 0110		RC	7,0101	009
0362	0201 0000 E50A		MVC	0(2,0),STAR	009
0364	1000		SR	0,0	009
0366	0400 0000		LM	0,0101	009
0368	0411 0000 0010		ZAP	ACCA,0(2,0)	009
036A	0411 0000 0010		ZAP	ACCA,0(2,0)	009
036C	0400 0000 0000		AM	0(2,0),ACCA	009
036E	0400 0000 0000		AM	0,0101	009
0370	0501 0000 E50A		CLC	0(2,0),STAR	009
0372	0770 0110		RC	7,0101	009
0374	0201 0000 E50A		MVC	0(2,0),STAR	009
0376	1000		SR	0,0	009
0378	0400 0000		LM	0,0101	009
037A	0411 0000 0010		ZAP	ACCA,0(2,0)	009
037C	0411 0000 0010		ZAP	ACCA,0(2,0)	009
037E	0400 0000 0000		AM	0(2,0),ACCA	009
0380	0400 0000 0000		AM	0,0101	009
0382	0501 0000 E50A		CLC	0(2,0),STAR	009
0384	0770 0110		RC	7,0101	009
0386	0201 0000 E50A		MVC	0(2,0),STAR	009
0388	1000		SR	0,0	009
038A	0400 0000		LM	0,0101	009
038C	0411 0000 0010		ZAP	ACCA,0(2,0)	009
038E	0411 0000 0010		ZAP	ACCA,0(2,0)	009
0390	0400 0000 0000		AM	0(2,0),ACCA	009
0392	0400 0000 0000		AM	0,0101	009
0394	0501 0000 E50A		CLC	0(2,0),STAR	009
0396	0770 0110		RC	7,0101	009
0398	0201 0000 E50A		MVC	0(2,0),STAR	009
039A	1000		SR	0,0	009
039C	0400 0000		LM	0,0101	009
039E	0411 0000 0010		ZAP	ACCA,0(2,0)	009
03A0	0411 0000 0010		ZAP	ACCA,0(2,0)	009
03A2	0400 0000 0000		AM	0(2,0),ACCA	009
03A4	0400 0000 0000		AM	0,0101	009
03A6	0501 0000 E50A		CLC	0(2,0),STAR	009
03A8	0770 0110		RC	7,0101	009
03AA	0201 0000 E50A		MVC	0(2,0),STAR	009
03AC	1000		SR	0,0	009
03AE	0400 0000		LM	0,0101	009
03B0	0411 0000 0010		ZAP	ACCA,0(2,0)	009
03B2	0411 0000 0010		ZAP	ACCA,0(2,0)	009
03B4	0400 0000 0000		AM	0(2,0),ACCA	009
03B6	0400 0000 0000		AM	0,0101	009
03B8	0501 0000 E50A		CLC	0(2,0),STAR	009
03BA	0770 0110		RC	7,0101	009
03BC	0201 0000 E50A		MVC	0(2,0),STAR	009
03BE	1000		SR	0,0	009
03C0	0400 0000		LM	0,0101	009
03C2	0411 0000 0010		ZAP	ACCA,0(2,0)	009
03C4	0411 0000 0010		ZAP	ACCA,0(2,0)	009
03C6	0400 0000 0000		AM	0(2,0),ACCA	009
03C8	0400 0000 0000		AM	0,0101	009
03CA	0501 0000 E50A		CLC	0(2,0),STAR	009
03CC	0770 0110		RC	7,0101	009
03CE	0201 0000 E50A		MVC	0(2,0),STAR	009
03D0	1000		SR	0,0	009
03D2	0400 0000		LM	0,0101	009
03D4	0411 0000 0010		ZAP	ACCA,0(2,0)	009
03D6	0411 0000 0010		ZAP	ACCA,0(2,0)	009
03D8	0400 0000 0000		AM	0(2,0),ACCA	009
03DA	0400 0000 0000		AM	0,0101	009
03DC	0501 0000 E50A		CLC	0(2,0),STAR	009
03DE	0770 0110		RC	7,0101	009
03E0	0201 0000 E50A		MVC	0(2,0),STAR	009
03E2	1000		SR	0,0	009
03E4	0400 0000		LM	0,0101	009
03E6	0411 0000 0010		ZAP	ACCA,0(2,0)	009
03E8	0411 0000 0010		ZAP	ACCA,0(2,0)	009
03EA	0400 0000 0000		AM	0(2,0),ACCA	009
03EC	0400 0000 0000		AM	0,0101	009
03EE	0501 0000 E50A		CLC	0(2,0),STAR	009
03F0	0770 0110		RC	7,0101	009
03F2	0201 0000 E50A		MVC	0(2,0),STAR	009
03F4	1000		SR	0,0	009
03F6	0400				

Construction Program

0344	4710 E14E	BC	1,SUR2	010
034E	1A99	SR	9.9	011
0350	4A90 ERCC	LM	9.TIND	011
0354	F811 EAA4 900A	ZAP	TTIM(2),R(2,9)	011
035A	F811 EA67 EAA4	ZAP	BINH*(2),TTIM	011
0360	F811 EAA6 EAA4	ZAP	BINE,TTIM	011
036A	F811 EAA6 900A	SP	BINE,R(2,9)	011
036C	F811 900A E886	ZAP	10(2,9),BINE	011
0372	4A90 EATE	AM	9.FTEN	011
037A	D501 9000 E56A	CLC	0(2,9),STAR	011
037C	4770 E20A	BC	7.TOTI	011
0380	D200 EAA0 E56A	MVC	BINH*12,STAR	011
0386	1A8A	SR	P.A	012
038A	4A80 E996	LM	A,THEG	012
039C	F911 A010 E861	CP	16(2,8),BINH(2)	012
039E	4770 E254	BC	7.GOOD	012
0396	F911 A012 E863	CP	18(2,8),BINH*2(2)	012
039C	4770 E254	BC	7.GOOD	012
03A0	D200 A010 E861	MVC	16(14,A),BINH	012
03A4	47F0 E25C	BC	15.INST	012
03A4	4A80 E56A	AM	A,THTY	012
03AE	47F0 E236	BC	15.HOPE	012
03B2	1A99	SR	9.9	012
03B4	4A90 ERCC	LM	9.TIND	012
03BA	1A8A	SR	A.A	012
03BA	4A80 E996	LM	A,TAEG	012
03BE	D200 EA61 9000	MVC	BINH(14),0(9)	013
03C4	F911 A010 EA61	CP	16(2,A),BINB(2)	013
03CA	4770 E24A	BC	7.NOGO	013
03CE	F911 A012 E863	CP	18(2,A),BINB*2(2)	013
03D4	4770 E2AA	BC	7.NOGO	013
03DA	D200 A010 E861	MVC	16(14,A),BINH	013
03DE	4A90 EATE	AM	9.FTEN	013
03E2	D501 9000 E56A	CLC	0(2,9),STAR	013
03EA	4780 E29A	BC	A.FIND	013
03EC	47F0 E262	BC	15.GOGO	013
03F0	1A8A	SR	A.A	013
03F2	4A80 E996	LM	A,TAEG	013
03F6	1899	SR	9.9	014
03FA	4A90 ERCC	LM	9.TIND	014
03FC	47F0 E132	BC	15.FIND	014
0400	4A80 E56A	AM	A,THTY	014
0404	D501 8000 E56A	CLC	0(2,A),STAR	014
040A	4780 E28C	BC	A,STEP	014
040E	47F0 E268	BC	15.GOGI	014
0412	9A01 EAA3	DI	SMC,X'01'	014
0416	4A80 E56A	SM	A,THTY	014
041A	F811 EAA0 8014	ZAP	ACCA,20(2,8)	014
0420	F811 EAA0 8016	AP	ACCA,22(2,8)	014
0426	F811 8018 E880	ZAP	24(2,8),ACCA	014
042C	F811 801A E8C6	ZAP	26(2,8),ZERO	015
0432	9B45 0001	CTO	1(0),X'45'	015
0436	4780 E004	BC	6.CONE	015
043A	4780 E20C	BC	4.SKIP	015
043E	9800 0078	HPR	X'878',0	015
		SUM		
		TOTI		
		HOPE		
		GOON		
		INST		
		GOGO		
		GOGI		
		FIND		
		NOGO		
		STEP		
		SKIP		

Figure CON-03

Construction Program

0002	0710 120C	RC	15.SKIP	015
0004	1AAA	SR	R,R	015
0006	0000 E994	LM	R,TRIG	015
0008	020E 007A 0000	MVC	11(MI1A),0'0)	015
0010	022F 0000 0000	MVC	DATA,MASK	015
0012	025A 0000 0010	MVC	DATA(12),10(0)	015
0014	012F 0000 0000	ED	DATA,DATA	015
0016	0277 0000 E9AD	MVC	PRT(120),PRT-1	016
0018	0277 0000 E05F	MVC	PRT(120),M000	016
0020	0000 0000	AS	10,PRT	016
0022	0000 0000	AM	R,THY	016
0024	0001 0000 E90A	CLC	0(2,A),STAR	016
0026	0270 0000	RC	7,INT0	016
0028	0101 0000	TM	SW,X'01'	016
0030	0100 0000	RC	R,SNRT	016
0032	1AAA	SR	R,R	016
0034	0000 E996	LM	R,TRFC	016
0036	0000 E9AC	NI	SW,X'00'	016
0038	0011 001A 0030	CP	24(2,A),50(2,0)	016
0040	0700 E34E	RC	13,COM0	016
0042	0001 E9AC	DI	SW,X'01'	017
0044	0210 0000 0000	MVC	0(130),0(0)	017
0046	0210 001E E043	MVC	0(30,A),30(0)	017
0048	0000 E9AA	MVC	30(30,0),0IN	017
0050	0001 001E E90A	AM	R,THY	017
0052	0770 E33E	CLC	30(2,A),STAR	017
0054	0101 E9AC	BC	7,LOPB	017
0056	0710 E330	TM	SW,X'01'	017
0058	1AAA	BC	1,SUR3	017
0060	0000 E996	SR	R,R	017
0062	0011 001A ER06	LM	R,TRFC	017
0064	0770 E33C	CP	26(2,A),Z000	017
0066	0000 E9AA	RC	7,STOP	018
0068	07F0 E37A	AM	R,THY	018
0070	0201 0000 E90A	RC	15,COM0	018
0072	1AAA	MVC	0(2,0),STAR	018
0074	0000 E996	SA	R,R	018
0076	0000 E9AC	LM	R,TRFC	018
0078	0011 0010 002E	NI	SW,X'00'	018
0080	0700 E33C	CP	16(2,A),40(2,0)	018
0082	0001 E9AC	BC	11,COM5	018
0084	0210 0000 0000	DI	SW,X'01'	018
0086	0210 001E E043	MVC	0(130),0(0)	018
0088	0000 E9AA	MVC	0(30,0),30(0)	018
0090	0001 001E E90A	MVC	30(30,A),0IN	019
0092	0770 E33C	AM	R,THY	019
0094	0101 E9AC	CLC	30(2,0),STAR	019
0096	0710 E302	BC	7,LUPC	019
0098	0000 0001	TM	SW,X'01'	019
0100	0700 E302	BC	1,SUR4	019
0102	0000 0323	BC	1(0),X'45'	019
0104	0700 E302	BC	0,PRR	019
0106	0000 0323	BC	0,SKIT	019
0108	07F0 E302	MFR	X'323',0	019
		BC	5,SKIT	019

Figure CON-04

Construction Program

051C	D277 E50E E50D	PRA	MVC	PRT(120),PRT-1	019
0542	D277 E50E E7C7		MVC	PRT(120),MOR5	019
054A	4040 E4E9		MAS	10,PRNT	020
054C	4040 E4D6		BAS	10,SPAC	020
0550	D277 E50E E50D		MVC	PRT(120),PRT-1	020
0552	D277 E50E E6D7		MVC	PRT(120),MOR2	020
055C	4040 E4E9		MAS	10,PRNT	020
05A0	D277 E50E E50D		MVC	PRT(120),PRT-1	020
05A2	D277 E50E E76F		MVC	PRT(120),MOR3	020
05AC	4040 E4E9		MAS	10,PRNT	020
0570	4040 E4D6		BAS	10,SPAC	020
0574	18AA		SR	A,A	020
0576	4A80 E996		LM	A,TRIG	020
057A	D20F E67B 8000		MVC	ITEM(16),0(16)	020
0580	D22F E6AA E8AA		MVC	DATA,MASK	021
0586	D20B E6AA 8010		MVC	DATB(12),16(16)	021
05AC	DE2F E6AA E8AA		ED	DATA,DATB	021
0592	D277 E50E E50D		MVC	PRT(120),PRT-1	021
059A	D277 E50E E65F		MVC	PRT(120),MOR0	021
059E	4040 E4E9		MAS	10,PRNT	021
05A2	F911 E88B 8010	LOPD	ZAP	8,INE,16(2,8)	021
05A8	D201 801C E841		MVC	28(2,A),PERC	021
05AE	4A80 E5A8	LOPE	AM	8,TMTY	021
05B2	D501 8000 E50A		CLC	0(2,A),STAR	021
05BA	4770 E4C2		BC	7,PRNN	022
05BC	18AA		SR	8,A	022
05BE	4A80 E996		LM	8,TRIG	022
05C2	D501 801C E841	COMT	CLC	28(2,A),PERC	022
05CA	4780 E480		BC	A,ADVA	022
05CC	D501 801C E83F		CLC	28(2,A),CRET	022
05D2	4780 E4AC		BC	8,LUPD	022
05D6	4A80 E56A	ADVA	AM	8,TMTY	022
05D8	D501 8000 E50A		CLC	0(2,8),STAR	022
05E0	4780 E4CE		BC	A,MALT	022
05E4	47F0 E48C		BC	15,COMT	022
05EA	F911 8012 E88B	PRNN	CP	18(2,8),8,INE	022
05EE	4770 E45A		BC	7,LOPE	023
05F2	D20F E67B 8000		MVC	ITEM(16),0(16)	023
05FA	D22F E6AA E8AA		MVC	DATA,MASK	023
05FE	D20B E6AA 8010		MVC	DATB(12),16(16)	023
0604	DE2F E6AA E8AA		ED	DATA,DATB	023
060A	D277 E50E E50D		MVC	PRT(120),PRT-1	023
0610	D277 E50E E65F		MVC	PRT(120),MOR0	023
0616	4040 E4E9		MAS	10,PRNT	023
061A	D201 801C E83F		MVC	28(2,8),CRET	023
0620	47F0 E45A		BC	15,LOPE	023
0624	9900 0444	MALT	MFA	X'444',0	024
0628	47F0 E4CE		BC	15,MALT	024
062C	9844 0001	SPAC	CTO	1(0),X'444'	024
0630	07CA		BCA	8,10	024
0632	4740 E4D6		BC	4,SPAC	024
0636	9900 2111		MFA	X'111',0	024
063A	47F0 E4D6		BC	15,SPAC	024
063E	D040 E50E 0678	PRNT	XIO	PRT(X'40'),120	024
0640	4780 E4FE		BC	8,PBSV	024

Figure CON-05

Construction Program

0000	0700 1000	NC	4.PPNT	024	
0001	0000 0000	HPR	X'777'0	024	
0002	0700 0000	NC	15.PPNT	024	
0003	0000 0000	PBSY	T100	024	
0004	0000 0000	T100	PRFR,X'010	024	
0005	0000 0000	RCA	15.10	025	
0006	0000 0000	HPR	X'773'0	025	
0007	0000 0000	NC	15.PPNT	025	
0008	0000 0000	DC	C' '	025	
0009	0000 0000	CDIN	DS	OCLAO	025
0010	0000 0000	DFSC	DS	CL10	024
0011	0000 0000	STAA	DS	CL3	024
0012	0000 0000	ENIE	DS	CL3	024
0013	0000 0000	MINK	DS	CL3	024
0014	0000 0000	DS	DS	CL55	024
0015	0000 0000	STAB	DS	CL2	024
0016	0000 0000	ENDI	DS	CL2	024
0017	0000 0000	MHWY	DS	CL2	024
0018	0000 0000	INTY	DC	M'30'	024
0019	0000 0000	STAM	DC	C'000'	024
0020	0000 0000	SM	DS	CL1	027
0021	0000 0000	DC	DC	C' '	027
0022	0000 0000	PRT	DS	CL120	024
0023	0000 0000	SNA	DS	CL1	024
0024	0000 0000	NDR1	DS	CL27	024
0025	0000 0000	DC	DC	C'ANTELOPE VALLEY'	024
0026	0000 0000	DC	DC	C'CONSTRUCTION CO'	024
0027	0000 0000	DC	DC	C'COMPANY ACTIVITIE'	024
0028	0000 0000	DC	DC	C'S SCHEDULE L.0.0'	024
0029	0000 0000	DC	DC	C'NUSTER'	024
0030	0000 0000	DS	DS	CL27	024
0031	0000 0000	NDR4	DS	OCL120	024
0032	0000 0000	DS	DS	CL20	024
0033	0000 0000	ITEM	DS	CL16	024
0034	0000 0000	DATA	DS	CL48	024
0035	0000 0000	DS	DS	CL28	024
0036	0000 0000	NDR2	DS	CL28	024
0037	0000 0000	DC	DC	C' ACTIVITY	024
0038	0000 0000	DC	DC	C' START END	024
0039	0000 0000	DC	DC	C' WORK START	024
0040	0000 0000	DC	DC	C' TOTAL SLACK	024
0041	0000 0000	DS	DS	CL28	024
0042	0000 0000	NDR3	DS	CL45	024
0043	0000 0000	DC	DC	C'EVENT EVENT	024
0044	0000 0000	DC	DC	C' TIME TIME	024
0045	0000 0000	DC	DC	C' TIME TIME	024
0046	0000 0000	DS	DS	CL28	024
0047	0000 0000	NDR5	DS	CL24	024
0048	0000 0000	DC	DC	C'ANTELOPE VALLEY'	024
0049	0000 0000	DC	DC	C' CONSTRUCTION CO'	024
0050	0000 0000	DC	DC	C'COMPANY CRITICAL'	024
0051	0000 0000	DC	DC	C' PATH ACTIVITIES'	024
0052	0000 0000	DS	DS	CL30	024
0053	0000 0000	CAET	DC	C'000'	024
0054	0000 0000	PERC	DC	C'88'	024
0055	0000 0000				024
0056	0000 0000				024
0057	0000 0000				024
0058	0000 0000				024
0059	0000 0000				024
0060	0000 0000				024
0061	0000 0000				024
0062	0000 0000				024
0063	0000 0000				024
0064	0000 0000				024
0065	0000 0000				024
0066	0000 0000				024
0067	0000 0000				024
0068	0000 0000				024
0069	0000 0000				024
0070	0000 0000				024
0071	0000 0000				024
0072	0000 0000				024
0073	0000 0000				024
0074	0000 0000				024
0075	0000 0000				024
0076	0000 0000				024
0077	0000 0000				024
0078	0000 0000				024
0079	0000 0000				024
0080	0000 0000				024
0081	0000 0000				024
0082	0000 0000				024
0083	0000 0000				024
0084	0000 0000				024
0085	0000 0000				024
0086	0000 0000				024
0087	0000 0000				024
0088	0000 0000				024
0089	0000 0000				024
0090	0000 0000				024
0091	0000 0000				024
0092	0000 0000				024
0093	0000 0000				024
0094	0000 0000				024
0095	0000 0000				024
0096	0000 0000				024
0097	0000 0000				024
0098	0000 0000				024
0099	0000 0000				024
0100	0000 0000				024
0101	0000 0000				024
0102	0000 0000				024
0103	0000 0000				024
0104	0000 0000				024
0105	0000 0000				024
0106	0000 0000				024
0107	0000 0000				024
0108	0000 0000				024
0109	0000 0000				024
0110	0000 0000				024
0111	0000 0000				024
0112	0000 0000				024
0113	0000 0000				024
0114	0000 0000				024
0115	0000 0000				024
0116	0000 0000				024
0117	0000 0000				024
0118	0000 0000				024
0119	0000 0000				024
0120	0000 0000				024
0121	0000 0000				024
0122	0000 0000				024
0123	0000 0000				024
0124	0000 0000				024
0125	0000 0000				024
0126	0000 0000				024
0127	0000 0000				024
0128	0000 0000				024
0129	0000 0000				024
0130	0000 0000				024
0131	0000 0000				024
0132	0000 0000				024
0133	0000 0000				024
0134	0000 0000				024
0135	0000 0000				024
0136	0000 0000				024
0137	0000 0000				024
0138	0000 0000				024
0139	0000 0000				024
0140	0000 0000				024
0141	0000 0000				024
0142	0000 0000				024
0143	0000 0000				024
0144	0000 0000				024
0145	0000 0000				024
0146	0000 0000				024
0147	0000 0000				024
0148	0000 0000				024
0149	0000 0000				024
0150	0000 0000				024
0151	0000 0000				024
0152	0000 0000				024
0153	0000 0000				024
0154	0000 0000				024
0155	0000 0000				024
0156	0000 0000				024
0157	0000 0000				024
0158	0000 0000				024
0159	0000 0000				024
0160	0000 0000				024
0161	0000 0000				024
0162	0000 0000				024
0163	0000 0000				024
0164	0000 0000				024
0165	0000 0000				024
0166	0000 0000				024
0167	0000 0000				024
0168	0000 0000				024
0169	0000 0000				024
0170	0000 0000				024
0171	0000 0000				024
0172	0000 0000				024
0173	0000 0000				024
0174	0000 0000				024
0175	0000 0000				024
0176	0000 0000				024
0177	0000 0000				024
0178	0000 0000				024
0179	0000 0000				024
0180	0000 0000				024
0181	0000 0000				024
0182	0000 0000				024
0183	0000 0000				024
0184	0000 0000				024
0185	0000 0000				024
0186	0000 0000				024
0187	0000 0000				024
0188	0000 0000				024
0189	0000 0000				024
0190	0000 0000				024
0191	0000 0000				024
0192	0000 0000				024
0193	0000 0000				024
0194	0000 0000				024
0195	0000 0000				024
0196	0000 0000				024
0197	0000 0000				024
0198	0000 0000				024
0199	0000 0000				024
0200	0000 0000				024
0201	0000 0000				024
0202	0000 0000				024
0203	0000 0000				024
0204	0000 0000				024
0205	0000 0000				024
0206	0000 0000				024
0207	0000 0000				024
0208	0000 0000				024
0209	0000 0000				024
0210	0000 0000				024
0211	0000 0000				024
0212	0000 0000				024
0213	0000 0000				024
0214	0000 0000				024
0215	0000 0000				024
0216	0000 0000				024
0217	0000 0000				024
0218	0000 0000				024
0219	0000 0000				024
0220	0000 0000				024
0221	0000 0000				024
0222	0000 0000				024
0223	0000 0000				024
0224	0000 0000				024
0225	0000 0000				024
0226	0000 0000				024
0227	0000 0000				024
0228	0000 0000				024
0229	0000 0000				024
0230	0000 0000				024
0231	0000 0000				024
0232	0000 0000				024
0233	0000 0000				024
0234	0000 0000				024
0235	0000 0000				024
0236	0000 0000				024
0237	0000 0000				024
0238	0000 0000				024
0239	0000 0000				024
0240	0000 0000				024
0241	0000 0000				024
0242	0000 0000				024
0243	0000 0000				024
0244	0000 0000				024
0245	0000 0000				024
0246	0000 0000		</		

Construction Program--Analysis

0104		CUNS	START	300	001
0106	0100		NASA	10.0	002
0108			USING	MRF,10	002
0110		HERE	LN	13,NAST	007
0112	0000 0000		USING	HERE+0000,1:	002
0114	0277 050F 050D	CUNE	MVC	PR11201,PRT-1	002
0116	0277 050F 0507		MVC	PR11201,MOM1	002
0118	0000 0000		RAS	10,SPAC	002
0120	0000 0000		RAS	10,PNT1	002
0122	0277 050F 050D		MVC	PR11201,PRT-1	002
0124	0000 0000		MVC	PR11201,MOM2	002
0126	0277 050F 0507		RAS	10,PNT1	002
0128	0000 0000		MVC	PR11201,PRT-1	002
0130	0277 050F 050D		MVC	PR11201,MOM3	003
0132	0000 0000		RAS	10,PNT1	003
0134	0000 0000		RAS	10,SPAC	003
0136	0000 0000		TM	5MC,X'01'	003
0138	0101 0003		AC	1,OUT	003
0140	0710 02F0		SA	0,0	003
0142	1000		LN	0,THEG	003
0144	0000 0000	READ	MVC	CDIN(01),COIN-1	003
0146	0200 0511 0510		XIO	CDIN(X'22'),00	003
0148	0022 0511 0090		AC	0,ANSY	003
0150	0700 0000		BC	0,READ	003
0152	0700 0000		MPR	X'230',0	003
0154	0000 0230		AC	15,READ	003
0156	0700 0000	ABSY	TIOB	0,X'20'	004
0158	0000 0000		TIOB	RDER,X'21'	004
0160	0000 0000		TIOB	LSCD,X'20'	004
0162	0000 0000		BC	15,FILE	004
0164	0000 0000	RDER	MPR	X'000',0	004
0166	0000 0000		AC	15,READ	004
0168	0000 0000	FILE	MVC	010,0).DESC	004
0170	0200 0000 0511		PACK	STAR(2),STAR(3)	004
0172	0512 0501 0521		ZAP	10(2,0),STAR(2)	004
0174	0511 0010 0501		PACK	END(12),END(13)	004
0176	0512 0503 0520		ZAP	10(2,0),END(12)	004
0178	0511 0012 0503		PACK	LJRY(2),WORK(3)	005
0180	0512 0505 0527		ZAP	20(2,0),WORK(2)	005
0182	0511 0010 0505		ZAP	22(2,0),ZERO	005
0202	0511 0010 0506		ZAP	24(2,0),ZERO	005
0208	0511 0010 0506		ZAP	20(2,0),ZERO	005
020E	0511 0010 0506		AM	0,HTY	005
0210	0000 0500		BC	15,READ	005
0218	07F0 0000	LSCD	MVC	012,0).STAR	005
021C	0201 0000 0500		BAS	10,OUT	005
0222	0000 02F0				

Figure CONA-01

This part of the program does the initialization and reads all the cards. Data is placed into a table in core storage by means of indexing a table. The index register used here is register 8. When all cards are read, a star (*) is placed at the end of the table. The table is then printed out by means of the BAS to OUT after the last card is read.

Construction Program--Analysis

0224	9A01 E5E6	30AT	OI	SW,X'01'	005
022A	18AA		SR	A,A	005
022C	4A40 E996		LH	8,TBEG	005
0230	9A00 E5AC		NI	SW,X'00'	005
0234	F911 A010 802E	LOOP	CP	16(2,A),46(2,8)	006
023A	4700 E3FE		BC	13,COMP	006
023E	9A01 E5AC		OI	SW,X'01'	006
0242	021D E843 8006		MVC	8(130),0(8)	006
0248	021D A000 801E		MVC	0(30,8),30(8)	006
024E	021D A01E E843		MVC	30(30,8),BIN	006
0256	4A40 E568	COMP	AH	8,TMTY	006
025A	D501 801E E56A		CLC	30(2,8),STAR	006
025E	4770 E00E		BC	7,LOOP	006
0262	9101 E56C		TM	SW,X'01'	006
0266	4710 E006		BC	1,5ORT	006
026A	1888		SR	8,8	006

The table is sorted on the basis of the start event number.

026C	4A80 E996		LH	8,TBEG	007
0270	1899		SR	9,9	007
0272	4A90 E8CC		LH	9,TTWO	007
0276	F811 8018 8016		ZAP	24(2,8),20(2,8)	007
027C	F811 801A E8CA		ZAP	24(2,8),ZERO	007
02A2	D201 801C E56A		MVC	28(2,8),STAR	007
0288	D501 801C E56A	FIND	CLC	28(2,8),STAR	007
028E	4770 E14E		BC	7,STOR	007
0292	4A80 E56A		AH	8,TMTY	007
0296	D501 8000 E56A		CLC	0(2,8),STAR	007
029C	4780 E28C		BC	8,STEP	007
02A0	47F0 E132		BC	15,FIND	007
02A4	D20D E861 8010	STOR	MVC	BINB(14),16(8)	008
02AA	1888		SR	8,8	008
02AC	4A80 E996		LH	8,TBEG	008
02B0	F911 A012 E861	FINA	CP	18(2,8),BINB(2)	008
02B6	4770 E176		BC	7,FINB	008
02BA	D20D E86F 8010		MVC	BINC(14),16(8)	008
02C0	D20D 9000 E86F		MVC	0(14,9),BINC	008
02C6	4A90 E87E		AH	9,FTEN	008
02CA	4A80 E56A	FIND	AH	8,TMTY	008
02CE	D501 8000 E56A		CLC	0(2,8),STAR	008
02D4	4770 E15A		BC	7,FINA	008
02DA	A981 8000 F56A		MVC	8(2,9),STAR	008

Figure CONA-02

FIND--A compare is made to determine if it is the end of the table. If not, the starting time is moved to BINB. The end event of the table is compared to the start event in BINB. If they are equal, the start event is moved to BINC and then BINC is moved to TTWO.

The registers are incremented and a compare is made for the star. If no star has been reached, a branch is made back to FINA. When a star is found for the end of the table, indexed by register 8, a star is placed in TTWO which is indexed by register 9.

Construction Program--Analysis

030E	1899		SUM	SA	9.9	011
0310	4490 EACC			LM	9.TTWO	011
0314	F911 E996 900A			ZAP	TTIM(2),A(2,9)	011
0318	F911 E997 E99A			ZAP	BINB(12),TTIM	011
0320	F911 E998 E99C		TOTI	ZAP	BINE,TTIM	011
0324	F911 E999 900A			SP	BINE,R(2,9)	011
0328	F911 900A E886			ZAP	10(2,9),BIME	011
0332	4490 E87E			AM	9.FTEN	011
0336	D501 9000 E96A			CLC	0(2,9),STAR	011
0340	4770 E20A			BC	7.TOTI	011
0344	D200 E880 E880			MVC	BINB(12),STAR	911
0348	189A			SA	8.8	012
0352	4890 E996			LM	8.TBEG	012
0356	F911 A010 E861		HOPE	CP	10(2,8),BINB(2)	012
0360	4770 E25A			RC	7.GOOD	012
0364	F911 8012 E863			CP	10(2,8),BINB(2)	012
0368	4770 E25A			BC	7.GOOD	012
0372	D200 A010 E861			MVC	10(14,8),BINB	012
0376	4770 E25C			BC	15.INST	012
0380	4890 E568		GOON	AM	8.TMTY	012
0384	4770 E236			BC	15.HOPE	012
0388	1899		INST	SA	9.9	012
0392	4890 E8CC			LM	9.TTWO	012
0396	189A		GOGO	SA	8.8	012
0400	4890 E996			LM	8.TBEG	012
0404	D200 E861 9000		GOGI	MVC	BINB(14),0(9)	013
0408	F911 8010 E861			CP	10(2,8),BINB(2)	013
0412	4770 E2AA			BC	7.NOGO	013
0416	F911 8012 E863			CP	10(2,8),BINB(2)	013
0420	4770 E2AA			BC	7.NOGO	013
0424	D200 8C10 E861			MVC	10(14,8),BINB	013
0428	4490 E87E			AM	9.FTEN	013
0432	D501 9000 E96A			CLC	0(2,9),STAR	013
0436	4780 E29A			BC	8.FINO	013
0440	4770 E242			BC	15.GOGO	013
0444	189A		FINO	SA	8.8	013
0448	4890 E996			LM	8.TBEG	013
0452	1899			SA	9.9	014
0456	4890 E8CC			LM	9.TTWO	014
0460	4770 E132			BC	15.FINO	014
0464	4480 E568		NOGO	AM	8.TMTY	014
0468	D501 8000 E36A			CLC	0(2,8),STAR	014
0472	4780 E28C			BC	8.STEP	014
0476	4770 E268			BC	15.GOGI	014
0480	9601 E883		STEP	OI	SMC,X'01'	014
0484	4880 E568			SM	8.TMTY	014
0488	F811 E880 8016			ZAP	ACCA,20(2,8)	014
0492	F411 E880 8016			AP	ACCA,22(2,8)	014
0496	F811 8018 E880			ZAP	24(2,8),ACCA	014
0500	F811 801A E8C0			ZAP	26(2,8),ZERO	015
0504	9845 0001		SKIP	CIO	1(0),X'45'	015
0508	4780 E00A			BC	8.CONE	015
0512	4740 E20C			BC	4.SKIP	015
0516	9900 0878			MVA	X'878',0	015

Figure CONA-04

SUM--Subtract ending time of each activity from the ending time of the longest activity in TTWO. Enter the difference, which is the slack time, in the slack time position. The index register is then reset and the beginning of the table is again loaded into register 8.

HOPE--A compare is made with BINB (the beginning time and the end time). If they match, we return it to the main table.

INST--Transfer table TTWO back to the main table. If it is the end of the main table a zero is zapped into slack time here because

there is no slack time in the last event. Skip to the next page and print out the headings and the table. This time the switch at the end of section CONE send us to OUT for the next sort and print.

Construction Program--Analysis

0442	4710 E20C	AC	15,SKIP	015	
0444	1AAA	OUT	SR	A,R	015
0448	4AA0 E99A	LH	A,TREG	015	
044C	0200 EAT8 A000	OUT0	MVC	ITFM(16),0(8)	015
044E	0270 EAP8 EAAA	MVC	DATA,MASK	015	
044A	0208 EAAA A010	MVC	DAT(12),16(8)	015	
044E	0120 EAAA EAAA	ED	DATA,DAT0	015	
0446	0277 E9AE E9A0	MVC	PRT(120),PAT-1	016	
0444	0277 E9AE E9A0	MVC	PRT(120),MOR0	016	
0470	40A0 E9E8	AS	10,PRT	016	
0474	4AA0 E9A8	AH	A,THTY	016	
0478	0401 A000 E96A	CLC	0(2,A),STAA	016	
047E	4770 E2FA	AC	7,OUT0	016	
0482	4101 E9FA	TM	SW,X'01'	016	
048A	47A0 E000	BC	A, SORT	016	
0484	1AAA	SOA3	SR	A,R	016
048C	4AA0 E99A	LH	A,TREG	016	
0490	4A00 E9AC	NI	SW,X'00'	016	
0496	F911 A01A A030	LOPB	CP	2A(2,A),50(2,0)	016
049A	4700 E99E	BC	13,COM0	016	
049E	4A01 E9AC	OI	SW,X'01'	017	
04A2	0210 EAA3 A000	MVC	BIN(30),0(8)	017	
04A8	0210 A000 A01E	MVC	0(30,A),30(8)	017	
04AE	0210 A01E E843	MVC	30(30,0),BIN	017	
04A4	4A80 E9A8	COM0	AH	A,THTY	017
04B8	D501 A01E E36A	CLC	30(2,0),STAA	017	
04BE	4770 E33E	BC	7,LOPB	017	
04C2	4101 E9FA	TM	SW,X'01'	017	
04C6	4710 E330	BC	1,SOA3	017	

Figure CONA-05

The table is sorted on the basis of the lowest slack time first.

04CA	1AAA	SR	A,R	017	
04CC	4AA0 E99A	LH	A,TREG	017	
04D0	F911 A01A E8C6	COMA	CP	26(2,0),ZEAO	017
04D6	4770 E33C	BC	7,STOP	016	
04DA	4A80 E9A8	AH	A,THTY	018	
04DE	47F0 E37A	BC	15,COMA	018	
04E2	0201 A000 E96A	STOP	MVC	0(2,0),STAA	018

Figure CONA-05 (Part-2)

COMR--A star is placed at the point where the zero slack time ends.

Construction Program--Analysis

04EA	1AAA						
04EB	4AA0 E990						
04EC	4AA0 E9AC						
04ED	4A11 A010 A02E						
04EA	47A0 E3AC						
04FC	4A01 E5AC						
0500	D210 E843 8000						
0506	D210 A000 A01E						
050C	D210 A01E E843						
0512	4A80 E5AB						
0516	D501 A01E E90A						
051C	4770 E3AC						
0520	4101 E5AC						
0526	4710 E3A2						
052A	4843 0001						
052C	47A0 E3E6						
0530	4740 E3D2						
0536	4900 0323						
053B	47FD E3D2						
053C	D277 E56E E56D						
0542	D277 E56E E7C7						
054A	4DA0 E4EA						
054C	4DA0 E4D6						
0550	D277 E56E E56D						
0556	D277 E56E E6D7						
055C	4DA0 E4EA						
0560	D277 E56E E56D						
0566	D277 E56E E74F						
056C	4DA0 E4EA						
0570	4DA0 E4D6						
0576	1AAA						
057E	4AA0 E990						
057A	D20F E67B 8000						
0580	D22F E68A E8AA						
0586	D20A E8AA A010						
05AC	DE2F E68B E8AA						
0592	D277 E56E E56D						
0598	D277 E56E E65F						
059E	4DA0 E4EA						
SOR4	SR	A.A					018
	LM	A.TDFC					018
	MI	SW,X*00*					018
LOPC	CP	1612.A,4612.81					018
	BC	11.COMS					018
	DI	SW,X*01*					018
	MVC	81N(301,018)					018
	MVC	0130.A,3018)					018
	MVC	30130.A,81N					018
COMS	AM	8.TMTY					019
	CLC	3012.-1,STAR					019
	BC	7.LOPC					019
	TM	SW,X*01*					019
	BC	1.SUR4					019
SKIT	C10	1101.X*65*					019
	BC	8.PRA					019
	BC	4.SKIT					019
	NPR	X*323*.U					019
	BC	15.SKIT					019
PRA	MVC	PRT(120),PAT-1					019
	MVC	PRT(120),MOR5					019
	BAS	10.PRNT					020
	BAS	10.SPAC					020
	MVC	PRT(120),PAT-1					020
	MVC	PRT(120),MOR2					020
	BAS	10.PRNT					020
	MVC	PRT(120),PAT-1					020
	MVC	PRT(120),MOR3					020
	BAS	10.PRNT					020
	BAS	10.SPAC					020
	SR	A.A					020
	LM	A.TBEG					020
	MVC	ITEM(16),018)					020
	MVC	DATA.MASK					021
	ED	DATB(12),1618)					021
	ED	DATA.DATB					021
	MVC	PRT(120),PAT-1					021
	MVC	PRT(120),MOR4					021
	BAS	10.PRNT					021

Figure CONA-06

SOR4--The new short table is sorted on the basis of the highest start event number first. This table will contain all the events with no slack time. This means we start at the back of the critical path drawing.

SKIT--Skip to a new page, print out headers, and space. The address of the beginning of the main table is loaded into register 8. The first activity printed out here is actually the last activity of the drawing.

Construction Program--Analysis

0402	F011 E00A A010	LOPD	ZAP	R1NF,1A1Z,A1	021
0404	0201 A01C E001		MVC	2A1Z,A1,PERC	021
0406	0400 E00A	LOPE	AM	A,TH1Y	021
0408	0501 A010 E00A		CLC	01Z,A1,STAR	021
040A	0770 E002		RC	7,PRNM	022
040C	1000		SR	A,N	022
040E	0400 E00A		LM	A,TRFG	022
0402	0501 A01C E001	COMT	CLC	2A1Z,A1,PFAC	022
0408	0700 E000		RC	A,ADVA	022
040C	0501 A01C E00F		CLC	2A1Z,A1,CAET	022
0402	0700 E00C		RC	A,LIMP0	022
0406	0400 E00A	ADVA	AM	A,TH1Y	022
040A	0501 A000 E00A		CLC	01Z,A1,STAR	022
040E	0700 E00E		RC	A,MALT	022
0400	0770 E00C		RC	15,COMT	022
0404	F011 A01Z E000	PRNM	CP	1A1Z,A1,B1NE	022
0408	0770 E00A		RC	7,L00PE	022
0402	0200 E07A A000		MVC	ITEM1A,0101	023
0406	0220 E00A E00A		MVC	DATA,MASK	023
040A	0200 E00A A010		MVC	DATA121,16101	023
0404	0220 E00A E00A		FD	DATA,DATA	023
0408	0277 E00E E000		MVC	PRT(120),PRT-1	023
040C	0277 E00E E00F		MVC	PRT(120),MOR6	023
040E	0040 E0EA		BAS	10,PRNT	023
0400	0201 A01C E00F		MVC	2A1Z,A1,CAET	023
0404	0700 E00A		BC	15,LOPE	023
0408	0000 E000	MALT	MPR	X'444',0	024
040C	0700 E0CE		BC	15,MALT	024
040E	0000 0001	SPAC	C10	110,X'000'	024
0400	0700		BCA	8,10	024
0404	0700 E000		BC	4,SPAC	024
0408	0000 0111		MPR	X'111',0	024
040C	0700 E000		BC	15,SPAC	024
040E	0000 E00E 0070	PRNT	X10	PRT(X'40'),120	024
0400	0700 E0FE		BC	8,PRSY	024

Figure CONA-07

LOPD--Index the first activity with a percent sign and add 30 to the main table. Reset the index register and check for a percent sign. If so, jump to ADVA; if not, check for a cret. If it is a cret, jump back to LOPD. This process now replaces each cret with a percent sign. The percent tells us how far we have gone, and the cret tells us the next adjacent event along the path that has no slack time.

Construction Program--Analysis

0000	0700 E0E8	BC	4,FRNT	024	
0001	0700 0222	NPR	X'222',0	024	
0002	0700 E0E8	BC	15,FRNT	024	
0003	0700 E0E8	PBSV	TIOB	024	
0004	0700 E0E8	TIOB	X'40'	024	
0005	0700 E308	BCR	PRER,X'00'	024	
0006	07FA	BCR	15,10	025	
0007	0700 0223	PRER	NPR	025	
0008	0700 E0E8	DC	X'223',0	025	
0009	00	DC	15,FRNT	025	
0010		DC	C'	025	
0011		CDIN	OS	0CLA0	025
0012		DESC	OS	CL10	026
0013		STAA	OS	CL3	026
0014		ENOE	OS	CL3	026
0015		WORK	OS	CL3	026
0016		DC	OS	CL55	026
0017		STAB	OS	CL2	026
0018		ENDI	OS	CL2	026
0019		WORY	OS	CL2	026
0020	001E	THTY	DC	M'30'	026
0021	3C3C	STAR	DC	C'00'	026
0022		SM	OS	CL1	026
0023	00	DC	C'	027	
0024		PRT	OS	CL120	027
0025		SMA	OS	CL1	028
0026		MDR1	OS	CL27	028
0027		DC	OS	CL27	028
0028		DC	OS	C'ANTELOPE VALLEY'	028
0029		DC	OS	C'CONSTRUCTION CO'	028
0030		DC	OS	C'COMPANY ACTIVITIE'	028
0031		DC	OS	C'S SCHEDULE LOW'	028
0032		DC	OS	C'WISTER'	028
0033		DC	OS	CL27	029
0034		MDR4	OS	CL120	030
0035		DC	OS	CL28	030
0036		ITEM	OS	CL16	030
0037		DATA	OS	CL48	030
0038		DC	OS	CL28	030
0039		MDR2	OS	CL28	030
0040	4040 4040 C1C3 E3C9 E5C9 E3E8 4040 4040	DC	C' ACTIVITY	030	
0041	40E2 E3C1 D9E3 4040 4040 C505 C440 4040	DC	C' START END	030	
0042	4040 E606 D9D2 4040 4040 E2E3 C109 E340	DC	C' WORK START	030	
0043	4040 E306 E3C1 D340 4040 E2D3 C1C3 D240	DC	C' TOTAL SLACK	030	
0044		DC	CL28	031	
0045		MDR3	OS	CL45	032
0046	C5E5 C505 E340 4040 C5E5 C505 E340 40	DC	C'EVENT EVENT	032	
0047	4040 E3C9 D4C5 4040 4040 E3C9 D4C5 4040	DC	C' TIME TIME	032	
0048	4040 E3C9 D4C5 4040 4040 E3C9 D4C5 4040	DC	C' TIME TIME	032	
0049		DC	CL2A	032	
0050		MDR5	OS	CL29	033
0051	C105 E3C5 D3D6 D7C5 40E5 C103 D3C5 E8	DC	C'ANTELOPE VALLEY'	033	
0052	40C3 D4D5 E2E3 D9E4 C3E3 C9D6 D340 C3	DC	C' CONSTRUCTION C'	033	
0053	D4D4 D7C1 D5E4 40C3 D4C9 E3C9 E3C1 D3	DC	C'COMPANY CRITICAL'	033	
0054	40D7 C1E3 C040 C1C3 E3C9 E3C9 E3C9 C5E2	DC	C' PATH ACTIVITIES'	033	
0055		DC	CL30	034	
0056	4C4C	CAET	DC	C'00'	035
0057	4C4C	PEAC	DC	C'00'	035
0058					
0059		RIN	OS	CL30	035
0060		RINH	OS	CL14	036
0061		RINC	OS	CL14	036
0062		FTIM	DC	M'14'	036
0063		ACCA	DC	X'000C'	036
0064		SMR	OS	CL1	036
0065		SMC	DC	X'00'	037
0066		TTIM	DC	X'000C'	037
0067		RIME	DC	X'000C'	037
0068		RASE	DC	YIMRE+40961	037
0069		MASK	DC	X'4070702020222222'	037
0070	4070 7070 7077 7777	DC	X'7777707070222222'	037	
0071	7777 7070 7077 7777	DC	X'7777707070222222'	037	
0072	7777 7070 7077 7777	DC	X'7777707070222222'	037	
0073	7777 7070 7077 7777	DC	X'7777707070222222'	037	
0074	7777 7070 7077 7777	DC	X'7777707070222222'	037	
0075	7777 7070 7077 7777	DC	X'7777707070222222'	037	
0076	7777 7070 7077 7777	DC	X'7777707070222222'	037	
0077	7777 7070 7077 7777	DC	X'7777707070222222'	037	
0078	7777 7070 7077 7777	DC	X'7777707070222222'	037	
0079	7777 7070 7077 7777	DC	X'7777707070222222'	037	
0080	7777 7070 7077 7777	DC	X'7777707070222222'	037	
0081		DATA	OS	CL12	037
0082		ZEMO	DC	X'000C'	039
0083		MULL	DC	C'0000'	039
0084		TTMO	DC	Y(TBLE)	039
0085		TBLE	OS	CL200	039
0086		TBEG	DC	Y(TABL)	041
0087		TABL	OS	CL240	041
0088		END	OS	CONS	043
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0090					
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0200					

Figure CONA-08

Construction Program--Output

Output Number 01

ANTELOPE VALLEY CONSTRUCTION COMPANY ACTIVITIES SCHEDULE

ACTIVITY	START EVENT	END EVENT	WORK TIME	START TIME	TOTAL TIME	SLACK TIME
CLEAN UP	23	30	6			
SIGN CONTRACT	24	30	3			
SEE LAWYER	24	31	2			
SEE	29	31	3			
MEET	29	32	1			
PRAY	30	32	3			
COLLECT	31	33	2			
FINISH	32	33	3			
END	33	34	1			
SECURE TIME	13	16	2			
EXCAVATE	13	17	5			
BRIDGE	12	16	4			
EXHUME	12	17	3			
RELATE	14	18	7			
COMPUTE	19	20	2			
COMPOSE	16	20	4			
COMPOSITE	15	18	1			
RUNWAY	15	19	3			
REDICULOUS	15	21	7			
JANUARY	17	20	2			
WEDNESDAY	20	26	3			
FEBRUARY	17	21	3			
MARCH	14	25	3			
APRIL	14	22	4			
MAY	14	23	5			
JUNE	19	25	4			
JULY	20	25	3			
WEDNESDAY	22	28	3			
AUGUST	19	23	5			
DECEMBER	21	26	4			
SEPTEMBER	21	23	1			
OCTOBER	20	24	3			
NOVEMBER	21	24	5			
DENSITY	2	5	4			
VERILY	13	14	3			
SUNDAY	24	26	1			
HALFWAY	1	3	3			
CONFUSION	1	4	2			
JUNK	4	9	3			
LOUSY	4	10	5			
EFFORT	2	6	2			
APPLESAUCE	1	2	1			
KLUNK	5	10	6			
NOTYET	7	11	5			
TUESDAY	25	28	4			
START		1				
UNDONE	14	15	1			
TANTALIZE	12	14	3			
OMITTED	10	11	2			
ROTTEN	10	13	3			
SILLY	11	14	4			
OPTIMUM	8	11	4			
POORLY	8	12	3			
QUIET	9	12	5			
MAYBE	6	11	7			
FINALLY	2	7	5			
GOOFY	3	7	1			
HARDLY	4	7	4			
IMAGINE	4	8	2			
SATURDAY	22	30	4			
NEARLY DONE	24	30	2			
THURSDAY	25	29	5			
FRIDAY	23	29	2			

Figure CON-07

Construction Program--Output

Output Number 02

ANTELOPE VALLEY CONSTRUCTION COMPANY ACTIVITIES SCHEDULE

ACTIVITY	START EVENT	END EVENT	WORK TIME	START TIME	TOTAL TIME	SLACK TIME
START		1			3	
RAILROAD	1	3	3		3	
CONFUSION	1	4	2		2	
APPLESAUSAGE	1	2	1		1	
DENSITY	2	5	4	1	5	
EFFORT	2	6	2	1	3	
FINALLY	2	7	5	1	6	
GOOFY	3	7	1	3	4	2
JUNK	4	9	3	2	5	
HARDLY	4	7	4	2	6	
IMAGINE	4	8	2	2	4	
KLUNK	5	10	6	5	11	
LOUSY	6	10	5	3	8	3
MAYBE	6	11	7	3	10	3
NOTYET	7	11	5	6	11	2
OPTIMUM	8	11	4	4	8	3
POORLY	8	12	3	4	7	
QUIET	9	12	5	5	10	
OMITTED	10	11	2	11	13	
ROTTEN	10	13	3	11	14	
SILLY	11	14	4	13	17	
DREAGE	12	16	6	10	16	2
FRUSTRATE	12	17	3	10	13	6
TANTALIZE	12	14	3	10	13	6
SECOND TIME	13	16	2	14	16	
EXCAVATE	13	17	5	14	19	
VERTLY	13	14	3	14	17	
INCORPSE	14	15	1	17	18	
COMPOSITE	15	18	1	18	19	6
RUNWAY	15	19	3	18	21	
REFRIGILIOUS	15	21	7	18	25	
PLATE	16	18	7	16	23	
COMPISE	16	20	4	16	20	3
JANUARY	17	20	2	19	21	2
FEBRUARY	17	21	3	19	22	3
MARCH	18	25	3	23	26	
APRIL	18	22	6	23	27	
MAY	18	23	5	23	28	
COMPUTE	19	20	2	21	23	
JUNE	19	25	4	21	25	1
AUGUST	19	23	5	21	26	2
MONDAY	20	26	3	23	26	5
JULY	20	25	3	23	26	4
OCTOBER	20	24	3	23	26	2
DECEMBER	21	24	4	25	29	2
SEPTEMBER	21	23	1	25	26	2
NOVEMBER	21	24	5	25	30	
WEDNESDAY	22	28	3	27	30	
SATURDAY	22	30	4	27	31	3
CLEAN UP	23	30	6	28	34	
FRIDAY	23	29	2	28	30	1
SUNDAY	24	26	1	30	31	
NEARLY DONE	24	30	2	30	32	2
TUESDAY	25	28	4	26	30	
THURSDAY	25	29	5	26	31	
SIGN CONTRACT	26	30	3	31	34	
SEE LAWYER	28	31	2	30	32	2
SHE	29	31	3	31	34	
HOPE	29	32	1	31	32	5
PRAY	30	32	3	34	37	
COLLECT	31	33	2	34	36	6
FINISH	32	33	3	37	40	
END	33	34	1	40	41	

Figure CON-08

Construction Program--Output

Output Number 03

ANTELOPE VALLEY CONSTRUCTION COMPANY CRITICAL PATH ACTIVITIES

ACTIVITY	START EVENT	END EVENT	WEEK TIME	START TIME	TOTAL TIME	SLACK TIME
END	33	34	1	40	41	
FINISH	32	33	3	37	40	
PRAY	30	32	3	34	37	
SIGN CONTRACT	26	30	3	31	34	
CLEAN UP	23	30	6	28	34	
SUNDAY	24	26	1	30	31	
NOVEMBER	21	24	5	25	30	
MAY	18	23	5	23	28	
REDCLOUS	15	21	7	18	25	
RELATE	16	18	7	16	23	
SECOND TIME	13	14	2	14	16	
UNDINE	14	15	1	17	18	
VERILY	13	14	3	14	17	
SILLY	11	14	4	13	17	
ROTTEN	10	13	3	11	14	
ROTTER	10	13	3	11	14	
OMITTED	10	11	2	11	13	
KLUNK	5	10	6	5	11	
KLUNK	5	10	6	5	11	
DENSITY	2	5	4	2	5	
APPLESAUCE	1	2	1		1	
START		1				

Figure CON-09

Antelope Valley Construction Co. Program

AVCC	10	0154	00	AVCC	10	0154	00	CDIN	10	0057	07	LIND	10	0065	07
AVCC	10	0154	11	DAIF	10	0054	11	DAYS	10	0700	11	DIRA	10	0074	07
AVCC	10	0154	07	DIRA	10	0054	07	DIRN	10	0700	08	FMS	10	0102	05
AVCC	10	0154	07	DIRA	10	0054	07	FIND	10	0700	0C	ETID	10	0050	04
AVCC	10	0154	05	DIRA	10	0164	05	FAPA	10	0047	07	FKID	10	0053	04
AVCC	10	0154	06	DIRA	10	0054	06	FLAK	10	0009	00	FLAP	10	0710	00
AVCC	10	0154	06	DIRA	10	0054	06	HUNG	10	0041	77	HURD	10	0702	01
AVCC	10	0154	07	DIRA	10	0054	07	HUNG	10	0000	0E	HURH	10	0014	01
AVCC	10	0154	07	DIRA	10	0054	07	HUNZ	10	075A	77	HINT	10	0100	03
AVCC	10	0154	07	DIRA	10	0054	07	INCH	10	07A3	0E	INCC	10	0713	04
AVCC	10	0154	07	DIRA	10	0054	07	ITUM	10	0702	0C	ITUD	10	004F	04
AVCC	10	0154	03	DIRA	10	01AC	03	MSKA	10	001E	0C	MSKH	10	0070	0E
AVCC	10	0154	03	DIRA	10	057A	03	PHFK	10	0504	03	PRNT	10	04E2	05
AVCC	10	0154	05	DIRA	10	054C	05	PUNA	10	00C3	01	PUNB	10	04C5	11
AVCC	10	0154	07	DIRA	10	004F	07	PUNF	10	00E7	07	PUNY	10	00EF	07
AVCC	10	0154	07	DIRA	10	004F	07	PUNY	10	053E	03	SLUP	10	034C	05
AVCC	10	0154	05	DIRA	10	017A	05	SKIP	10	0554	03	SW	10	00A3	00
AVCC	10	0154	03	DIRA	10	054C	03	STEP	10	034E	05	WEAK	10	070A	01
AVCC	10	0154	11	DIRA	10	0771	11	TUTA	10	01E0	05				

0154	00F0	AVCC	START	340	001
0154	00F0	AVCC	RASR	15.0	002
0154	00F0	AVCC	USING	0.15	002
0154	0040 F3FE	AVCC	RAS	10,SKIP	002
0154	0277 F423 F422	AVCC	MVC	PRT(120),PRT=1	002
0160	0277 F423 F400	AVCC	MVC	PRT,MORA	002
0166	0040 F3AC	AVCC	RAS	10,PRNT	002
0166	0277 F423 F513	AVCC	MVC	PRT,MORA	002
0166	0040 F3AC	AVCC	RAS	10,PRNT	002
0170	0040 F3AC	AVCC	RAS	10,SPAC	002
0174	0040 F3AC	AVCC	MVC	CDIN(01),CDIN=1	002
0174	0247 F6FC F4FB	AVCC	XIO	CDIN(X'22'),00	002
017E	0027 F4FC 0030	AVCC	AC	R,RHSY	002
0184	4740 F03E	AVCC	HC	4,READ	003
0184	4740 F022	AVCC	MUR	X'1A1',0	003
0184	9900 01A1	AVCC	HC	15,READ	003
0190	4740 F072	AVCC	TIUB	0,X'20'	003
0194	9A70 F03E	AVCC	TIUB	RISER,X'21'	003
0194	9A21 F04E	AVCC	TIUB	LSCD,X'24'	003
019C	9A24 F054	AVCC	HC	15,MINT	003
01A0	4740 F054	AVCC	MUR	X'999',0	003
01A4	9900 0999	AVCC	BC	15,HEAD	003
01A4	4740 F022	AVCC	UI	SM,X'01'	003
01A4	9A01 F74D	AVCC	TH	SM,X'01'	003
01A4	9101 F74D	AVCC	HC	1,CHPK	003
01A4	4710 F04A	AVCC	CLI	FLAG,C'T'	003
01A4	95E3 F747	AVCC	HC	M,THIA	003
01A4	4740 F047	AVCC	CLI	FLAG,C'T'	004
01C0	95C4 F747	AVCC	HC	M,INC	004
01C4	4740 F100	AVCC	CLI	FLAG,C'T'	004
01C4	95C3 F742	AVCC	HC	M,EXP	004
01C4	4740 F20E	AVCC	TH	SM,X'01'	004
01D0	9101 F74D	AVCC	HC	1,ENS	004
01D4	4710 F24C	AVCC	BC	15,HEAD	004
01D4	4740 F022	AVCC	CLI	FLAG,C'T'	004
01D4	95E3 F7C7	AVCC	HC	R,INNE	004
01E0	4740 F374	AVCC	BC	15,ENS	004
01E4	4740 F24C	AVCC	MVC	FLAP,FLAG	004
01FA	0700 F7C7 F742	AVCC	MVC	WEAK(2),WEFK	004
01FE	0701 F5A4 F4FC	AVCC	IVC	DAYS(14),DATE	004
01FA	0211 F5M9 F4FE	AVCC	MVC	PRT,MUNC	005
01FA	0277 F423 F5A8	AVCC	RAS	10,PRNT	005
020C	0040 F3AC	AVCC	MVC	MURZ(120),MURZ=1	005
0204	0277 F404 F403	AVCC	MVC	TITL(11),MORD	005
0204	0211 F418 F47C	AVCC	PACK	INCC,INCA	005
0210	F247 F74D F710	AVCC	MVC	INCR(13),MSKA	005
0214	0700 F47D F7CA	AVCC	EU	INCR(13),INCC	005
021C	0100 F47D F74D	AVCC	MVC	PRT,MURZ	005
0222	0277 F473 F404	AVCC	RAS	10,PRNT	005
0224	0040 F3AC	AVCC	MVC	MURZ(120),MURZ=1	005
022C	0277 F404 F403	AVCC	MVC	TITL(11),MORD	006
0232	0211 F418 F47C	AVCC	PACK	INCC,EXPA	006
0234	F247 F74D F71A	AVCC	MVC	INCR(13),MSKA	006
023E	020C F47D F7CA	AVCC	EU	INCR(13),INCC	006
0244	0100 F47D F74D	AVCC	MVC	PRT,MURZ	006
024A	0277 F423 F404	AVCC			

Figure AVC-01

AVC Program

039A	D277 F423 F588	MVC	PRT,MORC	012
039C	F247 F75D F718	PACK	EXTU,EXPA	012
0392	FA44 F762 F75D	AP	ETUT,EXTO	012
039A	FA44 F767 F75D	AP	ETOD,EXTO	012
039E	47F0 F022	BC	IS,READ	012
03A2	D277 F423 F588	MVC	PRT,MORC	012
03A8	D24F F76D F76C	MVC	COOD(80),COOD-1	012
03AE	D201 F76D F584	MVC	PUNA(21),WEAK	012
0384	D211 F76F F589	MVC	PUNH(1A),DAYS	012
038A	4DA0 F38C	BAS	10,PRNT	012
038E	D277 F604 F603	MVC	MORZ(120),MORZ-1	013
03C4	F374 F7A1 F753	UNPK	PUNG,ITOT	013
03CA	D20C F62D F7C8	MVC	INCB(13),MSKA	013
03D0	DE0C F62D F753	ED	INCB(13),ITOT	013
03D6	D211 F618 F67C	MVC	TITL(1A),MORU	013
030C	D277 F423 F604	MVC	PRT,MORZ	013
03E2	4DA0 F38C	BAS	10,PRNT	013
03E6	D277 F604 F603	MVC	MORZ(120),MORZ-1	013
03EC	F374 F789 F762	UNPK	PUNG,ETOT	013
03F2	D20C F62D F7C8	MVC	INCB(13),MSKA	013
03FA	DE0C F62D F762	ED	INCB(13),ETOT	014
03FE	D211 F618 F6AE	MVC	TITL(1A),MORU	014
0404	D277 F423 F604	MVC	PRT,MORZ	014
040A	4DA0 F38C	BAS	10,PRNT	014
040E	D277 F604 F603	MVC	MORZ(120),MORZ-1	014
0414	FA44 F7C2 F753	ZAP	ACCA,ITOT	014
041A	FA44 F7C2 F762	SP	ACCA,ETOT	014
0420	D20E F62D F7D5	MVC	INCB,MSKH	014
0426	DE0E F62D F7C2	ED	INCB,ACCA	014
042C	D211 F618 F6A0	MVC	TITL(1B),MORF	015
0432	D277 F423 F604	MVC	PRT,MORZ	015
0438	4DA0 F38C	BAS	10,PRNT	015
043C	D277 F604 F603	MVC	MORZ(120),MORZ-1	015
0442	F374 F799 F75A	UNPK	PUNF,ITOD	015
0448	D20C F63C F7C8	MVC	ITDB,MSKA	015
044E	DE0C F63C F758	ED	ITDB,ITOD	015
0454	D211 F618 F6B2	MVC	TITL(1A),MORC	015
045A	D277 F423 F604	MVC	PRT,MORZ	015
0460	4DA0 F38C	BAS	10,PRNT	015
0464	D277 F604 F603	MVC	MORZ(120),MORZ-1	015
046A	F374 F7A1 F767	UNPK	PUNG,ETOD	016
0470	D20C F649 F7C8	MVC	ETDB,MSKA	016
0476	DE0C F649 F767	ED	ETDB,ETOD	016
047C	D211 F618 F6C4	MVC	TITL(1B),MORH	016
0482	D277 F423 F604	MVC	PRT,MORZ	016
0488	4DA0 F38C	BAS	10,PRNT	016
048C	D277 F604 F603	MVC	MORZ(120),MORZ-1	016
0492	F844 F7C2 F758	ZAP	ACCA,ITOD	016
0498	F844 F7C2 F767	SP	ACCA,ETOD	016
049E	D20E F656 F705	MVC	OTDB,MSKB	017
04A4	DE0E F656 F7C2	ED	OTDB,ACCA	017
04AA	D200 F783 F74C	MVC	FLAK,T	017
04B0	D211 F618 F6D6	MVC	TITL(1B),MORJ	017
04B6	D277 F423 F604	MVC	PRT,MORZ	017
04BC	4DA0 F38C	BAS	10,PRNT	017
		SLOP		
		ENS		

Figure AVC-03

AVCC--Output

ANTELOPE VALLEY CONSTRUCTION COMPANY
CASH REQUIREMENTS PREDICTION

WEEK NUMBER 01	03MAY70 TO 09MAY70		
INCOME	3,000.00		
EXPENSE	5,000.00		
DIFFERENCE	2,000.00 -		
INCOME TO DATE		3,000.00	
EXPENSE TO DATE			5,000.00
DIFFERENCE TO DATE			2,000.00 -
WEEK NUMBER 02	10MAY70 TO 16MAY70		
INCOME	8,500.00		
EXPENSE	15,000.00		
DIFFERENCE	6,500.00 -		
INCOME TO DATE		9,500.00	
EXPENSE TO DATE			20,000.00
DIFFERENCE TO DATE			10,500.00 -
WEEK NUMBER 03	17MAY70 TO 23MAY70		
INCOME	12,000.00		
EXPENSE	3,500.00		
DIFFERENCE	8,500.00		
INCOME TO DATE		21,500.00	
EXPENSE TO DATE			23,500.00
DIFFERENCE TO DATE			2,000.00 -
WEEK NUMBER 04	24MAY70 TO 30MAY70		
INCOME	8,945.67		
EXPENSE	6,750.00		
DIFFERENCE	2,195.67		
INCOME TO DATE		30,445.67	
EXPENSE TO DATE			30,250.00
DIFFERENCE TO DATE			195.67
END AVCC	MUSTEA		

Figure AVC-06

SECTION II

DATA STRUCTURES

II DATA STRUCTURES

The Swanson Study was presented to help the reader become aware of a need for other programming tools other than the higher level languages. If he is going to be able to function at the level of this study, it will be well to master certain aspects of Data Structures.

The object of this section of the book is to provide the material so that the reader can develop his own information system, or at least work at the level of proficiency the Swanson Study demands.

This introductory section presents a definition of data structures by depicting the recommendations of the Committee of the Association of Computing Machinery for a college level course in data structures. This allows the reader to choose other subjects he might desire and pursue them by using the bibliography in the appendix.

HISTORY

Prior to the 1950's programmers were concerned with both memory allocation and the algorithm. Each programmer accepted the fact that he must exercise adequate control over fast-access storage, if the program were too large for the storage and he had to use auxiliary.

During the 1950's higher level languages were introduced, and the programmer became more involved with problem solving than with the machine and its various aspects. The programs grew in size and so did the problem of storage control and storage allocation.

Due to the shielding effect of the higher level languages, the programmer moved further and further from the reality of the efficiency

and inefficiency of the machine, and the requirement was for larger and more expensive memories. When the cost became so great, the realization came that it was time for the programmer to face reality and go back and work with such elements as information structures and storage allocation.

It would appear that we have come back to a point that has characteristics similar to the pre-1950's.

It will now be necessary for the programmer to put forth the extra effort to master the areas he avoided even though they may appear to be more difficult in some instances.

The term "INFORMATION STRUCTURE" as used in this book relates to the contents of the course of study presented by the Association of Computing Machinery in 1968. This term is used to infer the structural relationships between various data elements.

The Curriculum Committee on Computer Science of the Association for Computing Machinery¹ has offered recommendations for a college level course, entitled "Data Structures".

1. Basic Concepts of Data

- a. Representation of information as data inside and outside the computer, bits, nodes and data elements
- b. Data files and tables
- c. Names, values, environments
- d. Use of pointer or linkage variables to represent data structures

¹"Curriculum 68, Recommendations for Academic Programs in Computer Science," Communications of the ACM, Vol. II, No. 3, March 1968, pp. 151-197.

- e. Identifying entities about which data is to be maintained
 - f. Selecting data nodes and structures which are to be used in problem solution
 - g. Storage media, storage structures, encoding data and transformations from one medium and/or code to another
 - h. Alternative representations of information and data
 - i. Packing, unpacking and compressing data
 - j. Data formats, data description languages, specifications of data transformations
2. Linear Lists and Strings
- a. Stacks, last-in first-out, first-in first-out, double ended, and other linear lists
 - b. Sequential versus linked storage
 - c. Single versus double link circular lists
 - d. Character strings of variable length
 - e. Word packing, part word addressing
 - f. Pointer manipulation
 - g. Insertion, deletion and accessing of list elements
3. Arrays and Orthogonal Lists
- a. Storage of rectangular arrays in a one dimensioned media
 - b. Storage mapping functions
 - c. Direct and indirect address computation
 - d. Space requirements
 - e. Set up time
 - f. Accessing time and dynamic relocation times
 - g. Storage and accessing triangular arrays
 - h. Tetrahedral arrays

- i. Space matrices
4. Tree Structures
- a. Trees, subtrees, ordered trees, free trees, oriented trees, binary trees
 - b. Representation of trees using binary trees
 - c. Sequential techniques
 - d. Threaded lists
 - e. Insertion, deletion and accessing of elements of trees
 - f. Relative referencing, finding successors and predecessors
 - g. Walking through trees
 - h. Examples of tree structures such as algebraic formulas, arrays and other heirarchic data structures (PLI/and COBOL)
5. Storage Systems and Structures
- a. Behavioral properties of Unit record (card)
 - b. Random access (core)
 - c. Linear (tape)
 - d. Intermediate (disk and drum)
 - e. Storage media including cost, size, speed, reusability
 - f. Inherent record and file structure
 - g. Deficiencies and interrelation of these properties
 - h. Influence and machine structure--addressing on data structuring
 - i. Hierarchies of storage, virtual memory, segmentation, paging and bucketing
 - j. Influence of data structures and data manipulation on

storage systems

k. Associative structures, both hardware and software

6. Storage Allocation and Collection

7. Multi-linked Structures

8. Sorting (ordering techniques)

a. Radix sorting

b. Radix exchange sorting

c. Merge sorting

d. Bubble sorting

e. Address table sorting

f. Topological sorting

g. Comparative efficiency of sorting techniques

h. Effect of data structures and storage structures on
sorting techniques

9. Symbol Table and Searching

10. Data Structures in Programming Languages

a. Compile time and run time data structures needed to
implement source language data structures of programming
languages

b. Linkage between partially executed procedures

c. Data structures of co-routines, scheduled procedures,
and control structures

d. Storage management of data structures in procedure
oriented languages

e. Examples of higher level languages which include list
processing and other data structure features

11. Formal Specification of Data Structures

- a. Specification of syntax of classes of data structures
 - b. Predicate selectors and constructors for data manipulation
 - c. Data definition facilities
 - d. Programs as data structures
 - e. Computers as data structures and transformations
 - f. Formal specification of semantics
 - g. Formal systems viewed as data structures
12. Generalized Data Management Systems
- a. Structures of generalized data management systems
 - b. Directory maintenance
 - c. User languages (query)
 - d. Data description maintenance
 - e. Job management
 - f. Embedding data structures in generalized data management systems
 - g. Examples of generalized data management systems and comparison of system features

II DATA STRUCTURES

If the reader will notice the numbering sequence used for this section, it does contain a certain structuring. The section begins with Linear Lists and ends with Virtual Memory.

There was a definite reason in presenting the subjects in this particular order, and that was one of logical development of the tools needed to create information systems.

The basic information starts with lists (and links), which will partially answer the problem of linking, created by the Swanson Study. If the reader will examine the definition of data structures, he will see that Files are listed at lb. Files are listed in this section at 2.4 because they fit logically into the development of the material for an information system.

The final culmination of this section ends with Dynamic Storage Allocation and Virtual Memory. They are presented as a little icing for the cake, since certain programming positions do not require an awareness of them, however they are presented for the reader who wants to do just a little bit more.

II DATA STRUCTURES

- 2.1 Introduction
- 2.2 Linear Lists
 - 2.2.1 Stacks, Deques and Queues
 - 2.2.2 Linked Lists
- 2.3 Trees
 - 2.3.1 Hierarchical Ring Structure
 - 2.3.2 Traversing Trees - Preorder and Post Order
 - 2.3.3 Directories
 - 2.3.3.1 Structured Tree Directory
 - 2.3.3.2 Random Ordered Directories
- 2.4 Files
 - 2.4.1 Sequential Files
 - 2.4.2 Inverted Files
 - 2.4.2.1 Creation, Maintenance and Deletion of Keys

- 2.4.2.2 Logical Operators
 - 2.4.2.2.1 Logical Operators
 - 2.4.2.2.2 Logical Operator "OR"
 - 2.4.2.2.3 Logical Operator "AND"
- 2.4.3 Multilist File
 - 2.4.3.1 Deletion of Records and Addition of Keys
- 2.4.4 Cellular Multilist File
- 2.5 Sorting
 - 2.5.1 Sorting Methods
 - 2.5.1.1 Key Sort
 - 2.5.1.2 Detached Key Sort
 - 2.5.1.3 Non Detached Key Sort
 - 2.5.2 Radix Sort
 - 2.5.3 Quicksort
 - 2.5.4 Topological Sort
 - 2.5.5 Quadratic Sort
 - 2.5.6 Calculated Address Sort
 - 2.5.7 Sorting by Insertion
 - 2.5.8 Interchange Sort
- 2.6 Merging
 - 2.6.1 Two Way Merge
 - 2.6.1.1 Merging Ordered Files with Subfiles
 - 2.6.1.1.1 Merging Step Downs
- 2.7 Search
 - 2.7.1 Linear Search
 - 2.7.2 Binary Search
- 2.8 Dynamic Storage Allocation

2.9 Virtual Memory

List Formats

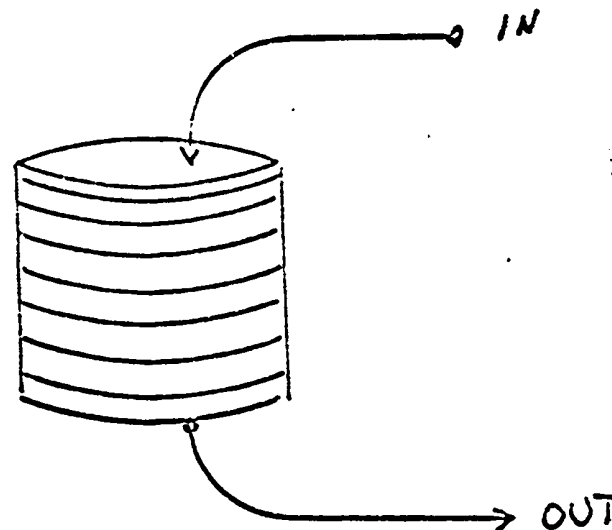
An important aspect of the list languages is that the original list of available space must be either set up automatically by the system or manually by the programmer. This original space amounts to a pool or block of storage which the program can draw from as needed.

Three kinds of fixed format lists have been adopted by various writers. The format presented here was originally named by Knuth:

1. QUEUES enable opposite end accessibility; input one end and output the other end.
2. DEQUES enable double end accessibility; a word or record can be added or deleted at either end.
3. STACKS enable double end accessibility; a word or record can be added to one end.

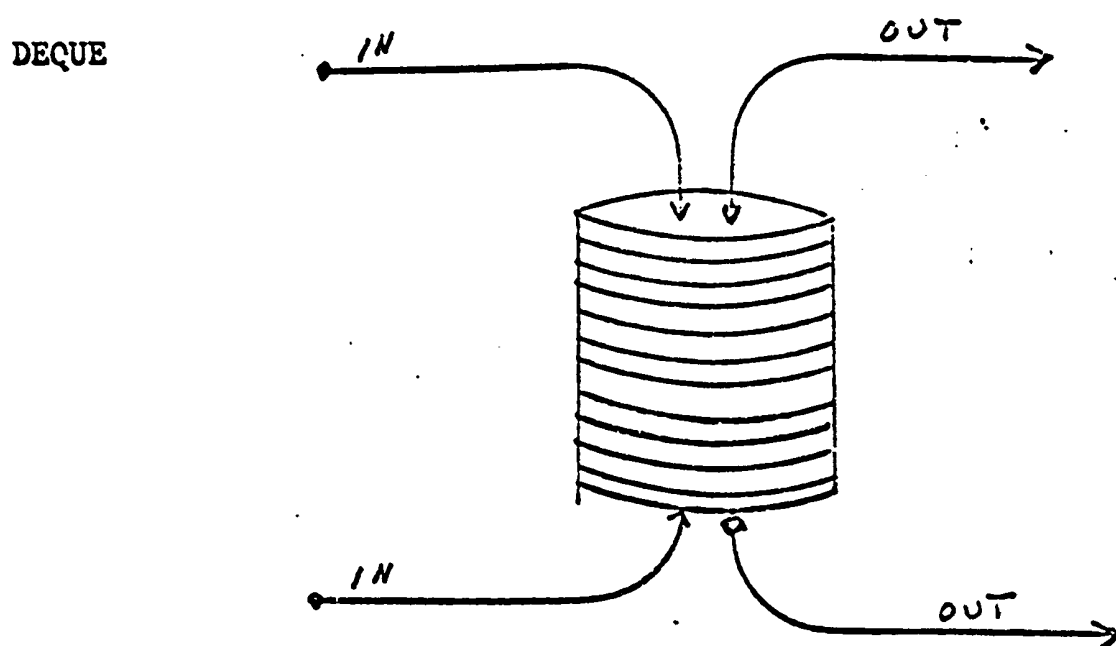
It is important to note that with the formats presented here, it is not possible to access records that are located in the center of either method.

QUEUE

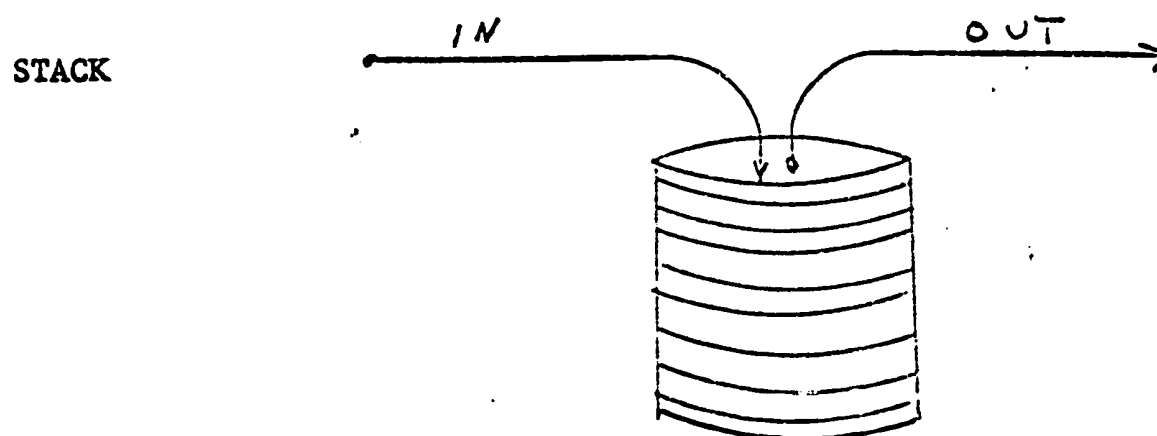


The Que has two ends, and the records are entered at the top

and withdrawn from the bottom. This method is generally used when records are to be serviced in the order of receipt, (FIFO, which is first-in first-out). This method also is used in accounting. A Queue is used for a list of areas to receive records from sequential files, and it is also used as a list of areas to be written on a sequential file or files.



It is possible to insert or delete a record from either the top or the bottom of the deque. This could be called last-in first-out or first-in first-out.

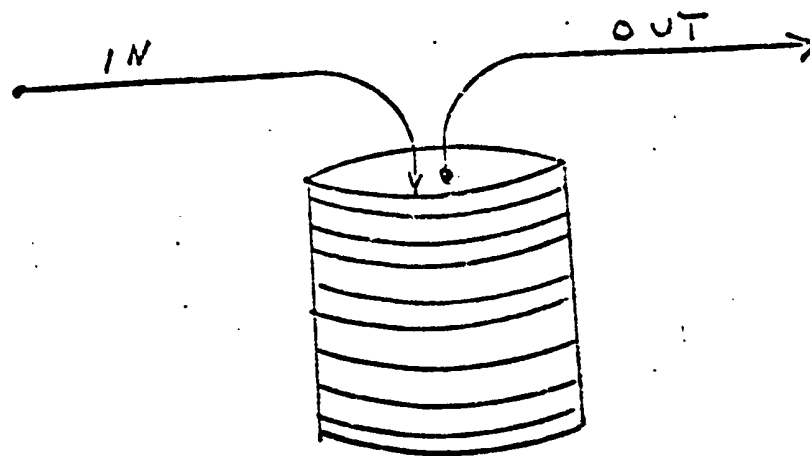


A stack takes out the most recent item that was put in. It uses the last-in first-out method. This method is one that is used more for list than the other methods. The list of available space is

one application for this method.

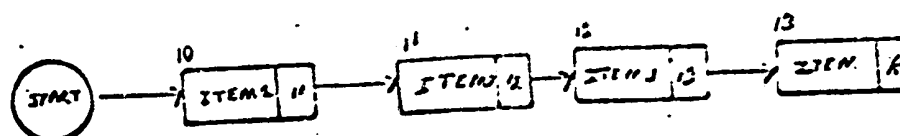
List of Available Space

A linked list indicates some type of instrument to control the space not in use. As evidenced by the language SLIP, written by Wizenbaum, a List of Available Space is quite necessary. It is essentially a stack.



The last (top) record is the first one out. All nodes that are not being used are linked by this list.

Linked Lists



ADDRESS	ITEM	LINK
10	ITEM 1	11
11	ITEM 2	12
12	ITEM 3	13
13	ITEM 4	K

LINK VARIABLE →

Go Through a Stack

The linked list is much more flexible than the sequential list. The addresses 10, 11, 12 and 13 are addresses in memory, and the figure (*) is a null link, which stops the action.

The program requires a link variable that points to the address number 10, and from address 10, all other items on the list are easily found.

```

*   G O T H R O U G H A S T A C K
*   REGISTER 9 POINTS TO TOP OF STACK
*   REGISTER 10 HAS RETURN ADDRESS FROM SUBR.
*   REGISTER 11 HAS RETURN ADDRESS FROM
*   SUBROUTINE VIST
*
A   START 0
    USING *,0
STAR HAS 11,VIST   GO TO SUB FOR WORK TO BE
*                   DONE
    AH 9,H4        POINT TO LINK
    LH 9,0(0,9)   LOAD REG 9 TO POINT TO
*                   NEXT ITEM IN STACK
    CH 9,H0TM     IS IT BOTTOM OF STACK =
    HCR 2,10     YES, RETURN FROM SUBR.
    HC 15,STAR
    END STAR

```

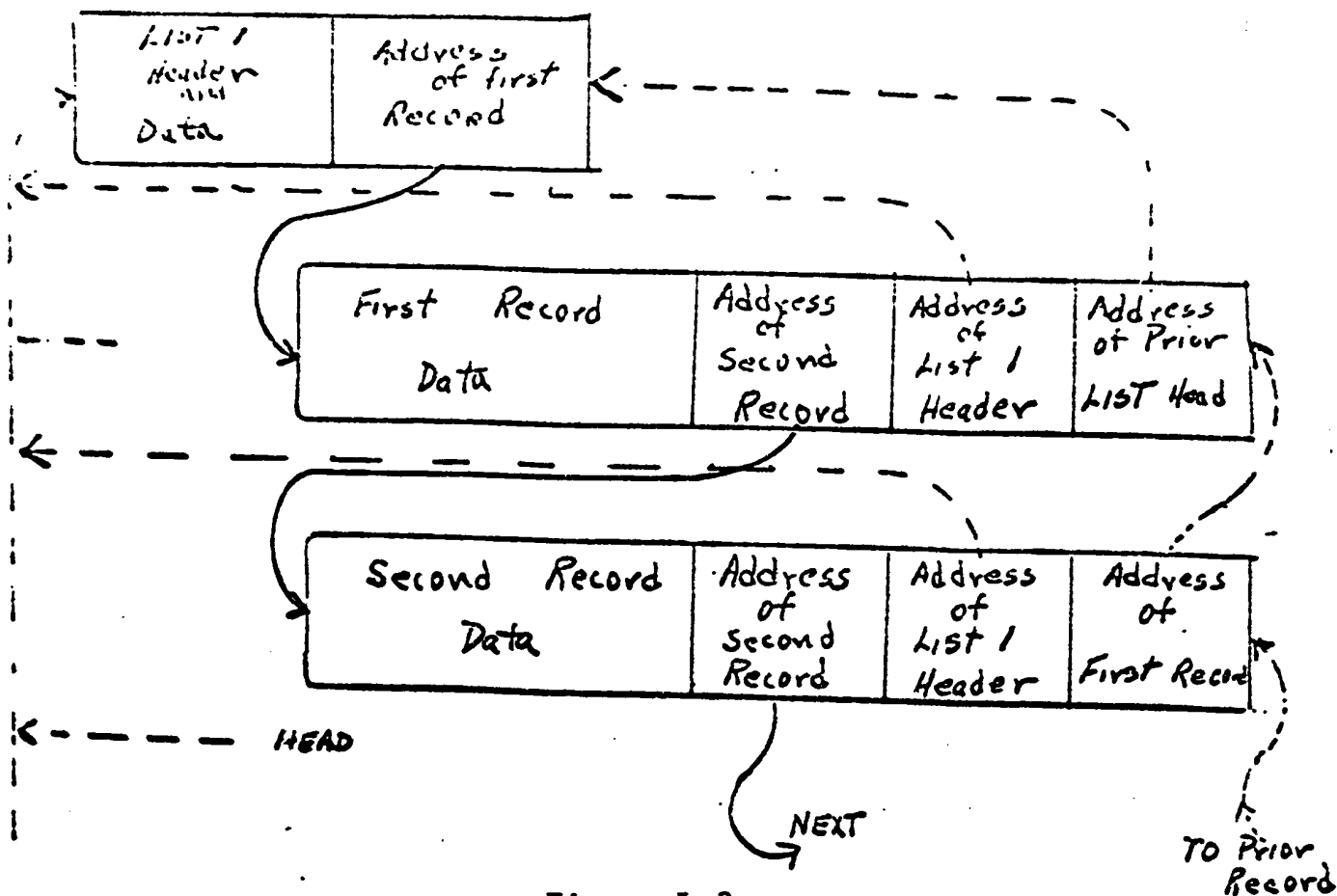
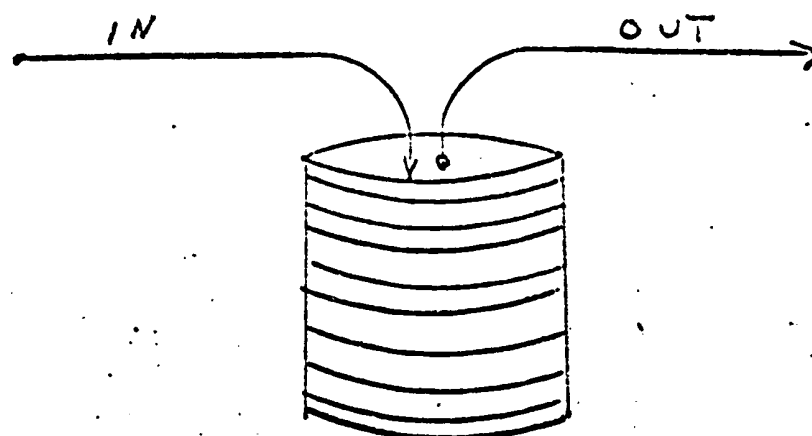


Figure L-3

Refer to the Swanson Study to Chart #9, which depicts the Bill of Material retrieval. Also refer to File Page #4, which links the various files together. The above drawing is intended to carry the ideas the two Swanson charts contain, and prepare the reader for the linking process.

Activities

1. Duplicate the drawing of the linkage of File Page #4.
2. Duplicate the drawing of Chart #9.
3. Listen to the tape cassettes that pertain to this section, and review the Basic Assembler set of instructions.



```

* ADD TO TOP OF A STACK
*
* REGISTER 8 POINTS TO AVAILABLE STORAGE
* REGISTER 9 POINTS TO TOP OF STACK
* REGISTER 10 HAS RETURN ADDRESS FROM SUBR.
*
A      START 0
      USING *,0
STAR  CH      8,LIMIT      ANY MORE AVAILABLE SPACE
      HC      2,OVFL      NO, BRANCH TO OVERFLOW
      MVC      0(4,8),DATA  STORE DATA IN AVAIL.
      STH      8,TEMP      SAVE POINTER FOR NEW TOP
*                                OF STACK
*      AH      8,H4      REG. 8 POINTS TO LINK OF
*                                AVAIL.
*      STH      9,0(0,8)  AVAIL LINKED TO TOP OF
*                                STACK
*      LH      9,TEMP      REG. 9 POINTS TO LINK OF
*                                AVAIL
*      AH      8,H2      REG. 8 POINTS TO NEW
*                                AVAILABLE STORAGE
*      BCR      15,10     RETURN FROM SUBROUTINE
      END      STAR

```

Figure L-4

The disadvantage of the extra space used by the links may be circumvented by including several items to each link, and a link need not take a large amount of space. It would take only one byte on a byte oriented computer. It is also possible to create an overlap of tables and share common parts. (This will be discussed later.)

This example of a one way link assumes that in any type of search, the direction can proceed forward only.

Delete From Top of Stack

```

* REG. 9 POINTS TO LINK OF ITEM A
* REG. 10 HAS RETURN ADDRESS FROM SUBROUTINE
*
A      START 0
      USING *,0
STAR  AH      9,H4      POINT TO LINK OF A
      LH      9,0(0,9)  REGISTER 9 POINTS TO B
*      CH      9,R01TM   ( NEW TOP OF STACK )
      BC      2,UNFL    BOTTOM OF STACK, =
      BCR     15,10     YES, GO TO UNDERFLOW
      END     STAR      RETURN FROM SUBROUTINE

```

Figure L-5

Delete From Inside a Stack

The deletion of an item (or items) is easy, since the deletion requires a change in the link to address 13. The same change with a sequential file would entail a deletion of the item in address 12, and a move up of the following items.

```

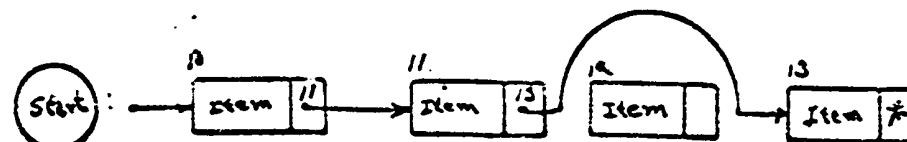
* REG. 9 POINTS TO LINK OF ITEM A
* REG. 10 HAS RETURN ADDRESS FROM SUBROUTINE
* REG. 11 IS USED AS A TEMPORARY REGISTER
*
A   START 0
    USING *,0
STAR AH  9,H4   POINT TO LINK OF A
    LH  11,0(0,9) REG. 11 POINTS TO B
    AH  11,H4   POINT TO LINK OF B
    MVC 0(2,9),0(11) LINK OF A POINTS TO C
    SH  9,H4   POINT TO A
    CH  11,BOTM BOTTOM OF STACK =
    BC  2,UNFL YES, GO TO UNDERFLOW
    BCR 15,10  RETURN FROM SUBROUTINE
    END STAR

```

Figure L-6

Insert Into Stack (or string) Between B and C

The insertion of an item into a list is easy because the inserted node can reside at any available location.



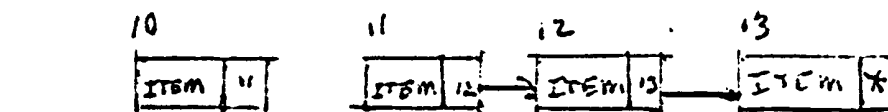
Address	Item	Link
10	Item	11
11	Item	13
12	Item	*
13	Item	*


```

*
* REG. 8 POINTS TO AVAILABLE STORAGE
* REG. 9 POINTS TO LEFT LINK OF ITEM C
* REG. 10 HAS RETURN ADDRESS FROM MAIN SUBR.
*
A      START 0
      USING *,0
STAR  CH  8,LIMT  ANY MORE AVAILABLE STORAGE
      RC  2,OVFL  NO, GO TO OVERFLOW
      MVC 0(2,8),0(9) STORE RIGHT LINK TO B
*
*           INTO LEFT LINK OF AVAIL
      AH  8,H2    POINT TO DATA OF AVAIL
      MVC 0(2,8),DATA STORE DATA IN AVAIL
      AH  8,H2    POINT TO RIGHT LINK OF AVAIL
      STH 9,0(0,8) STORE LINK TO C INTO
*
*           RIGHT LINK OF AVAIL
      AH  8,H2    POINT TO NEW AVAILABLE
*
*           STORAGE
      MVC 0(2,9),0(8) STORE LINK TO AVAIL
*
*           INTO LEFT LINK OF C
      SH  8,H6    POINT TO LEFT LINK OF AVAL
      LH  9,0(0,8) REG. 9 WILL POINT TO
*
*           RIGHT LINK OF B
      STH 8,0(0,9) RIGHT LINK OF B WILL
*
*           POINT TO LEFT LINK OF AVAL
      AH  8,H4    POINT TO RIGHT LINK OF AV.
      LH  9,0(0,8) REG. 9 WILL POINT TO
*
*           LEFT LINK OF C
      AH  8,H2    POINT TO NEW AVAILABLE ST.
      BCR 15,10  RETURN FROM SUBROUTINE
      END  STAR

```

Delete From Left End of List



```

*      DELETE FROM THE LEFT END OF A LIST
*
*      REG. 9 POINTS TO THE LEFT END OF THE LIST
*      REG. 10 HAS RETURN ADDRESS FROM MAIN SUBR.
*
Δ      START 0
        USING #,0
STAR  AH  9,H4    POINT TO RIGHT LINK OF A
        LH  9,0(0,9) POINT TO NEW LEFT END OF
*
*      CH  9,REND IS IT THE RIGHT END OF
*
*      HC  2,UNFL YES, GO TO UNDERFLOW
        BCR 15,10 RETURN FROM SUBROUTINE
        END  STAR

```

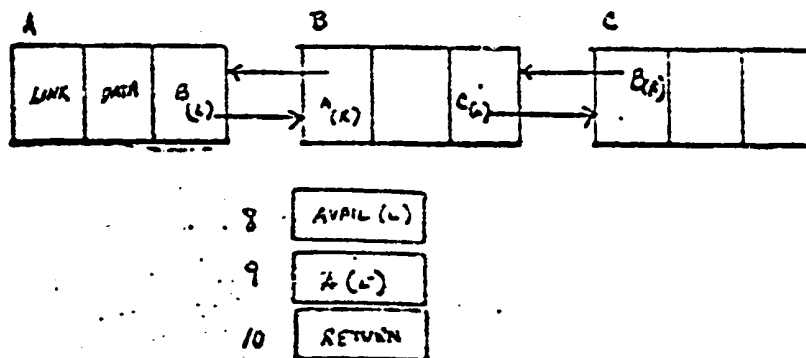
Go Through a Doubly Linked List

Reg 9---points to left end of the list

Reg 10---contains the return address from the main subroutine

Reg 11---contains the return address from the subroutine Visit.

(This subroutine performs some function and then returns us to the main program or subroutine we were in before the jump to Visit.)



```

*
* REG. 9 POINTS TO THE LEFT END OF THE LIST
* REG. 10 HAS RETURN ADDRESS FROM MAIN SUBR.
* REG. 11 HAS RETURN ADDRESS FROM SUB. VIST
*
A      START 0
      USING *,0
STAR  AH  9,H2    POINT TO DATA
      HAS  11,VIST GO TO SUBR. FOR WORK TO BE
*
*
*      AH  9,H2    POINT TO RIGHT LINK
*      LH  9,0(0,9) LOAD REG. 9 TO POINT TO
*
*      CH  9,REND  IS IT THE RIGHT END OF
*
*
*      BCR 2,10   YES, RETURN FROM SUBR.
*      BC  15,STAR NO, CONTINUE THROUGH LIST
*      END  STAR

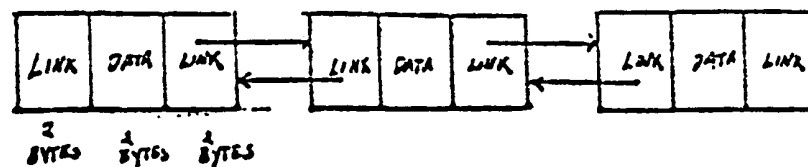
```

Add to Left End of Doubly Linked List

Reg 8---points to available storage

Reg 9---points to left end of list

Reg 10---contains the return address from the subroutine



```

*      ADD TO LEFT END OF A DOUBLY LINKED LIST
*
*      REG. 8 POINTS TO AVAILABLE STORAGE
*      REG. 9 POINTS TO THE LEFT END OF THE LIST
*      REG. 10 HAS RETURN ADDRESS FROM MAIN SUBR.
*
A      START 0
      USING #,0
STAR  CH  8,LIMIT  ANY MORE AVAILABLE STORAGE
      BC  2,OVFL  NO, GO TO OVERFLOW
      STH 8,TEMP  SAVE POINTER FOR NEW LEFT
*
      AH  8,H2    POINT TO DATA OF AVAIL
      MVC 0(2,8),DATA  STORE DATA IN AVAIL
      STH 8,0(0,9) STORE LINK THAT POINTS
*
      LH  9,TEMP  POINT TO NEW LEFT END OF
*
      AH  8,H2    POINT TO NEW AVAILABLE
*
      BCR 15,10  RETURN FROM SUBROUTINE

```

Figure L-8

Doubly Linked Lists

Reg 8---points to left link of available storage

Reg 9---points to left link of an item in list

Reg 10---contains the return address from the subroutine

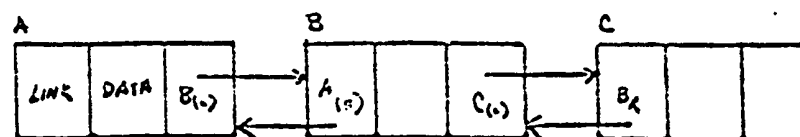
Insert Into List Between B and C (when Reg 9 points to left link of G)

```

*   INSERT INTO A LIST BETWEEN ITEMS B AND C
*
*   REG. 8 POINTS TO AVAILABLE STORAGE
*   REG. 9 POINTS TO LEFT LINK OF ITEM C
*   REG. 10 HAS RETURN ADDRESS FROM MAIN SUBR.
*
A   START 0
    USING *,0
STAR CH  8,LIMIT ANY MORE AVAILABLE STORAGE
    AC   2,OVFL NO, GO TO OVERFLOW
    MVC  0(2,8),0(9) STORE RIGHT LINK TO B
*
    AH  8,H2 POINT TO DATA OF AVAIL
    MVC  0(2,8),DATA STORE DATA IN AVAIL
    AH  8,H2 POINT TO RIGHT LINK OF AVAIL
    STH 9,0(0,8) STORE LINK TO C INTO
*
    AH  8,H2 POINT TO NEW AVAILABLE
*
    MVC  0(2,9),0(8) STORE LINK TO AVAIL
*
    SH  8,H6 POINT TO LEFT LINK OF AVAIL
    LH  9,0(0,8) REG. 9 WILL POINT TO
*
    STH 8,0(0,9) RIGHT LINK OF B WILL
*
    AH  8,H4 POINT TO RIGHT LINK OF AVAIL
    LH  9,0(0,8) REG. 9 WILL POINT TO
*
    AH  8,H2 POINT TO NEW AVAILABLE ST.
    BCR 15,10 RETURN FROM SUBROUTINE
    END  STAR

```

Delete From Inside of List



```

REG 8  [AVAIL]
REG 9  [A(L)]
REG 10 [RETURN]
REG 11 [TEMP]

```

```

*      DELETE FROM INSIDE OF A LIST
*
*      REG. 9 POINTS TO THE LEFT END OF THE LIST
*      REG. 10 HAS RETURN ADDRESS FROM MAIN SUBR.
*
A      START 0
      USING #,0
STAR  STH  9,TEMP  SAVE REG. 9
      AH   9,H4    POINT TO RIGHT LINK OF A
      LH   11,0(0,9) POINT TO LEFT LINK OF B
      AH   11,H4    POINT TO RIGHT LINK OF B
      MVC  0(2,9),0(11) RIGHT LINK OF A WILL
*                               POINT TO LEFT LINK OF C
*      LH   9,0(0,11) REG 9 WILL POINT TO THE
*                               LEFT LINK OF C
      SH   11,H4    POINT TO LEFT LINK OF B
      MVC  0(2,9),0(11) LEFT LINK OF C WILL
*                               POINT TO RIGHT LINK OF A
*      CH   9,REND  IS IT THE RIGHT END OF
*                               THE LIST
      BC   2,UNFL  YES, GO TO UNDERFLOW
      LH   9,TEMP  RESTORE REG. 9 0 POINT
*                               TO LEFT LINK OF A
      BCR  15,10  RETURN FROM SUBROUTINE
      END  STAR

```

Figure L-9

Activities

1. Flowchart example #1 "Go Through Stack"
2. Flowchart "Add To Top of Stack"
3. Flowchart "Delete From Top of Stack"
4. Flowchart "Insert Into Stack Between B and C"
5. Flowchart "Add to Left of Doubly Linked List"
6. Flowchart "Delete From Inside of List"
7. Flowchart "Doubly Linked Lists"

Trees

Lists and strings are understood easier, developing the hierarchical structure idea with trees.

Lists**** (((d))) (Q, R,), L, () (PTA)

The comma and parentheses become quite important in depicting atoms from the list.

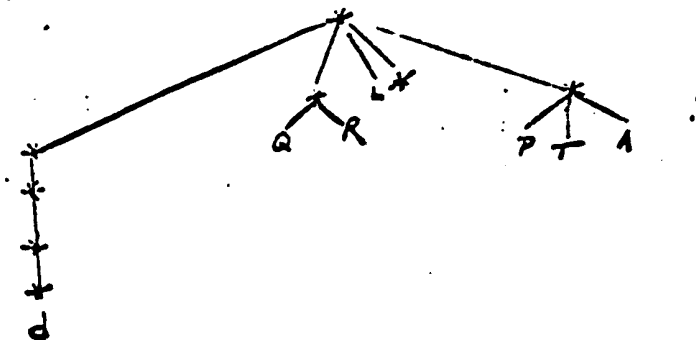


Figure T-1

The asterisks help to define the recursion. Here, the list contains the list (((d))) which consists of the list ((d)), which consists of list (d), which consists of the atom d. The asterisks indicate a blank or null.

This idea exists in the Dewey Decimal System for the libraries:

- 000-099 General Works
- 100-199 Philosophy, Psychology, Ethics
- 200-299 Religion and Mythology
- 300-399 Sociology (Economics, Civics, Education, Vocations)
- 400-499 Philology (Language, Dictionaries, Grammar)
- 500-599 Science
- 600-699 Useful Arts (Medicine, Engineering, Agriculture, Radio, Aviation, etc.)
- 900-999 History, Geography, Biography, Travel

Each of these ten main classes is broken up into more specialized fields. For example, class 600-699, Useful Arts, is broken into ten classifications including Medicine, Engineering, and Manufacturing. Each of these divisions is further subdivided.

This hierarchical tree idea exists in the paging scheme of this text. For example:

- 1.1
- 1.2
- 1.1.2

The idea also exists in set theory. For example it exists in nested sets:

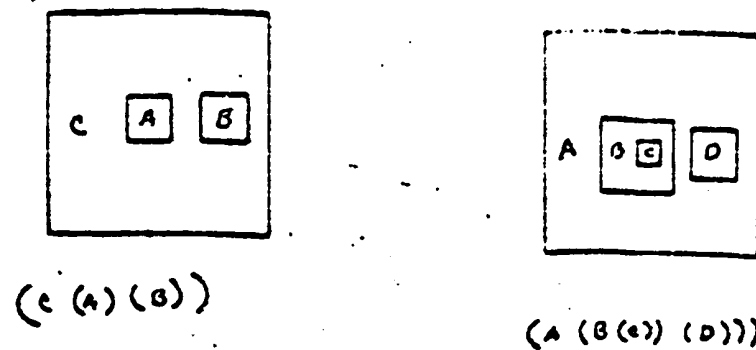


Figure T-2

A Combination of Sets and Trees

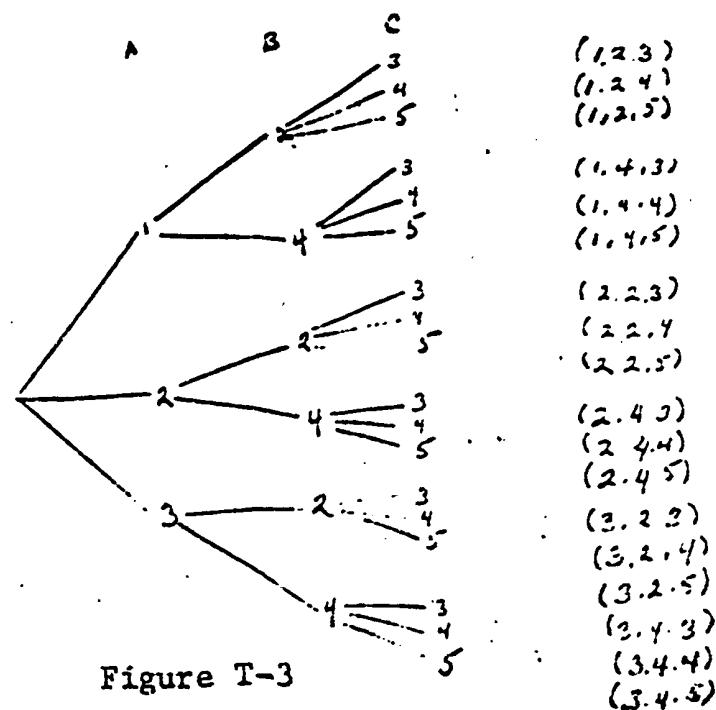


Figure T-3

The hierarchical idea allows us to combine sets and trees. For

example:

Let A -- the set 123

Let B -- the set 242

Let C -- the set 345

With the help of a hierarchical tree, we can find $A \times B \times C$.

We see that $A \times B \times C$ consists of the ordered triplets to the right of the tree.

Tree Hierarchical Structures

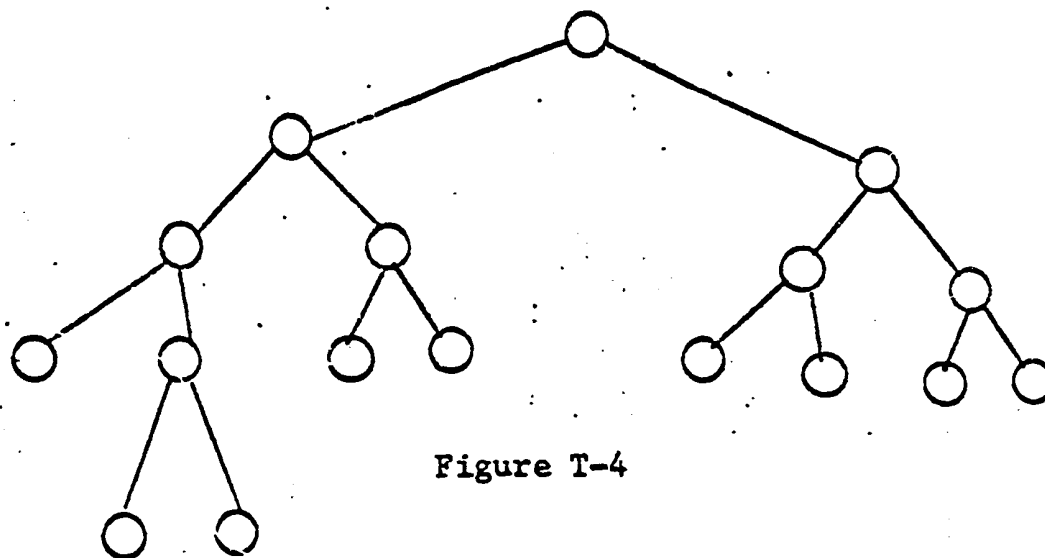


Figure T-4

The idea of a tree is presented here to form a concept of structure that might exist within core memory.

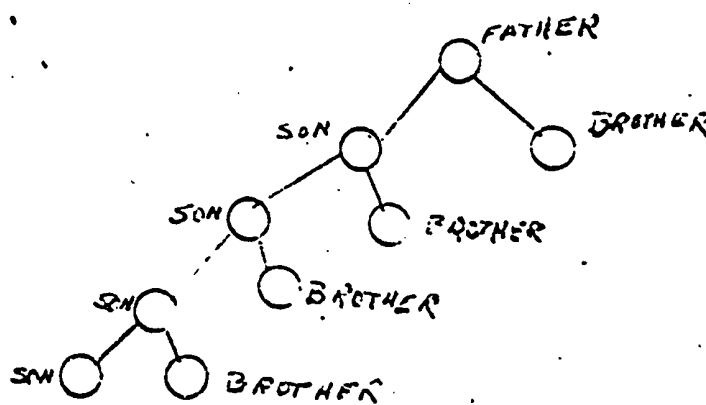


Figure T-5

2.3.1 Hierarchical Ring Structure

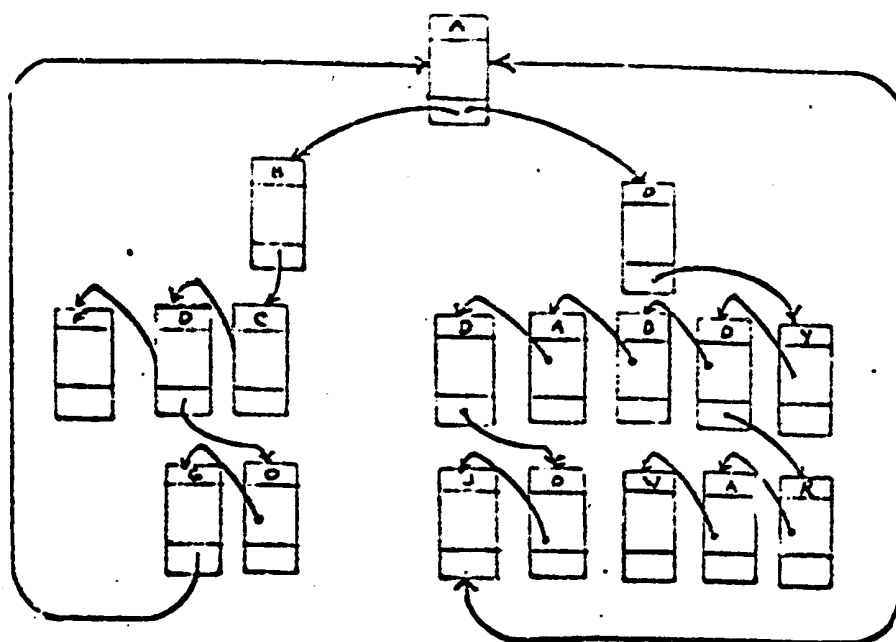


Figure H-1

The tree nodes are divided into three parts, which can be adjusted to fit most computers. This search procedure scans the descendants of the head (root) by following the chain of part two of each element (the pointer in part two). If a match is found on one level a branch can be made to the next level by means of the address of the pointer of part three of the element.

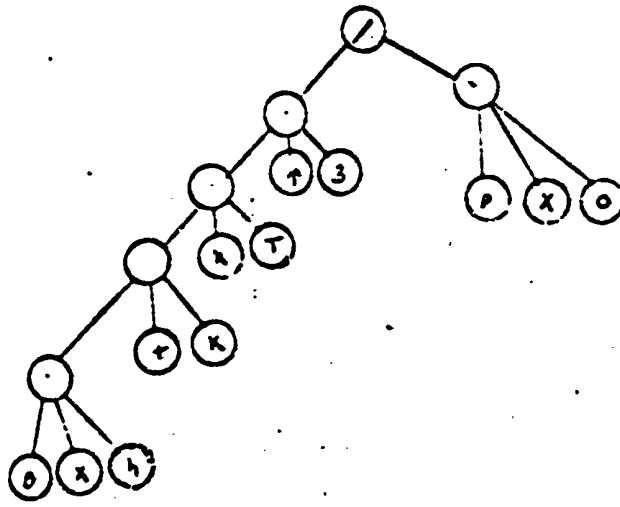
It is easy to insert a new node here because the chain can be broken and then the pointers can be changed to effect a new link.

This arrangement of the node, Figure H-1, amounts to the insertion of an atom in the first element of the node. The second part links to the brother on the same level, and the third part of the node links to the next structural level of descendancy.

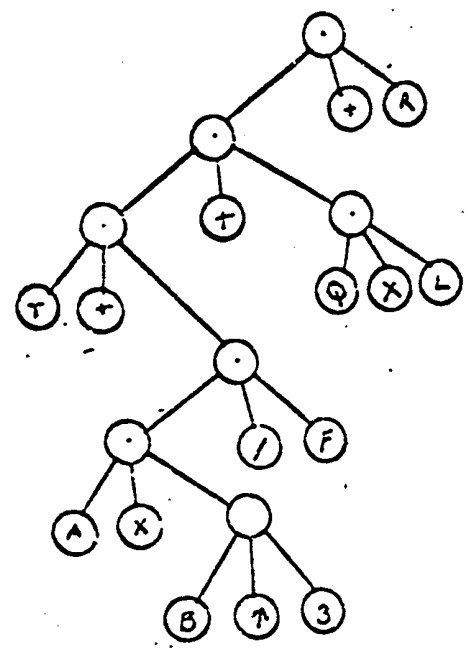
The complete ring is made, which allows the return to the head of the list, and at the same time facilitates entering the ring at any point.

Algebraic Formulas

Algebraic formulas are better understood with the help of a tree hierarchy structure.



$$(b \times h + k) + 3 / (p \times o)$$



$$(t + a \times b^3 / f + q \times l + r)$$

Figure H-2

Double Ring Link to Tree

This type of ring structure represents a double link type structure, which allows a much freer movement, but does take more space due to the pointers.

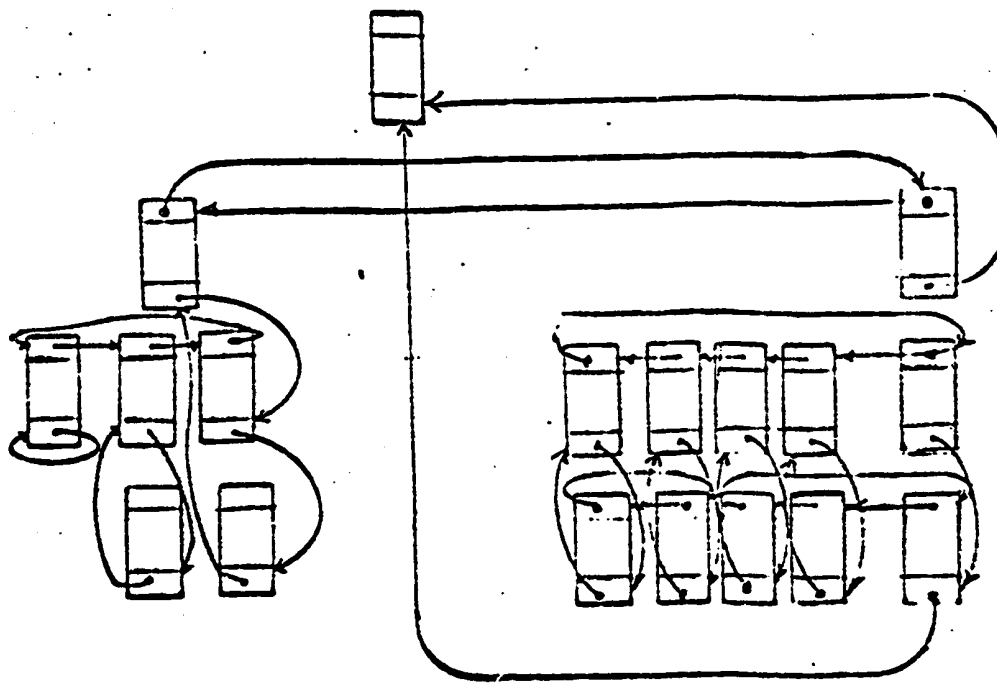


Figure H-3

Traversing Trees

Traverse a Tree in Preorder Sequence

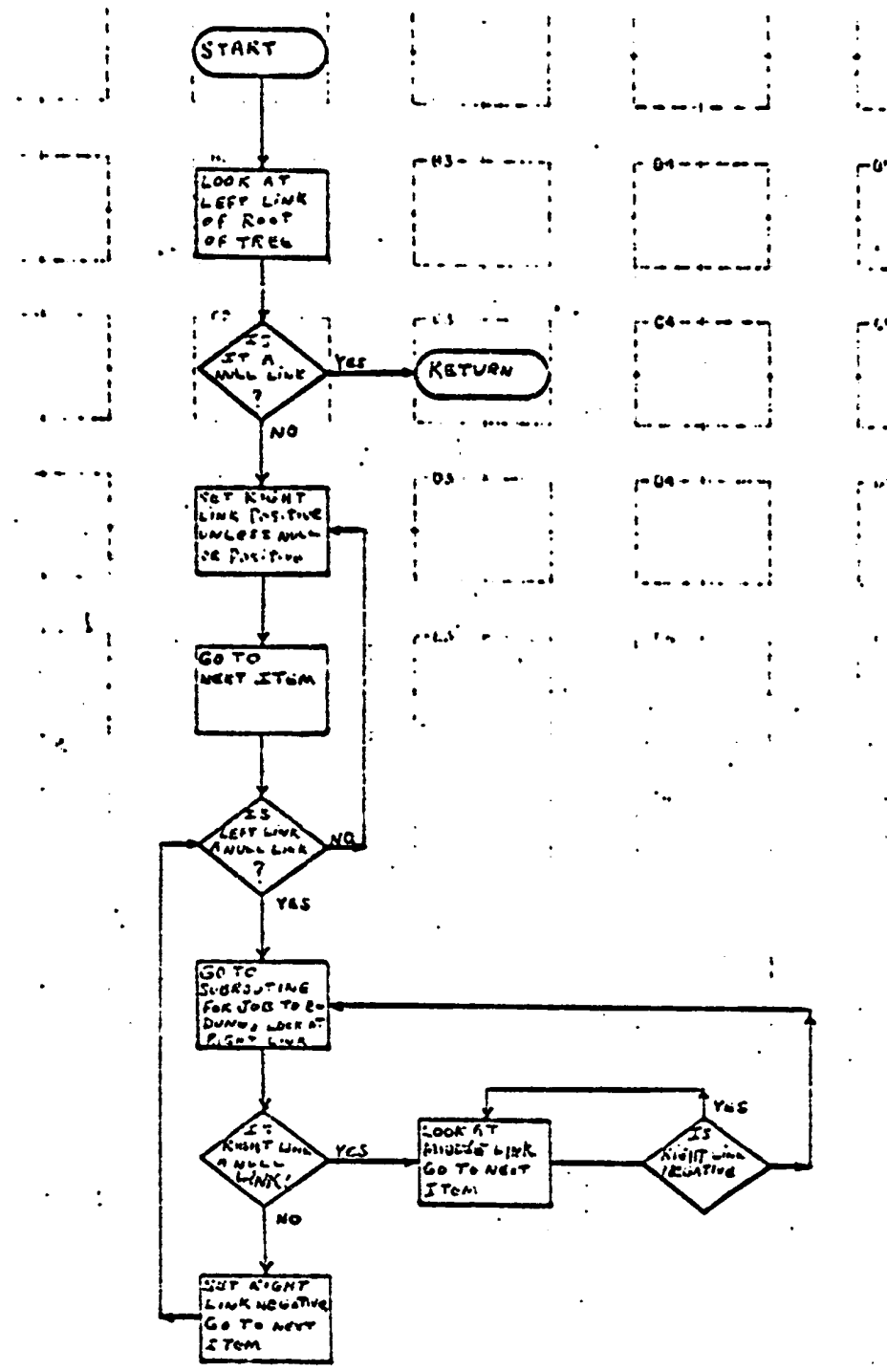


Figure TT-1

Preorder

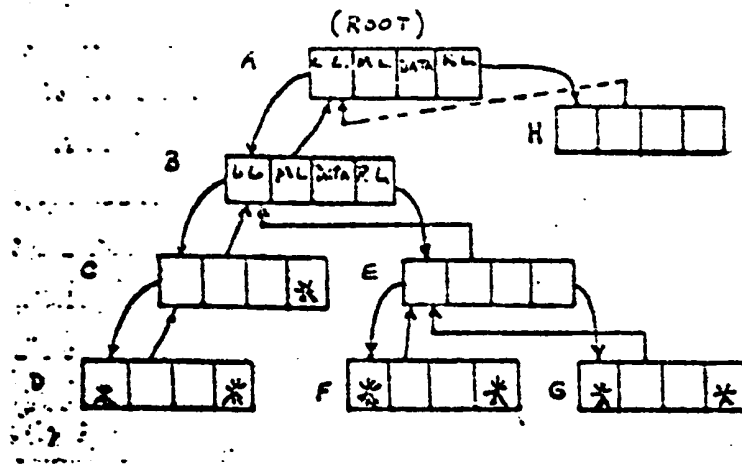


Figure TT-2

Traverse the Tree in Preorder

A, B, C, D, E, F, G, H

Left Link = 2 bytes

Right Link = 2 bytes

Middle Link = 2 bytes

Data = 2 bytes

Null Link = *

Traverse a Tree in Preorder

```

*
* REG. 9 POINTS TO LEFT LINK OF ROOT OF TREE
* REG. 10 IS USED TO HANDLE DATA
* REG. 11 IS USED AS AN AUXILLARY REGISTER
* REG. 12 CONTAINS RETURN ADDRESS FROM SUBR.
*
A      START 0
      USING #,0
*STAR  SR      11,11      CLEAR REG. 11
      LH      10,0(0,9)  LOAD REG. 10 FROM LEFT
*                               LINK OF ROOT
      CH      10,NULL    IS IT A NULL LINK =
      HCR     8,12      YES, RETURN FROM SUBR.
      AH      9,H2      NO, POINT TO MIDDLE LINK
      STH     12,0(9)   STORE RETURN ADDRESS IN
*                               MIDDLE LINK OF ROOT
      SH      9,H2      POINT TO LEFT LINK
      HAS     13,VIST   GO TO SUBR. FOR WORK
*                               TO BE DONE
*
POS    AH      9,H6      POINT TO RIGHT LINK
      LH      10,0(0,9)  LOAD REG.10 FROM R. LINK
      CH      10,NULL    IS IT NULL OR POSITIVE
      HC      11,*+6     YES
      SR      11,10     NO, MAKE R.L. POSITIVE
      STH     11,0(9)   STORE POS. NO. IN R.LINK
      SR      11,11     CLEAR REG. 11
      SH      9,H6      POINT TO LEFT LINK
      LH      9,0(0,9)  POINT TO LEFT LINK OF
*                               NEXT NODE
*
NEXT   BAS     13,VIST   GO TO SUBR. FOR WORK
*                               TO BE DONE
*
      LH      10,0(0,9)  LOAD REG. 10 FROM L.LINK
      CH      10,NULL    IS IT A NULL LINK =
      BC      7,POS     NO, GO TO MAKE RIGHT
*                               LINK POSITIVE
*
NXTR   AH      9,H6      POINT TO RIGHT LINK
      LH      10,0(0,9)  LOAD REG. 10 FROM R.LINK
      CH      10,NULL    IS IT A NULL LINK =
      HC      13,MIDL   BRANCH IF REG.10 NOT POS
      SR      11,10     NEGATE
      STH     11,0(9)   STORE THE NEGATIVE NO.
*                               IN RIGHT LINK
*
      LH      9,0(0,9)  POINT TO NEXT NODE
      HC      15,NEXT
*
MIDL   SH      9,H4      POINT TO MIDDLE LINK
      LH      9,0(0,9)  POINT TO NEXT NODE
      HC      15,NXTR
      END      STAR

```

Figure TT-3

Traverse a Tree in Post Order

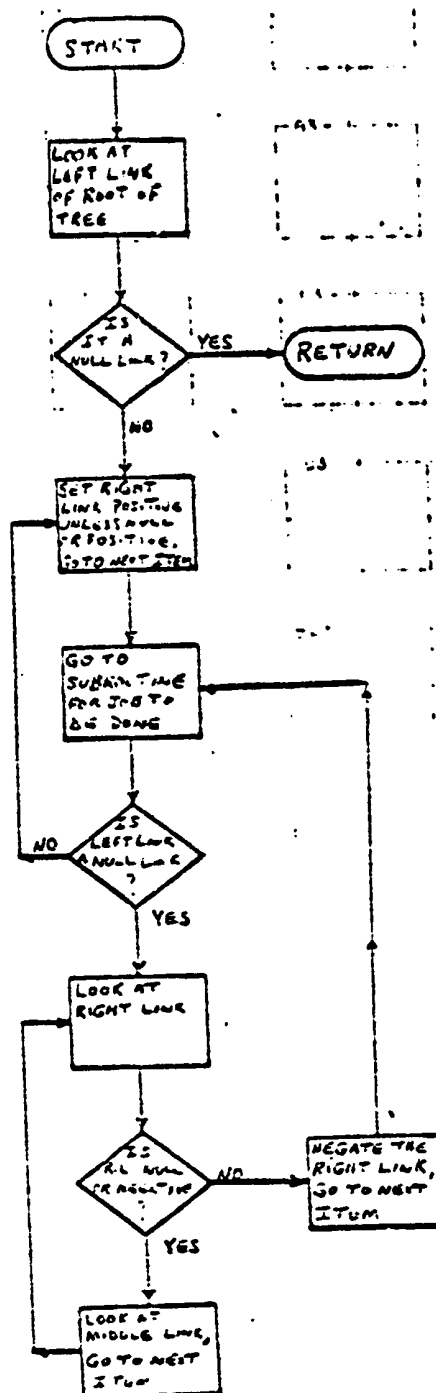


Figure TT-5

Post Order

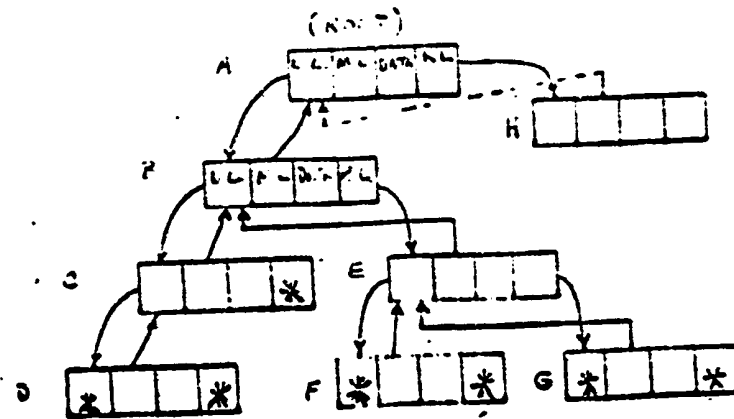


Figure TT-6

Traverse the Tree in Post Order

D, C, B, F, E, G, A, H

Traverse a Tree in Post Order

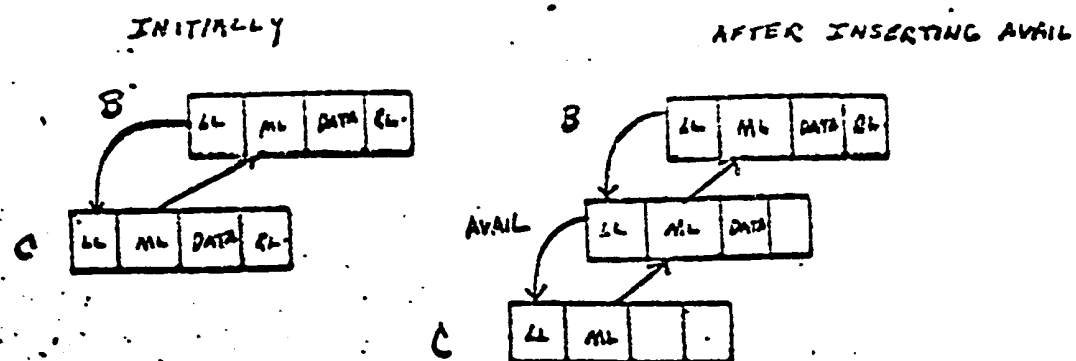
```

* REG. 9 POINTS TO LEFT LINK OF ROOT OF TREE
* REG. 10 IS USED TO HANDLE DATA
* REG. 11 IS USED AS AN AUXILLARY REGISTER
* REG. 12 CONTAINS RETURN ADDRESS FROM SUBR.
*
A      START 0
      USING *,0
STAR  SR      11,11      CLEAR REG. 11
      LH      10,0(0,9)  LOAD REG. 10 FROM LEFT
*                               LINK OF ROOT
      CH      10,NULL    IS IT A NULL LINK =
      RCR      8,12      YES, RETURN FROM SUBR.
      AH      9,H2      NO, POINT TO MIDDLE LINK
      STH     12,0(0,9)  STORE RETURN ADDRESS IN
*                               MIDDLE LINK OF ROOT
      SH      9,H2      POINT TO LEFT LINK
      AH      9,H6      POINT TO RIGHT LINK
      LH      10,0(0,9)  LOAD REG.10 FROM R. LINK
      CH      10,NULL    IS IT NULL OR POSITIVE
      RC      11,*+6     YES
      SR      11,10     NO, MAKE R.L. POSITIVE
      STH     11,0(0,9)  STORE POS. NO. IN R.LINK
      SR      11,11     CLEAR REG. 11
      SH      9,H6      POINT TO LEFT LINK
NEXT  LH      9,0(0,9)  POINT TO LEFT LINK OF
*                               NEXT NODE
      LH      10,0(0,9)  LOAD REG. 10 FROM L.LINK
      CH      10,NULL    IS IT A NULL LINK =
      RC      7,POS     NO, GO TO MAKE RIGHT
*                               LINK POSITIVE
      VISI    BAS, 13,VIST YES, GO TO SUBR. FOR
*                               WORK TO BE DONE
      AH      9,H6      POINT TO RIGHT LINK
      LH      10,0(0,9)  LOAD REG.10 FROM R.LINK
      CH      10,NULL    IS IT NULL OR NEGATIVE
      RC      2,NEG     NO, GO TO SET R.LINK NEG
MIDL  SH      9,H4      POINT TO MIDDLE LINK
      LH      9,0(0,9)  POINT TO NEXT NODE
      AH      9,H6      POINT TO RIGHT LINK
      LH      10,0(0,9)  LOAD REG.10 FROM R.LINK
      CH      10,NULL    IS IT NULL OR NEGATIVE
      RC      4,MIDL    YES, GO TO MIDDLE LINK
      SH      9,H6      NO, POINT TO LEFT LINK
      RC      15,VISI   GO FOR WORK TO BE DONE
      SH      11,10     NEGATE
NEG   STH     11,0(0,9)  STORE NEG.NO. IN R.LINK
      RC      15,NEXT   GO TO NEXT NODE
*                               NEXT NODE
*
      END) STAR

```

Figure TT-8

Add to a Tree Between B and C

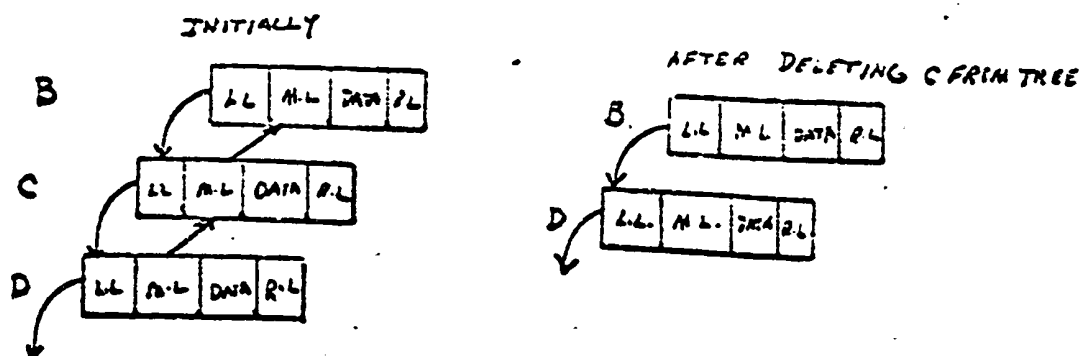


* ADD TO A TREE BETWEEN NODES B AND C
 *
 * REG. 8 POINTS TO AVAILABLE STORAGE
 * REG. 9 POINTS TO NODE B
 *

STRT	START 0	
	USING #, 0	
STAR	LH 10,0(0,9)	SAVE LEFT LINK OF B
	STH 8,0(0,9)	LEFT LINK OF B WILL
*		POINT TO AVAIL
	STH 10,0(0,8)	LEFT LINK OF AVAIL
*		WILL POINT TO C
	LH 9,0(0,8)	REG. 9 POINTS TO
*		LEFT LINK OF C
	AH 9,H2	REG. 9 POINTS TO
*		MIDDLE LINK OF C
	LH 10,0(0,9)	SAVE MIDDLE LINK OF C
	STH 8,0(0,9)	MIDDLE LINK OF C POINT
*		TO LEFT LINK OF AVAIL
	AH 8,H2	REG. 8 POINTS TO MIDDLE
*		LINK OF AVAIL
	STH 10,0(0,8)	MIDDLE LINK OF AVAIL
*		POINTS TO L.L. OF B
	END STAR	

Figure TT-9

Delete Node C From A Tree



```

*   DELETE NODE C FROM A TREE
*
*   REG.9 POINTS TO LEFT LINK OF B
*
STRT  START 0
      USING *,0
STAR  LH      8,0(0,9)   REG.8 POINTS TO L.L. C
      MVC     0(2,9),0(8) L.L. OF B WILL POINT
*                                     TO LEFT LINK OF D
      LH      9,0(0,8)   REG.9 POINTS TO L.L. D
      AH      8,H2       REG.8 POINTS TO M.L. C
      AH      9,H2       REG.9 POINTS TO M.L. D
      MVC     0(2,9),0(8) M.L. OF D POINTS TO
*                                     LEFT LINK OF B

      END    STAR

```

Figure TT-10

2.3.3 STRUCTURED TREE DIRECTORY

This section, which includes Random Organization, and the complete File Section, which follows, are an integral part of an information management system.

The directory, files and type of structure, which is chosen for each, will have a great bearing on the speed and efficiency of the system. In the Swanson study, one approach was used. These two parts of the book will allow us to widen our area of selection of approaches we might choose.

We have developed the hierarchic structure concept thus far in the book. Now we will make use of it and others in developing index systems to files for search and retrieval purposes. As with the Swanson study, these methods are slanted toward disk systems.

2.3.3.1 STRUCTURED TREE DICTIONARY

It is important to note that each node (leaf) represents a complete record in the system. The use of three letters for a key is used here as an example. Actually a three-level tree can accommodate several million keys, if there are enough characters in the addressing method. Addressing method here means the KEY/ADDRESS/LIST LENGTH, which amounts to a multiword method.

The first level of the tree is maintained in core storage, and the other two levels are maintained in auxiliary storage. A typical disk notation is used in which T12 means track 2 of disk 1.

The key fragments across the bottom of the tree stand for complete names, but were truncated in order to form a fixed length key

format. This fixed length format could be imposed by the hardware or could be the choice of the programmer. Any number of letters could be used, but as mentioned earlier, the number of letters in the key affects the number of records possible. Three were used here for the sake of brevity.

The search format for the tree is the one mentioned above:

NAME/ADDRESS/LIST LENGTH

The name, address and list length exist in auxiliary storage, and the addressing scheme proceeds from low to high as do the IBM 360 computers.

Level 1 resides in core storage, and it directs us to level 2, which resides in auxiliary storage. This address at level 2 is a track address. The 0 denotes that it is track 0 and 1 denotes that cylinder 1 is used. The highest alphabetic letter here is denoted by COW/T12/*, which signifies that COW marks an upper limit for this section of the tree. The asterisk denotes that there is no list length here, since the branch is not made to the actual list as yet. Level 2 then takes us to T12, which resides on level 3. This level contains the actual lists. T12 takes us to track 2 of cylinder 1. This level contains COW, which is the name part of the address we want. The actual address and length of the record is given here.

The process of decoding this type of dictionary proceeds by the following manner. For example, assume that we have the word COWBOY. The first three letters will be truncated (COW). This key is then compared to the dictionary file in core storage. It is smaller than HOT so COW is selected. This key takes us to track T10, track 0, disk 1, and then proceeds to compare at this level. COW is larger than CAM,

but it does match COW.

This causes a branch to track 2 of disk 1 (T12), where a match is made. The link address takes us to the actual address of COWBOY, and we are given the length of the record that resides at that location.

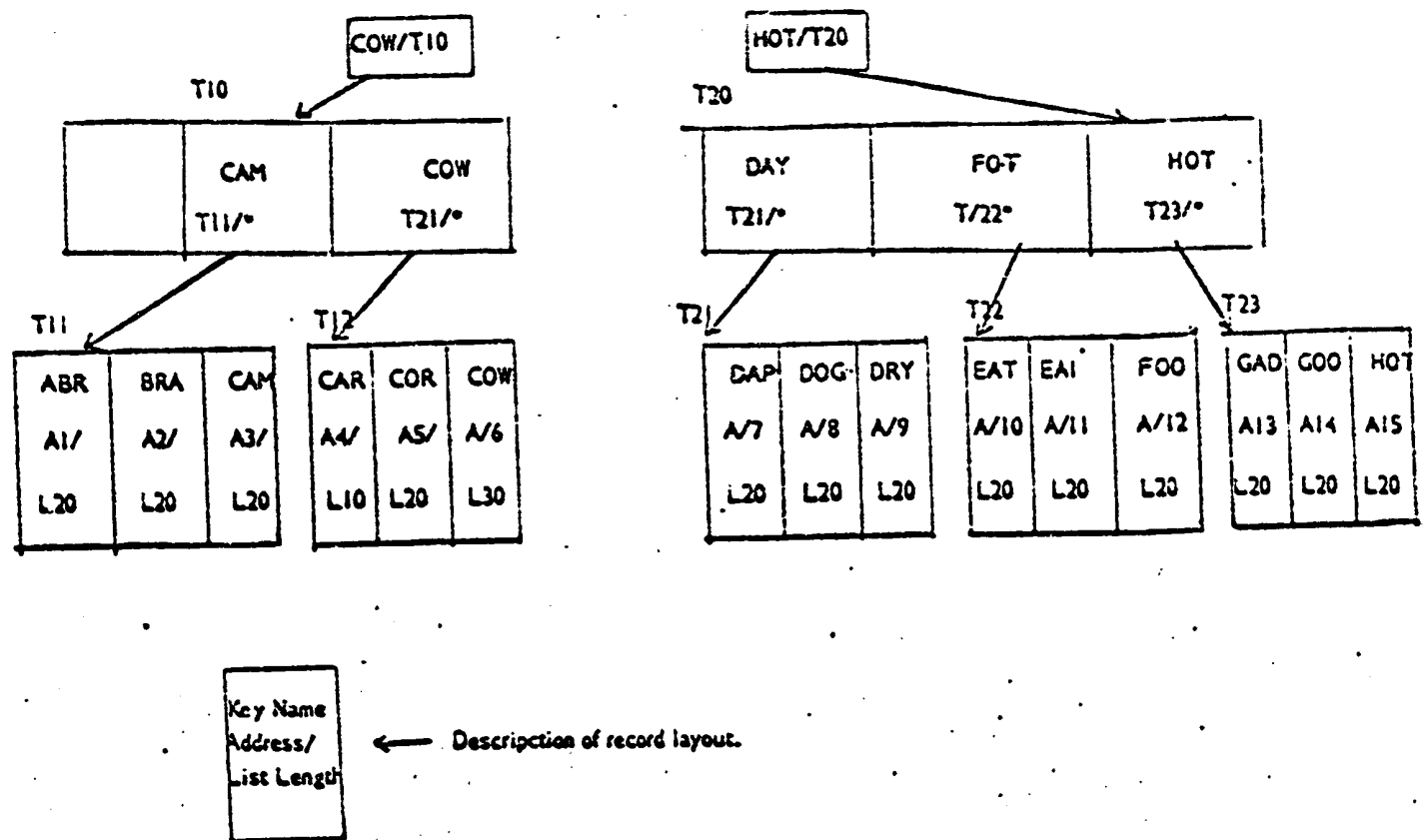


Figure STD-1

2.4 FILES

The file organization can be divided into three major groups:

1. Sequential File Organization
2. Inverted Files (or inverted lists)
3. String structures

The other types of organizational methods listed in this section

can be considered to be combinations of the basic three methods.

A slight emphasis is placed on the Inverted File for reasons of comparison to the RPG methods which occur later in the book.

Figure F.1. depicts the structural organization from the most common at the base, Sequential Organization, to the more sophisticated and complex at the top.

The idea of file partitioning is introduced in this section. The partitioning concept is intended to make access to the disk records faster. The two basic methods of partitioning we will use are Inverted Lists and Multiple Threaded Lists. Finally, we will introduce combinations of these methods, and develop a Page Partition Multilist.

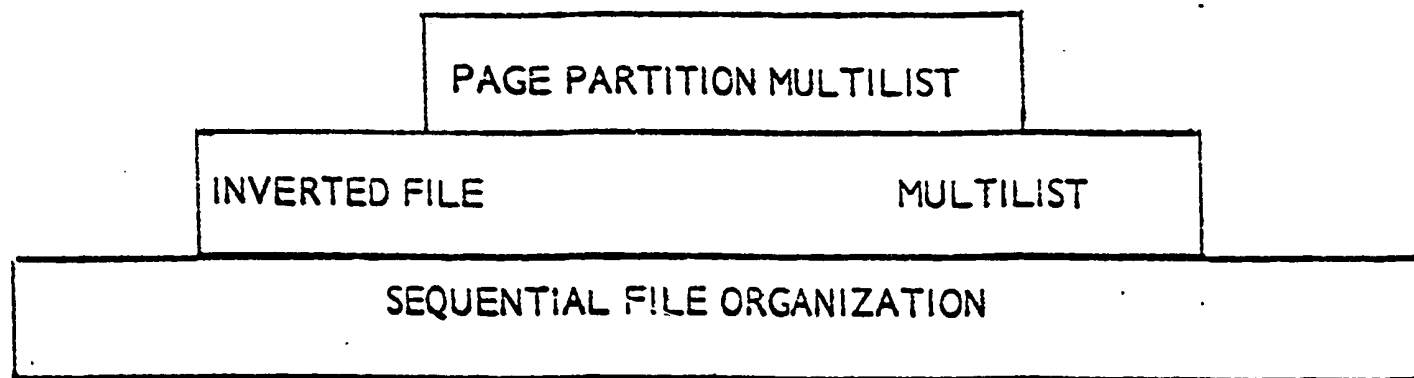


Figure F-1

The Inverted File and the Multilist exist at the extremes of a continuum as depicted by Figure F.1. They represent opposing or contrasting methods of file organization for search and retrieval purposes. The other methods represented here were developed as faster methods were needed. The Multilist and Page Partitioned Multilist evolved as

a result of the use of file structures in the auxiliary and core storage. As will be seen in a later part of the book, the Virtual Memory techniques evolved at a later date.

2.4.1 SEQUENTIAL FILES

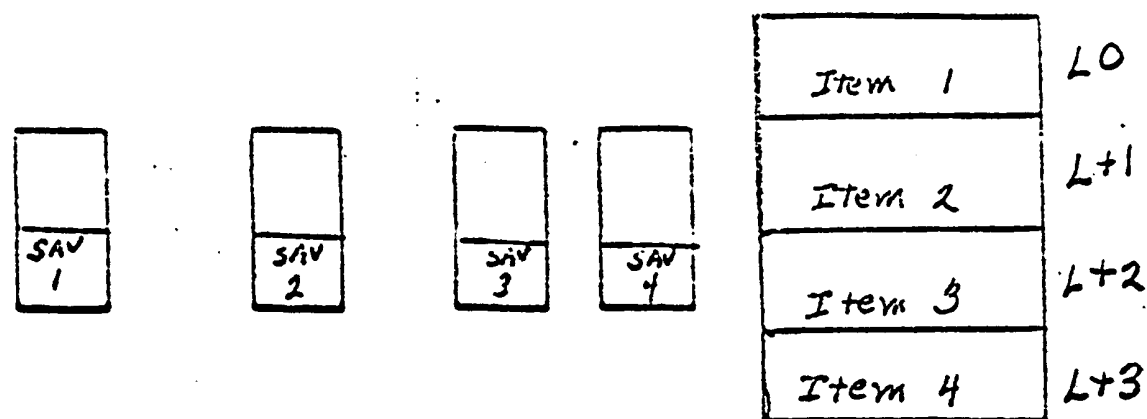


Figure F-2

Sequential data organization is the most common type used. For example, an inventory file is made up of many records, and each record in turn contains many fields:

Stock number
 Unit of measure
 On-hand quantity
 Material description
 On-order-quantity

The relationship of these fields of the record is that they do relate to a particular stock item. An inventory record for a hand saw

would have a stock number for hand saws, a unit of measure, a material description of hand saws, and a measure for which the information about on-hand, on-order and reorder quantities for hand saws is kept.

The sequential file has certain advantages in that there is a fast access for each relationship. There is a disadvantage when a file is to be searched until a record having a particular key is found. This requires an examination of the first record, if the key is wrong, the next record, etc. This process goes on until the correct record is found. The updating process is also rather difficult with a sequential file. If the new record is shorter or longer than the original record in the file, the adjacent records may be affected or destroyed. The update process then becomes very costly when there is only one record to update. The update process is generally used when there are a number of records to update.

If we have a file of records, it becomes necessary to discover certain features that exist in the file. For example, if a saw is sold, the inventory record for saws needs to be isolated so that the value of the on-hand quantity field can be reduced.

The common attribute here is called a key. By selecting a different key for a file and sorting on the basis of that particular key, the sequence of records in the file may change, but we can obtain the desired information.

It is not necessary to store records with keys. There are times when the sequential arrangement is based solely on the arrival within the system.

Activities

Linkages

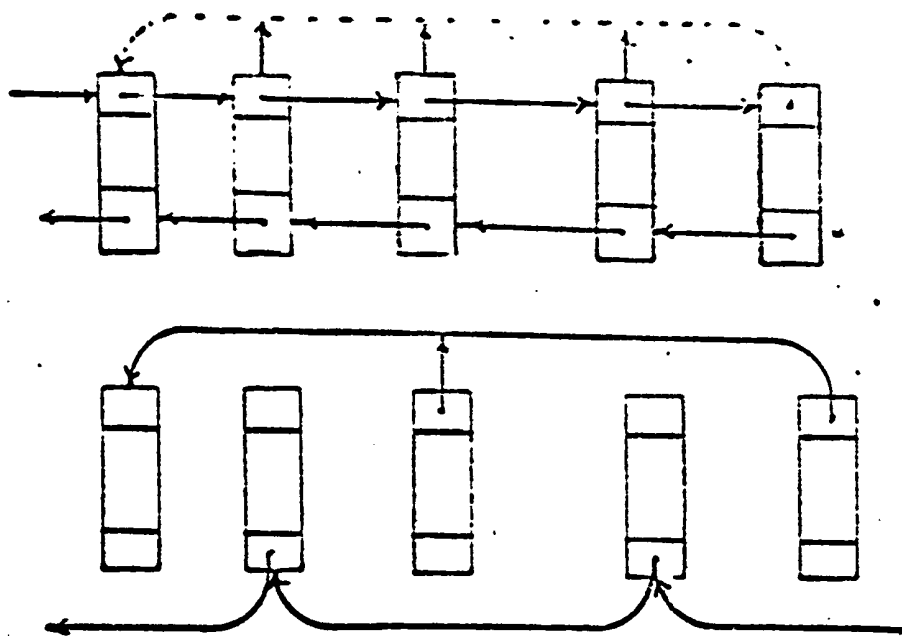


Figure A-4

Both these methods can be used for the string structure of a file.

1. Compare them to the Swanson Study.
2. Compare them to the linkages in Section II of Data Structures.
3. Where do the actual linkage processes take place, in core or on disk?

2.4.2 INVERTED FILES

KEY	AD(1,1)
KEY	AD(1,1)
KEY	AD(2,1)
KEY	AD(3,4)
KEY	AD(7,1)

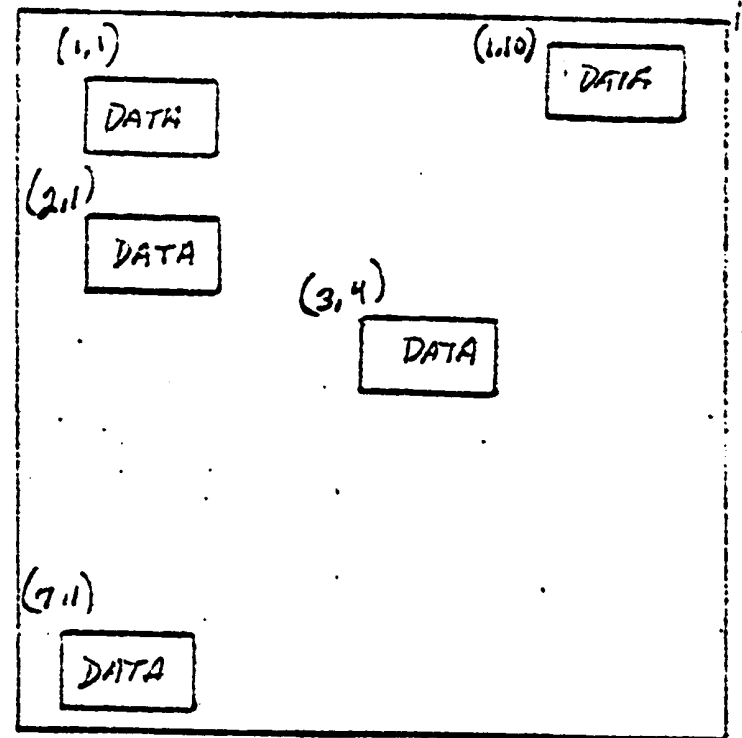


Figure IF-1

The above example IF-1 is a drawing of the Inverted File structure taken to its ultimate, because the list has a length of 1. This type of structure requires the index to take up as much or more room as the list. The pointer AD(1,1) is a disk address to take us to track 1 of cylinder 1 to a particular location, which has a list length of 1.

A comparison could be made between State and City to determine which ones belong together. This could be done without linking to the list, which means a comparison of the indexes.

The main disadvantage to this method is the large directories needed. It works best as a partially inverted file combined with a random method, such as a hierarchical tree method or sequential method. This makes it possible to invert on just a few keys.

2.4.2.1 CREATION, MAINTENANCE and DELETION OF KEYS

If there is a master file, an inverted file may be created from it. This entails searching through the master file, making a directory from the selected keys, and applies to both the inverted file and the multilist file.

Every time a record is added to the master file, it becomes necessary to update the directories it pertains to.

The Inverted File Algorithm:

1. After the record is entered in the master file, prepare it for entry into the inverted file.
2. Assign the auxiliary address AAD, which may entail some form of the list of available space.
3. Key n in the directory is decoded to the variable length inverted file index.
4. Place Key n in its proper sequence into the directory; if the space is exhausted, use a link to the next page or block.
5. Add one to the list length in the index.
6. Continue steps 3 to 5 for all keys of the new record.
7. Store the new auxiliary address at AAD.

Deletion of Keys

The deletion of keys causes no problem, as far as overflow is concerned, since there is a reduction of space being used.

The Deletion algorithm:

1. Retrieve record from auxiliary storage and delete Key n.

2. Decode Key n in the proper directory of the inverted file list.
3. Delete the address from the list (this address is the address AAD residing in the inverted file list).
4. Repeat steps 2 and 3 until all keys of the record have been deleted that are required to be deleted.

Deletion of Records

The Record Deletion algorithm:

1. Transfer record from auxiliary storage location to core.
2. Set the record delete bit.
3. Decode every key of the record in the index and remove the record address (AAD) from all inverted file lists.
4. Decrement each file list, which has been affected by 1.

Addition of Keys

A problem arises here pertaining to record length. After the addition of a key, it is possible that the track could overflow. If so, the record is deleted and a link is inserted to another track where this whole record is inserted.

The Addition of Keys algorithm:

1. Decode the Key n in the proper directory to the variable length inverted file.
2. Determine the proper sequence, insert AAD of the record of Key A, and add 1 to the list length.
3. Continue steps 1 and 2 for all keys to be added.
4. Transfer the record from AAD (auxiliary address), and

add new keys to record.

5. If the record doesn't cause overflow of the track after packing it, return it to AAD.
6. If the record causes overflow, insert a link to another track, and place the whole record there.

2.4.2.2 LOGICAL OPERATOR "OR"

The logical operator "OR" effects a union of Directory A and Directory B. A key that resides in either Directory A or Directory B will be included in Directory C, and if the key appears in both directories, it appears only once in Directory C.

$$C = A \text{ OR } B$$

The Logical Operator "OR" algorithm:

1. Set 3 index registers (one register each for A, B, and C) to 1.
2. Compare unit 1 of A to unit 1 of B (the keys). If the key of A is smaller than the key of B, move the key of A to Directory C and increment index A and index C. If the key of A is equal to the key of B, move the key of A to Directory C and increment index A and index C. If the key of A is larger than the key of B, move the key of B to Directory C and increment index A and index C.
3. Repeat step 2 until Directories A and B have both been exhausted.
4. Stop.

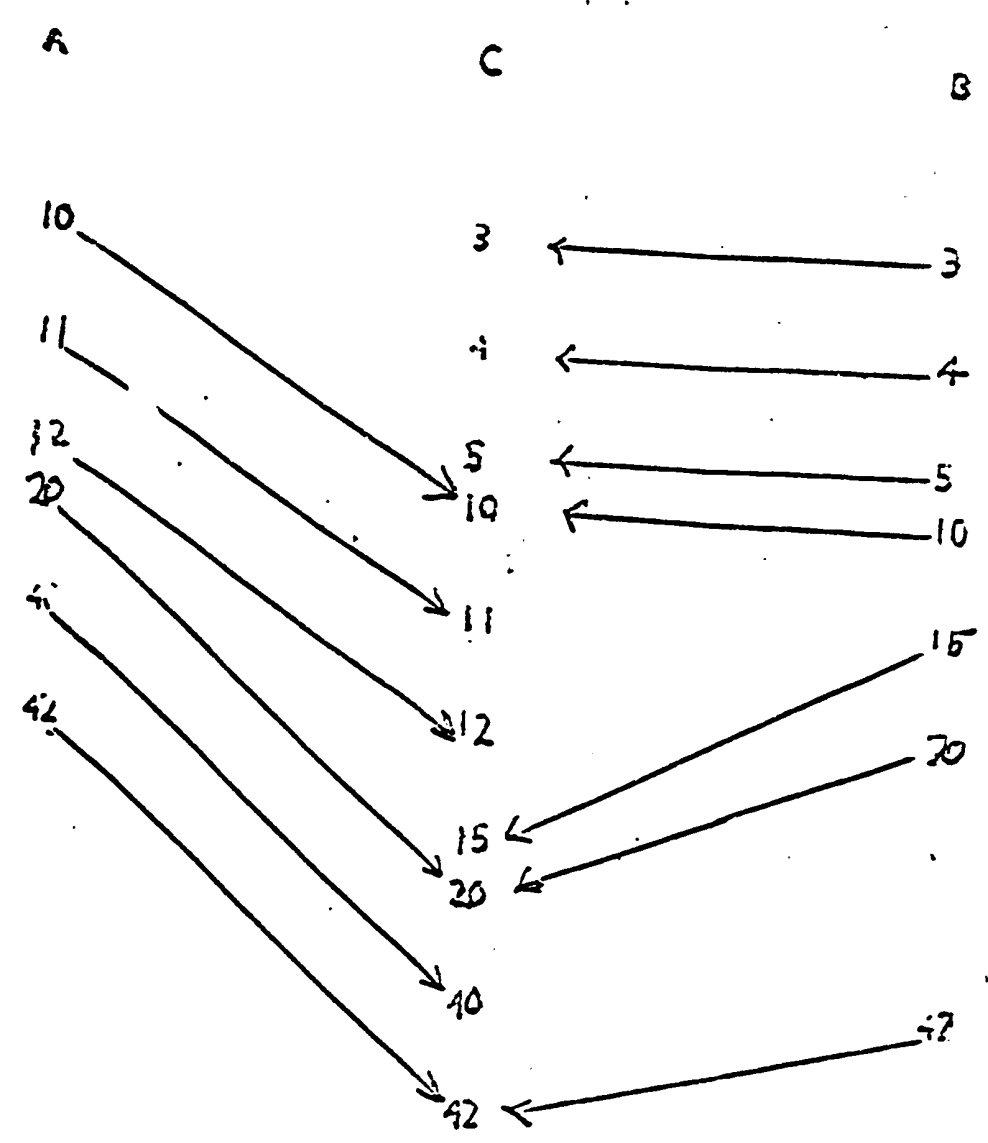


Figure IF-2

C = A "OR" B

Label	Operation	Op. and	Statement	Comments
*				
STAR	LH	9, AAA		
	LH	10, BBB		
	LH	11, CCC		
GO	CLC	0(2, 9), 0(10)	COMPARE FILE A TO FILE B	
	BC	4,ALO	A IS LOW	
	BC	8, EQAL	BRANCH TO EQAL	
	MVC	0(2, 11), 0(10)	B IS LOW, MOVE FILE B TO C	
	AR	10, H2	INCREMENT R10(B)	
	BC	15, INCC	INCREMENT FILE C	
*				
*				
ALO	MVC	0(2, 11), 0(9)	MOVE FILE A TO C	
	AH	9, H2	INCREMENT REG 9 - FILE A	
	CH	REND, 0(9)	IS FILE A FINISHED	
	BC	8, OUT	YES BRANCH OUT OF SUBROUTINE	
*				
*				
INCC	AH	11, H2	INCREMENT 11 - FILE C	
	BC	15, GO	BRANCH BACK TO GO.	

Label	Operation	Op. and	Statement	Comments
EQAL	MVC	0(2, 11), 0(10)	MOVE FILE B TO FILE C	
	AH	10, H2	INCREMENT REG 10 - FILE B	
	MVC	0(2, 11), 0(9)	MOVE FILE A TO FILE C	
	AH	9, H2	INCREMENT REGISTER 9	
	CH	9, REND	IS FILE	
	BC	8, OUT		
	BC	15, INCC	INCREMENT FILE C	
OUT	BCR	15, 14	BRANCH BACK TO MAIN PRG.	
	END	STAR		

Figure IF-3

2.4.2.2.2 A AND NOT B

Directory C will contain only the keys of Directory A that are not contained in Directory B.

The algorithm for C - A "AND NOT" B

1. Set 3 index registers (one each) for A, B and C to 1.
2. If the key of A is smaller than the key of B, move the key of A to Directory C, increment indexes A and C.
If the key of A is greater than the key of B, increment the index of B.
3. Repeat step 2 until Directories A and B have been exhausted.
4. Stop.

C = A "AND NOT" B

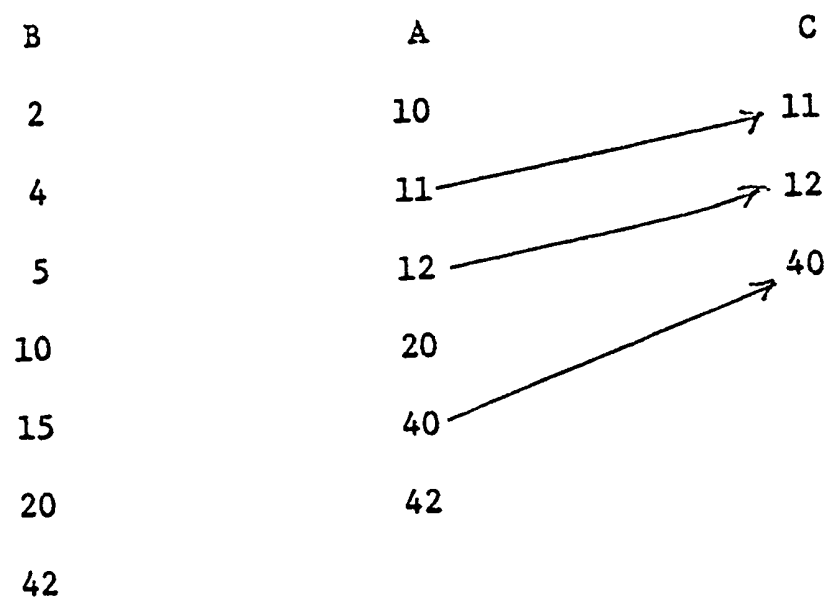


Figure IF-4

A "AND NOT" B

PROGRAM		FUNCTIONING INSTRUCTIONS		GRAPHIC		PAGE 1 OF 1	
PROGRAMMER	DATE			PUNCH		CARD SEQUENCE NUMBER	
STATEMENT							
21	22	23	24	25	26	27	28
Label	Operation	Address	Operand	Address	Operand	Comments	Identification Sequence
STAR	LH	9,	AAA			LOAD ADDRESS OF A	
	LH	10,	BBB			LOAD ADDRESS OF B	
	LH	11,	CCC			LOAD ADDRESS OF C	
GO	CLC	0(2,9),	0(10)			COMPARE FILE A & B	
	BC	4,	ALD			IF A IS LOW BRANCH	
	AH	9,	H2			EQUAL - SO GO TO NEXT A	
	CH	9,	REND			IS FILE A FINISHED	
	BC	8,	OUT			BRANCH TO OUT	
	BC	15,	GO			BRANCH BACK TO GO	
*							
*							
ALD	AH	10,	H2			INCREMENT FILE B	
	CH	10,	REND			IS FILE B FINISHED	
	BC	7,	GO			NO - BRANCH BACK TO GO	
	MVC	0(2,11),	0(9)			YES - SO MOVE A TO C	
*							
*							
	AH	9,	H2			INCREMENT A	
	AH	11,	H2			INCREMENT C	
	BC	15,	GO			BRANCH BACK TO GO	
OUT	BCR	15,	14			RETURN TO MAIN PROGRAM	

Figure IF-5

2.4.2.2.3 LOGICAL OPERATOR "AND"

Let us examine the operations it takes the system to create a resultant file, if we take two indexed files and request a retrieval of keys for which both files apply. Let's scan file A and file B.

It is written this way in set theory:

$$\text{INDEX C} = \text{A} \cup \text{B}$$

This allows us to retrieve the keys from Directory A that also reside in Directory B.

The algorithm for Logical Operator "AND"

1. Set 3 index registers (one each) for A, B, and C to 1.
2. Compare unit 1 of A to unit 1 of B.

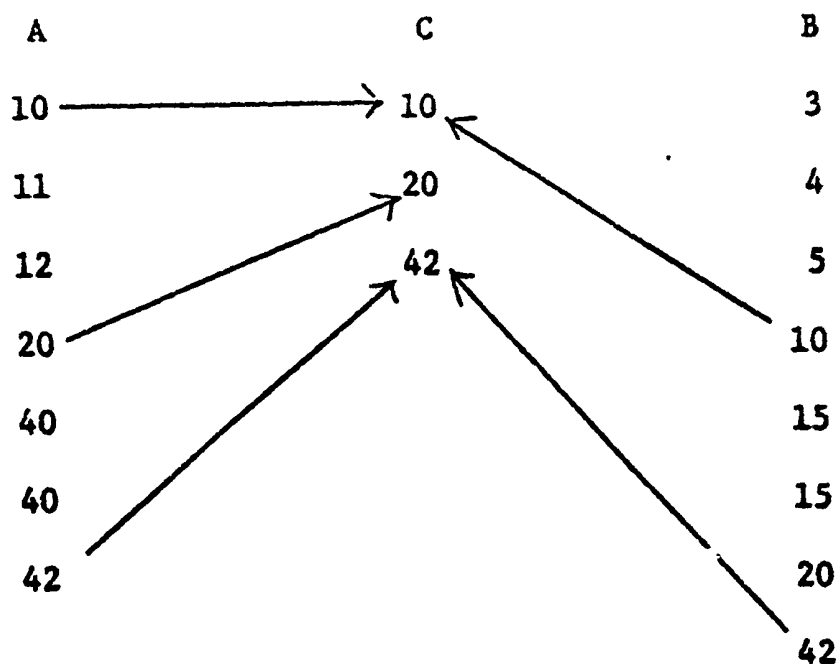
If the key of A is equal to the key of B, increment A to the next key of A, and increment B to the next key of B, move A to C and increment C.

If the key of A is smaller than the key of B, increment A to the next key of A.

If the key of A is larger than the key of B, increment B to the next key of B.

3. Repeat step 2 until all the keys in B have been investigated.

A "AND" B



Address	Operation	Comments
0	START	
1	USING	* , 0
2	STAR	LH 9, AAA
3		LH 10, BBB
4		LH 11, CCC
5	GO	CLC 0(2, 9), 0(10) COMPARE FILE A & FILE B.
6		BC 7, ALO BRANCH IF A IS LOW OR HI
7		MVC 0(2, 11), 0(9) A & B EQUAL MOVE A TO C
8		AH 9, H2 INC
9		BC 15, GO BRANCH BACK TO GO.
10		*
11		*
12	ALO	AH 10, H2 INCREMENT REGISTER 10 (FILE A)
13		CH 9, REND IS IT THE END OF A.
14		BC 8, OUT BRANCH TO MAIN PROGRAM.
15		CH 10, REND IS IT THE END OF 10.
16		BC 7, GO NO - NOT THE END
17		LH 10, BBB LOAD START ADDRESS OF FILE B
18		BC 15, GO BRANCH BACK TO GO
19		*
20		*
21	OUT	BCR 15, 14 BRANCH OUT OF SUBROUTINE
22		END START

Figure 1F-6

2.4.3 MULTILIST FILE

The use of "cells" or "pages" allows us to use fixed-size blocks to implement a method called the Multilist File. A sequential index contains the key values by which records are indexed. The pointer, which is associated with each key, points to the list of records that has the particular key value in question. The addressing method uses a page number and record number within the page.

A group of records can be retrieved at one time. This implies that access time can be reduced, however the saving is brought about only when a number of records are desired.

See figure M.F.-1, the Multilist File. The index, or directory contains keys L, M, N, and O. If we had a key that started with the letter L, the directory would decode it and take us to (3,3)/11, which translates to page 3, record 3 and tells us that there are 11 records in the list.

The other two methods discussed here use the same records and the same scheme for presenting them. There is a key, an address, and a list length sequence for each system.

The other two methods discussed here are presented with the same type drawing for ease of understanding.

The same logical operations performed on the inverted files can be performed on the multilist files.

Multilist

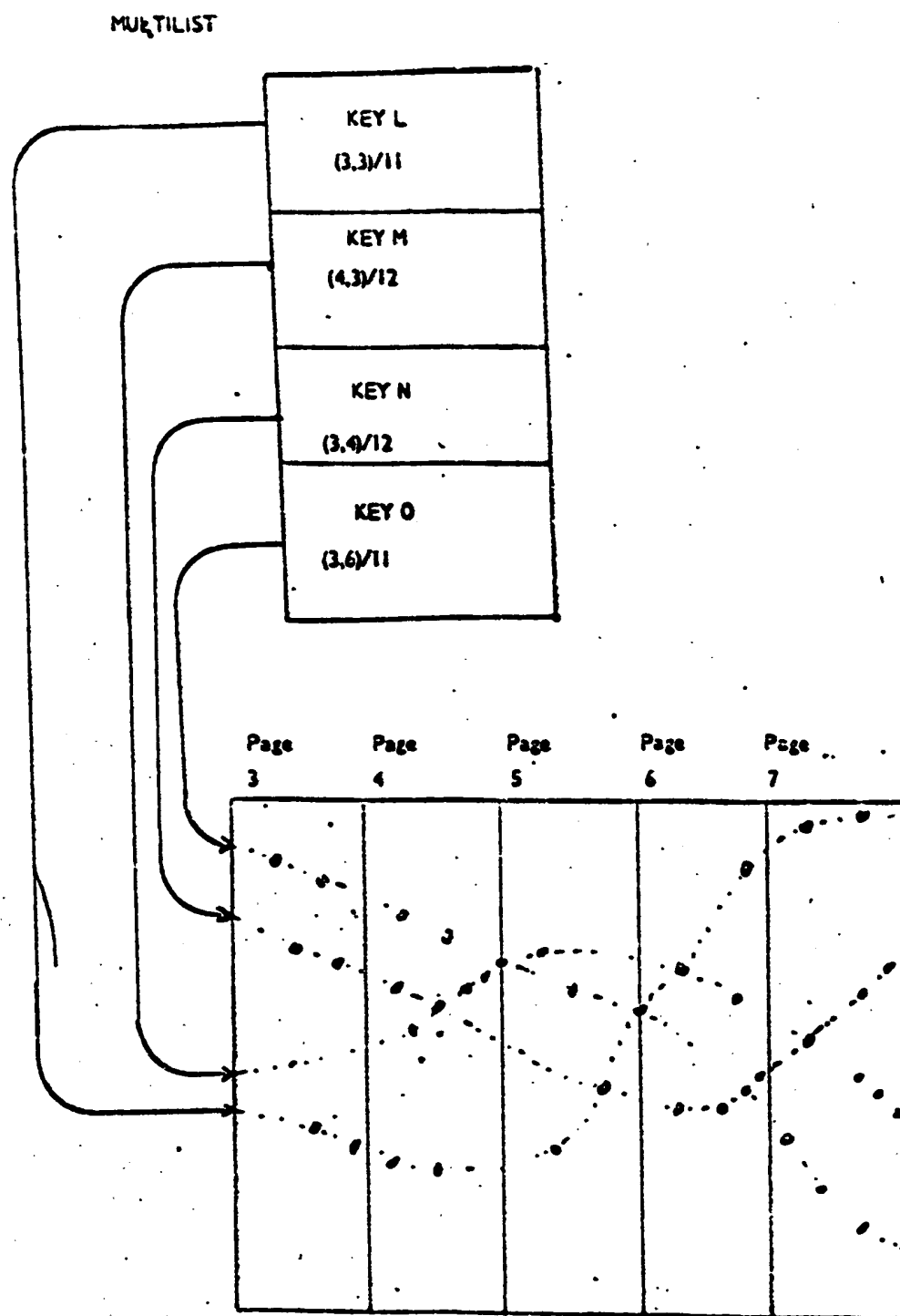
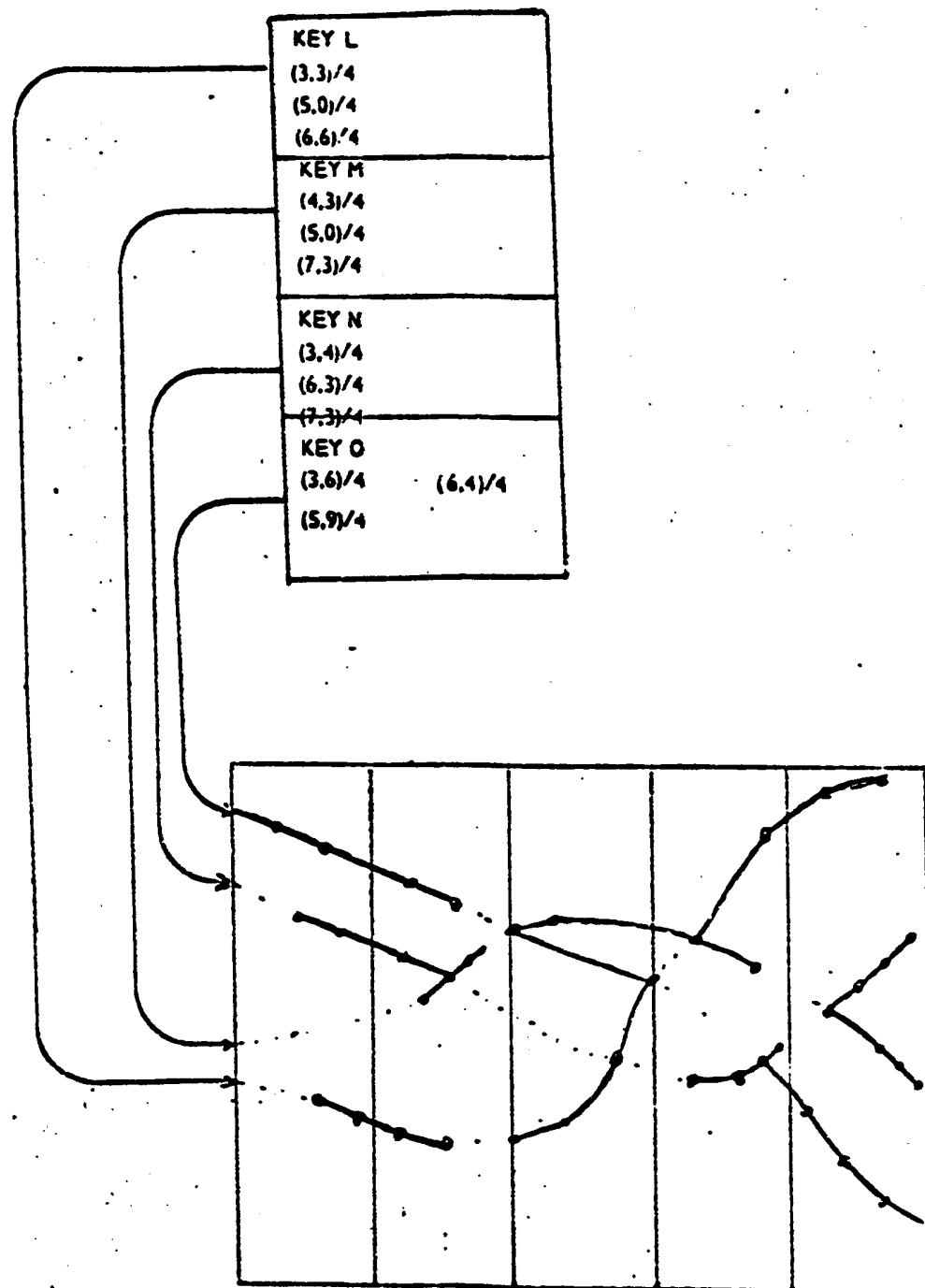


Figure M-1

Controlled Multilist Lengths

The same pattern is followed here with the exception that the list lengths are not more than four records long. The number in the lists is arbitrary. This avoids the long lists that the Multilist method had.

See Figure M.F.2, the Controlled Multilist File. The index contains the same keys that the Multilist File did, but we notice that there are more divisions within the key group. This time L has three divisions and the same number of records.



Multilist Organization for Pages

The Page Multilist files are linked so that they do not cross page boundaries. The advantage to this is that it does not require several pages of information when a key is used to retrieve certain information. A Multilist File might require several pages to be accessed.

The cells or pages become the partition instrument. This allows the list structure and the partition (page partition) to function.

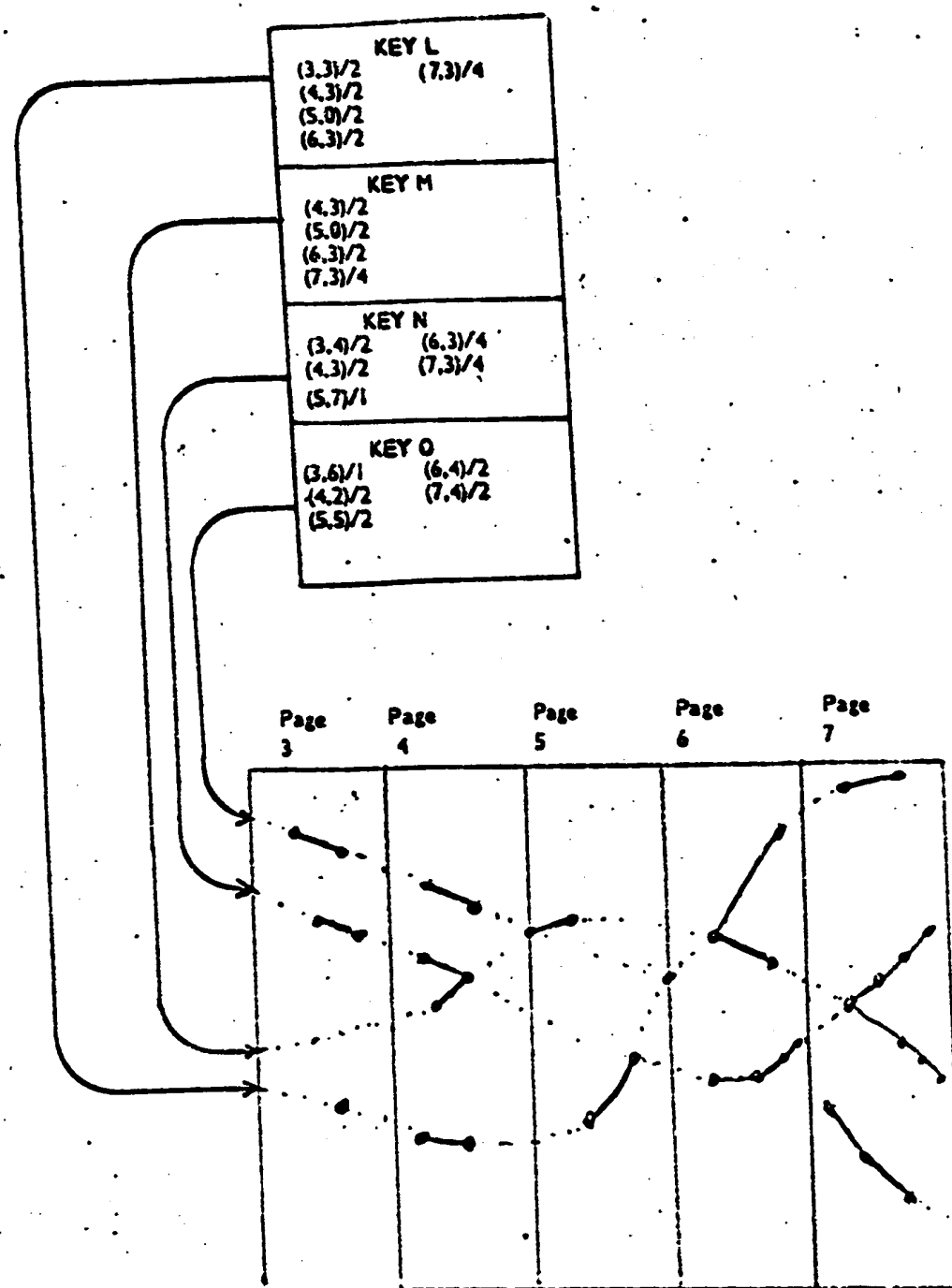


Figure M-3

2.5 SORTING

The object of the sort section is to present various methods of sorting to allow the reader a wider choice of methods from which to select.

The activity section at the end of this section suggests examples that are contained in the various Basic Assembler Language programs in the book.

SORTING METHODS

RECORD SORT

The sort algorithm is applied directly to the complete record by using the key in the record. This means the complete record is moved when a change is indicated. The end result is that each record is in its proper place when the file is sorted.

2.5.1.1 KEY SORT

The key sort works with the keys only, as opposed to the whole record. This includes an associative table of address pointers. Associative in this instance refers to a memory device, in which each cell contains information pertaining to the key and entry. The entry part tells us whether a cell has an entry pertaining to the key that accessed the particular cell. If there is no entry, this signal is returned as "not found" or something similar.

There are two types of key sorts, the detached and the non-detached key sort.

2.5.1.2 DETACHED KEY SORT

The algorithm for the Detached Key Sort:

1. The record keys are placed in an associative table with address pointers.
2. Apply the internal sorting algorithm to the associative table only.
3. If a transfer is to be made, make it include only the key and associated record address.
4. When the associative table is sorted, the complete records in the file are moved to output according to the sorted order of the address pointers.

This algorithm is quite forward, if we separate the key and address from the rest of the record to form a new table.

2.5.1.3 THE NONDETACHED KEY SORT

The nondetached key sort forms a table which contains pointers to the keys of the associated records. This method applies the sort algorithm indirectly through the associated address pointers in the table. The address pointers (not the keys) are moved during the sort. Both the table of addresses and the stored file are referenced during the sorting procedure.

When the sorting algorithm is finished, only the record address pointers have been sorted. The retrieval process requires that the sequence of the record address associative table be followed.

2.5.2 RADIX SORT

This method is used for sorting punched cards mechanically, but can be used for the computer as well.

The least significant digit of each key is examined. Then the

record is assigned to a pocket which depends on the value of the digit. After all the records have been examined, the records are distributed again, and so on, until all the digits of the key have been used for distribution. After the collection of the pass following the most significant digit, the records are in order. For decimal keys, ten pockets are needed, or else each pass could be replaced by two passes, in which case fewer pockets would be required.

The total number of passes is equal to the number of digits in the key, and the capacity of each pocket must be sufficient for all the records that might end up there.

Radix Sort

*	*	79	38	72	56	66	13	57	13	40	23	91	24	45	28	03	21
---	---	----	----	----	----	----	----	----	----	----	----	----	----	----	----	----	----

79	28	57	66	45	94	23	72	91	40
	38		56		24	03			
						23			
						13			
9	8	7	6	5	4	3	2	1	0

79	38	28	57	56	66	45	24	94	13	23	03	23	72	91	40
----	----	----	----	----	----	----	----	----	----	----	----	----	----	----	----

91		72	66	56	40	38	23	13	03
94		79		57	45		23		
							24		
							28		
9	8	7	6	5	4	3	2	1	0

94	91	89	72	66	57	56	45	40	38	28	24	23	23	13	03
----	----	----	----	----	----	----	----	----	----	----	----	----	----	----	----

Figure RS-1

2.5.3 QUICKSORT

The quicksort algorithm can be classified as a partitioning type algorithm, and it was developed by Hoare. The main aspects of the algorithm are:

1. The first pass located the item that occupies the middle spot of the list and classifies it as a bound.
2. This item is copied into a temporary location and replaced with the bottom of the list.
3. The list is scanned from the top, and each item is compared in sequence to the bound until an item is found that is larger than, or equal to, the bound.
4. If the new item is larger than the bound item, the new item is moved to the bottom of the list, creating a vacant spot in the top part of the list.
5. The top-down scan is then continued, alternating the scan each time a transfer is made, until all the items have been scanned. When the two top pointers coincide, the list will be partitioned into two parts.
6. The count is then placed in the vacant spot between the two partitions. This is the spot where the two pointers coincide. It now occupies its sorted position, because all items above it are smaller, and all items below it are larger or equal to it.

The partition that has the most elements is then stacked for later processing, while the algorithm proceeds to partition the smaller list in order. This continues until each item has been placed into its

proper sorted position either by being selected as a bound or by remaining a single item in a partition.

2.5.4 TOPOLOGICAL SORT

The topological sort is used in situations that require partial ordering. Pert and Critical Path methods are good examples of this process. Partial ordering occurs in mathematics in situations such as $A = B$ (between real numbers), and also between A and B as A \leq B in set theory.

The critical path program presented in a previous chapter has many good examples of the topological sort. Presented here is the algorithm (in part).

5. Sort the table on the basis of lowest starting event number.
6. Locate activity of the lowest starting event number.
7. Locate and place in table T10 all activities with end event numbers that match the start events of (6).
8. Sort T10 on the basis of ending times.
9. Subtract ending time of each activity from the ending time of the longest activity in table T10.
10. Enter difference in slack time position.
11. Enter longest event time from table T10 into start time for activity (6).
12. Compute end time for activity (6).
13. Index activity (6).
14. Transfer table T10 and activity (6) back to the main table.

15. Test for last event.

16. NO---back to (6).

2.5.5 QUADRATIC SORT

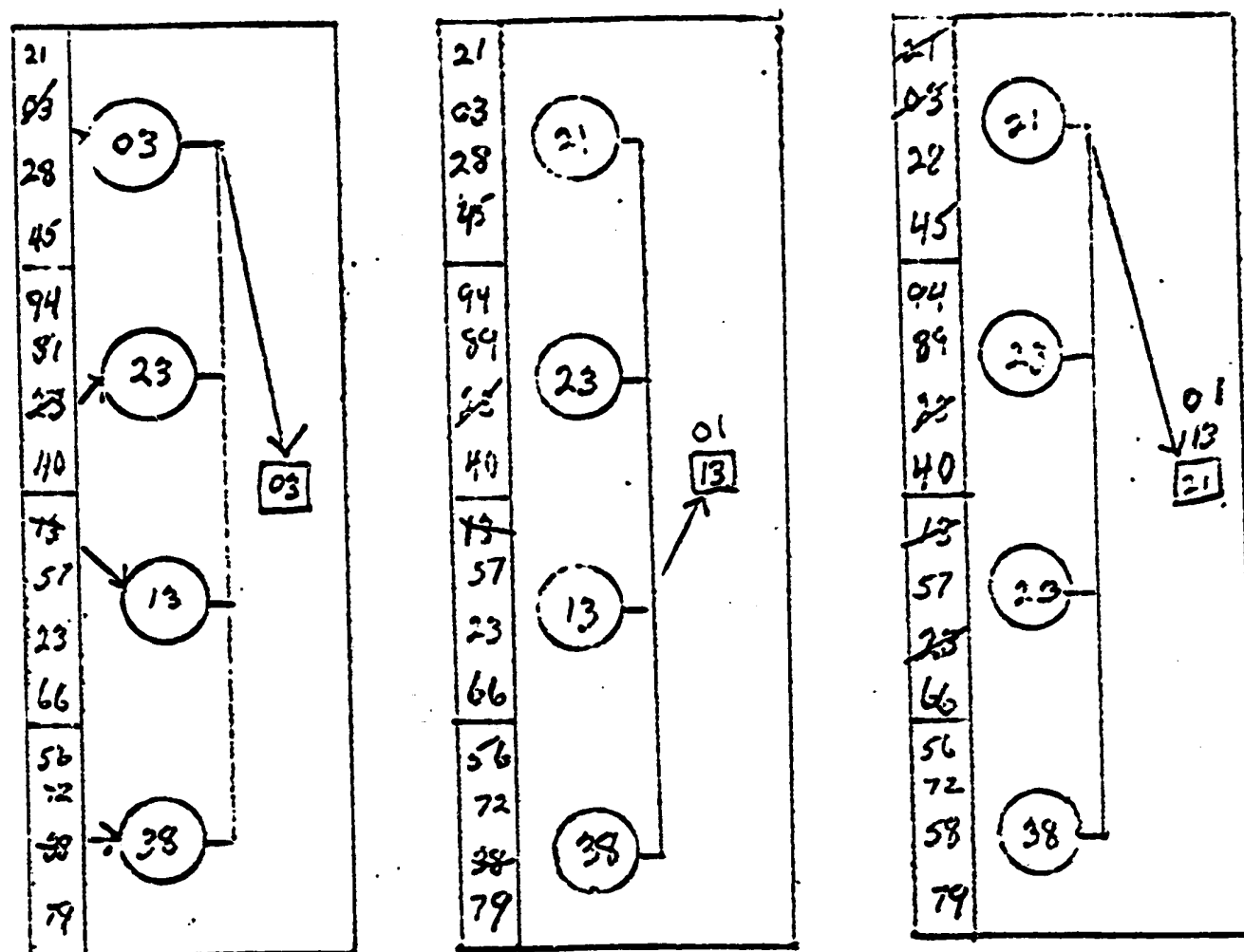


Figure QS-1

The quadratic sort divides the array into four equal parts and then searches each of the four groups for the smallest data. This number is erased and moved to a temporary area, where each of four are compared. The smallest of these is transferred to the output array. The process, with the exception of the division into four units, is combined until all the original array is blank, and all the intermediate

2.5.7 SORTING BY INSERTION

14		14	02	02	02	02	02	02
21		21	14	10	10	10	3	03
10			21	14	14	14	10	10
18				21	18	16	14	14
16					21	18	16	16
3						21	18	18
80							21	21
60							80	60
70								80
79								

Figure SI-1

Each key is examined in turn and inserted into the correct place. The earlier members are pushed down when the need arises. The mechanics of moving the earlier records down is also used as a method of insertion with lists. The actual number is $n/4$. This method of sorting is useful when there isn't much core storage.

Sorting by Insertion

Line	Op	Op	Op	Op	Op	Op	Op	Op	Op
TAB	CP	SCOR, TAB	2, 17	(5, 8)	COMPARE THE SCORE				
	BC	10, DWN			OF 2 HI OR EQUAL - BRANCH TO DWN				
	AH	8, STEP			INCREMENT REGISTER 8				
	BC	15, TAB			UNCONDITIONAL BRANCH BACK TO TAB				
*					MOVE THE TABLE DOWN				
DWN	STN	8, MARK			STORE THE ADDRESS IN REG 8 IF MARK				
	LH	9, TLM			LOAD THE TABLE END IN REG 9				
	MVC	TBL 2(20, 9), TBL 1-20(9)			MOVE LAST ITEM IN				
*					TABLE DOWN ONE (IN A360				
**					TRAINING THIS IS REALLY TO A				
*					HIGHER ADDRESS				
	SH	9, STEP			DECREMENT REGISTER 9				
	CH	9, MARK			DOES 9 EQUAL THE ADDRESS OF MARK				
	BC	3, DWN+2			IF 9 IS STILL HIGHER BRANCH TO				
*					THE MVC STATEMENT BELOW DWN				
X					THIS IS DWN+2.				
	MVC	TBL 2(20, 8), INP			PUT IT IN TABLE				
	BC	15, READ			READ ANOTHER CARD				

Figure S1-2

2.5.8 INTERCHANGE SORT

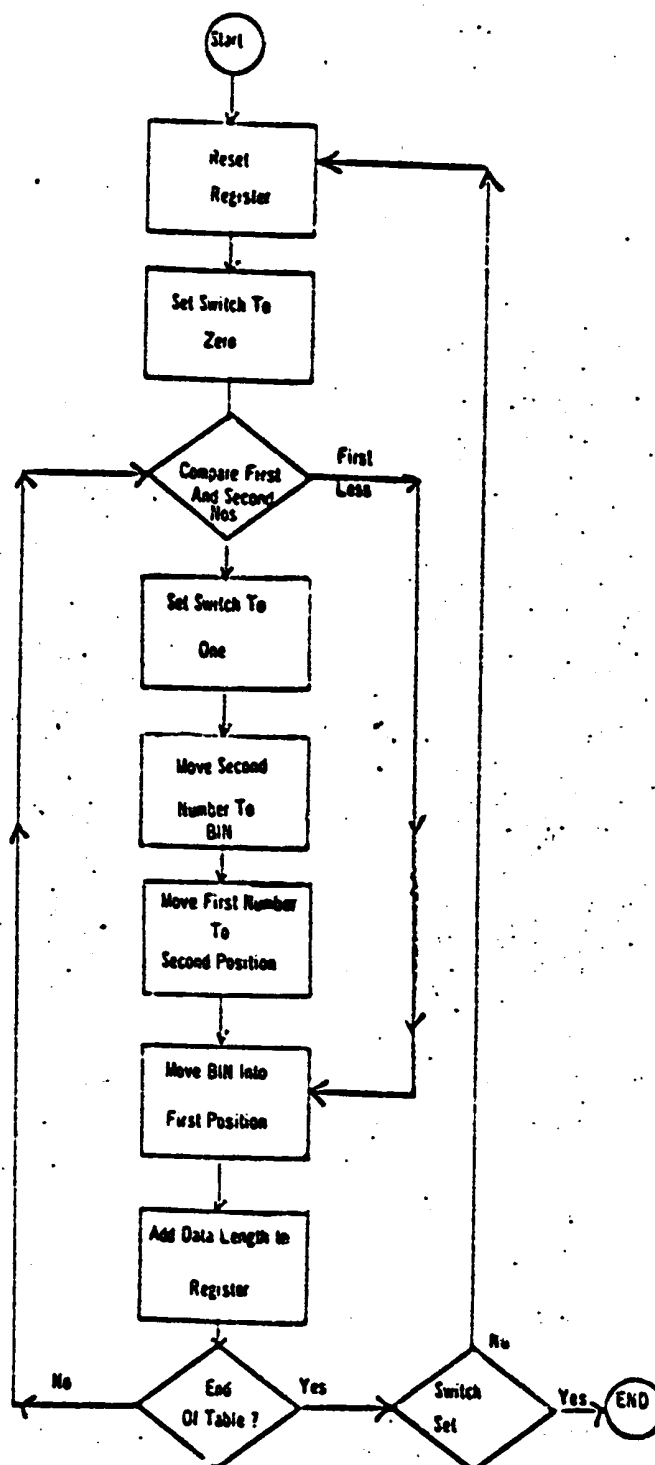


Figure IS-1

Interchange Sort

START	CI	SWA, X'02'		SET SWITCH SWA
	SR	8, 8		CLEAR REGISTER 8
	LH	9, TREG		LOAD START ADDRESS OF TABLE
	NI	SW, X'00'		TURN SW OFF
LOOP	CP	16(2, 8), 46(2, 8)		COMPARE FIRST TWO NUMBERS
	BC	13, 20H4		BRANCH IF EQUAL CALLW-1STOP
	OZ	SW, X'01'		SET SWITCH SW
	MVC	BIN(30), 0(8)	SDRT	FIRST NR TO BIN - 30 BYTES
	MVC	0(30, 8), 30(C)	SDRT	SECOND NR TO REPLACE FIRST.
	MVC	30(30, 8), BIN	MOVE	BIN TO REPLACE SECOND NR.
COMP	AH	8, THTY		INCREMENT REG R BY 30
	CLC	30(2, 8), STAR		IS IT THE END OF THE TABLE
	BC	7, LOOP		BRANCH IF UNEQUAL
	TM	SW, X'02'		TEST IF SW HIGH BRANCH TO SORT
	BC	2, SORT		OTHERWISE PROCEED

Address	Operation	Operand	Comments
	LH	8, TBEG	
	SR	9, 9	
	LH	9, TWO	
	ZAP	24(2, 8), 20(2, 8)	
	MVC	22(2, 8), STAR	
END	CLC	25(2, 8), STAR	
	BC	7, STOP	
	AH	8, THTY	
	CLC	0(2, 8), STAR	
	BC	8, STEP	
	BC	15, FIND	
STOP	MVC	BINB(24), 16(2)	
	SR	8, 8	
	LH	8, TBEG	
FINA	CP	18(2, 8), BINB(2)	
	BC	7, FINB	
	MVC	BINC(14), 16(8)	
	MVC	0(14, 9), BINC	
	AH	9, FTEN	
FINB	AH	8, THTY	
	CLC	0(2, 8), STAR	
	BC	7, FINA	
	MVC	0, (2, 9), STAR	

Figure M-2

2.6 MERGING

2.6.1 TWO WAY MERGE

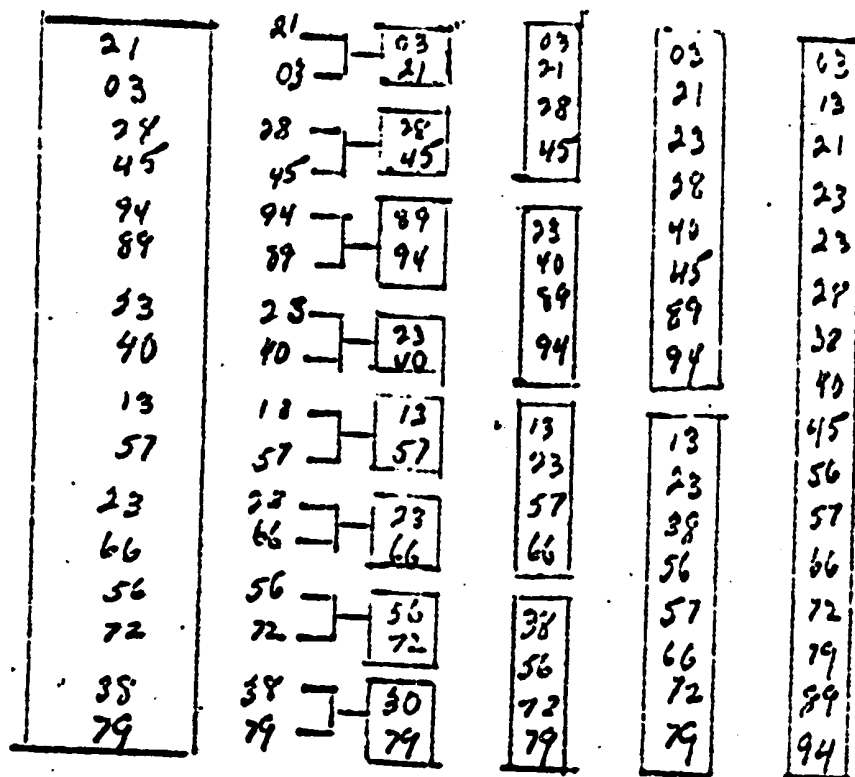


Figure M-1

Many different types of sorting are descendants of this basic approach to the merge. The two way merge compares pairs of keys, and each pair, of which the smaller key is placed first. After one pass, the initial group N consists of $N/2$ strings of length two. These pairs of strings are combined to form $N/4$ strings of length four. These pairs are combined to form $N/8$ strings, until one string results. If N is of the form 2^p , it will require p passes to complete the sort.

Merging Two Sorted Files

The comparison is made between the keys in the records. A merge can be made with the order ascending or descending. For our example, we will use the ascending merge, since our two input files are sorted in ascending order.

This merge effects the union of File A and File B into File C. If a number appears in both files, the number in File A will receive first precedence.

The Merge Algorithm

1. Set two indexes (1 each) for A and B, also set an index for C. (Set these indexes to 1.)
2. Compare the key of File A to the key of File B, if the key of A is smaller than the key of B, move the key of A to File C and increment index A and index C.
 - a. If the key of A is equal to the key of B, move the key of A to File C.
 - b. If the key of File A is larger than the key of File B, move the key of B into File C, and increment index B and index C.
3. Repeat step 2 until all of Files A and B have both been exhausted.

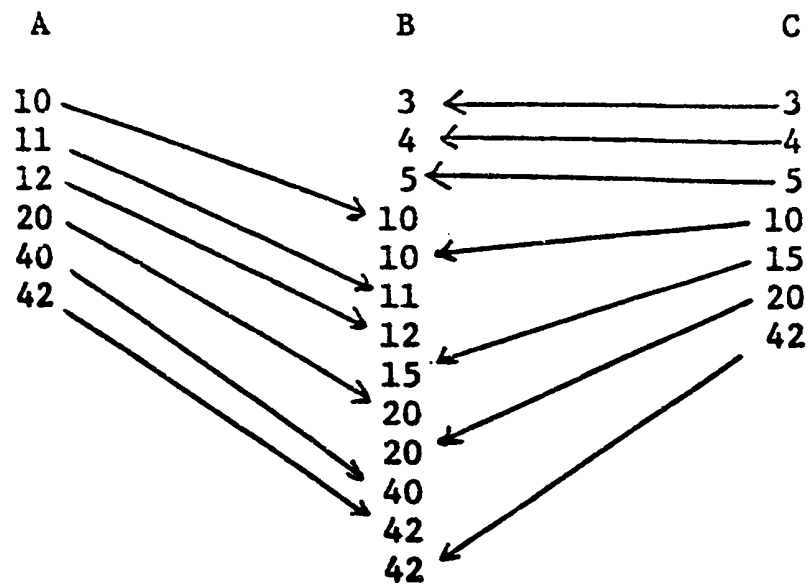


Figure M-3

2.6.1.2 MERGING ORDERED FILES WITH SUBFILES

There are four possibilities that can exist here:

1. There can be no problem, such as the preceding example, where no changes are necessary.
2. A step down condition can exist (a single step down).
3. A double step down condition can exist.
4. It may be necessary to perform a roll out.

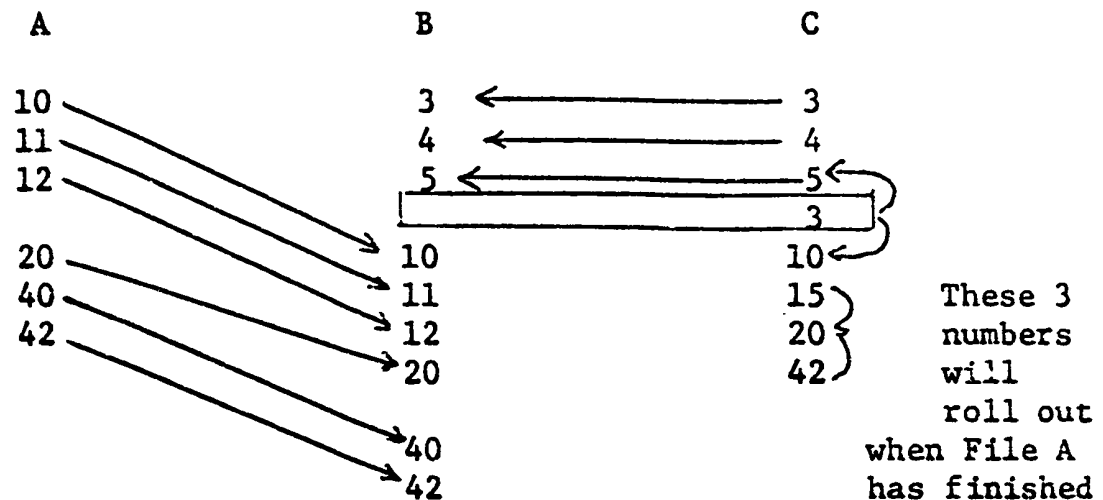


Figure M-5

The single step down is given in the example above, M 1. The number 3 in File B is a step down, because it is smaller than the number preceding it in the file, and it is also smaller than the number that follows it in File B. This means that File B will not be used again until File A has a step down also. The sequential process continues with File A until A results in the step down mentioned above. If there are no further step downs in File A, then File B will continue sequentially. This is called a roll out.

Merging Ordered Files with Subfiles

Line	Op	Opnd	Statement
STAR	LH	9, AAA	LOAD ADDRESS OF FILE A
	LH	10, 3BB	LOAD ADDRESS OF FILE B
	LH	11, 2CC	LOAD ADDRESS OF FILE C
* GO3	MVC	TEMP, 0(9)	MOVE 9-2 TO COMPARE
	AH	9, H2	
	CH	9, TEMP	COMPARE 9 TO 9-2
	BC	4, R0L2	BRANCH TO SET SWITCH
	SH	9, H2	RETURN 9 TO ORIGINAL NR
* * GO4	MVC	TEMP, 0(10)	MOVE 10-2 COMPARE
	AH	10, H2	REDUCE 10
	CH	10, TEMP	COMPARE 10 TO 10-2
	BC	4, R0L2	BRANCH TO SET SWITCH 2
	SH	10, H2	

Line	Op	Opnd	Statement
GO	CLC	0(2, 9), 0(10)	COMPARE FILE A TO FILE B.
	BC	8, EQUAL	BRANCH IF EQUAL.
	BC	4, ALD	
	MVC	0(2, 11), 0(10)	B IS LOW MOVE B TO C.
	AH	10, H2	INCREMENT B.
	BC	15, INCC	BRANCH TO INCREMENT C.
* * ALD	MVC	0(2, 11), 0(9)	MOVE FILE A TO C.
	AH	9, H2	INCREMENT FILE A.
INCC	AH	11, H2	INCREMENT FILE C.
	BC	15, GO3	
* * EQUAL	MVC	0(2, 11), 0(10)	MOVE FILE B TO FILE C.
	AH	10, H2	INCREMENT FILE B.
	MVC	0(2, 11), 0(9)	MOVE FILE A TO FILE C.
	AH	9, H2	INCREMENT FILE A.
	AH	11, H2	INCREMENT FILE C.
	CH	9, REND	IS FILE A FINISHED
	BC	8, CDMP	YES - BRANCH OUT.
	BC	15, GO3	GO TO BEGINNING
* * OUT	BC R	15, 14	BRANCH BACK TO MAIN PROG.

Figure M-6

Merging Ordered Files with Subfiles

Statement	Address	Operation	Comments
ROLL	MVC	0(2, 11), 0(9)	MOVE FILE A TO FILE C
	AH	11, A2	INCREMENT FILE C
	AH	9, A2	INCREMENT FILE A
	MVC	TEMP, 0(9)	MOVE FILE A TO TEMP
	AH	9, A2	INCREMENT FILE A
	CH	9, TEMP	COMPARE A TO AH2
	BC	4, ROLL	IF AH2 IS LOW, BRANCH TO ROLL2
	CH	9, REND	IS IT THE END OF FILE A?
	BC	4, ROLL	IF NO - LOOP TO ROLL
	MVC	SW, ONE	YES - SET SWITCH ONE
	BC	15, COMP	BRANCH TO COMP
ROLL2	MVC	0(2, 11), 0(10)	MOVE FILE B TO FILE C
	AH	11, A2	INCREMENT FILE C
	AH	10, A2	INCREMENT FILE B
	MVC	TEMP, 0(9)	MOVE AREA INDEXED BY 10 TO TEMP
	AH	9, A2	INCREMENT FILE B
	CH	9, TEMP	COMPARE TO B+2
	BC	4, ROLL	IF B+2 LOW - BRANCH TO ROLL
	CH	9, REND	END OF FILE B?
	BC	4, ROLL2	NO - BRANCH BACK TO ROLL2
	MVC	SW1, ONE	YES - SET SWITCH 1
	BC	15, COMP	BRANCH TO COMPARE

Statement	Address	Operation	Comments
ROLL3	CLL	SW1, SW	
	BC	2, ROLL2	
	BC	4, ROLL	
	BC	8, OUT	
	END		

Figure M-6 (Part-2)

Double Step-Down

The double step-down is treated as if there were no step down at all. In fact it is handled like the straight merge.

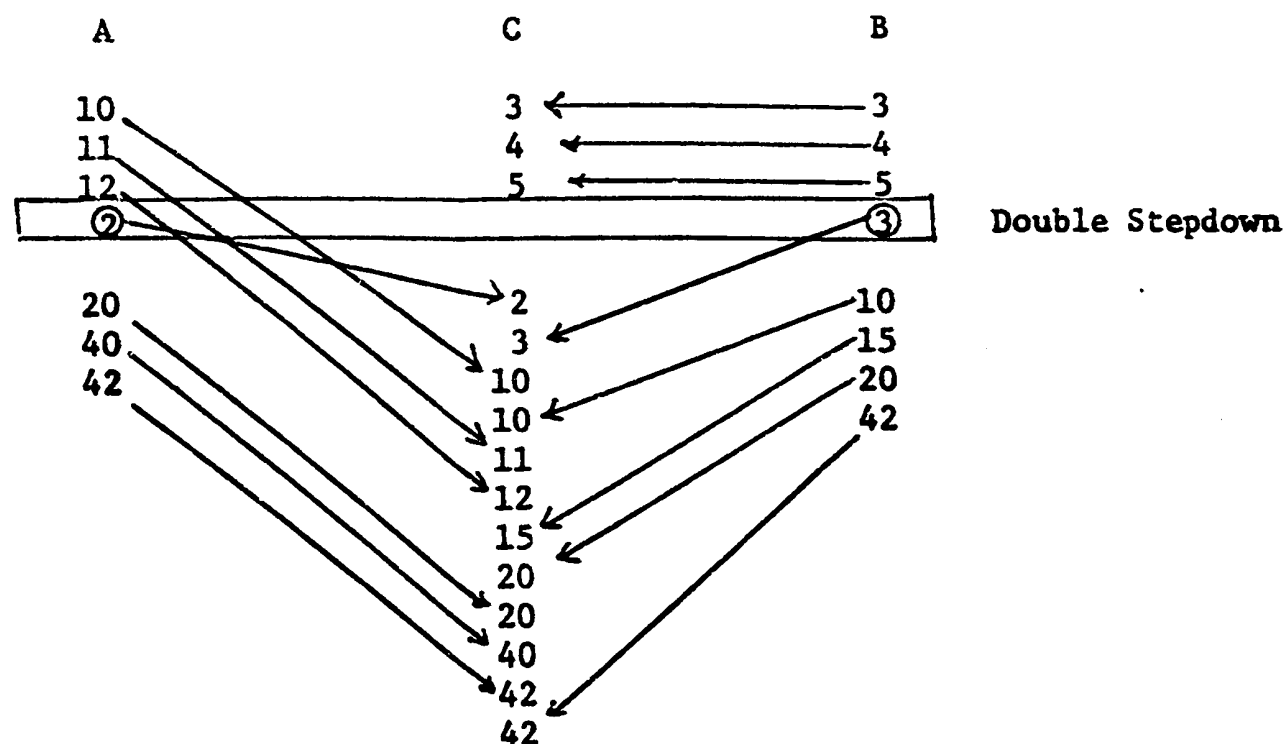


Figure M-7

The double step-down is a step down which occurs in both files, so it operates in the same manner as a merge.

Activities

1. What type of sort does the construction program use? If the reader will notice, the sort flow chart was taken out of context for emphasis.
2. What type of sort does the STAT program use? Compare it to the example presented in this chapter.
(Answer) A displacement insertion type.
3. What type of sort does cause problems when there are many

digits to sort?

(Answer) Interchange Sort.

Merging--Double Step-Down

Line	Op	Address	Comment
STAR	LH	9, AAA	LOAD START OF FILE A
	LH	10, BBB	LOAD START OF FILE B
	LH	11, CCC	LOAD START OF FILE C
GO3	MVC	TEMP, 0(10)	MOVE A TO TEMP
	AH	9, H2	INCREMENT A
	CH	9, TEMP	COMPARE A TO A+2
	BC	4, SWCH	A+2 IS LOW - BRANCH TO SWCH
	SH	9, H2	SET A BACK 2 BYTES
GO4	MVC	TEMP1, 0(10)	MOVE B TO TEMP
	AH	10, H2	INCREMENT B
	CH	10, TEMP1	COMPARE B TO A+2
	BC	4, SWLO	STEP DOWN IF B+2 LOW
	SH	10, H2	SET FILE B BACK 2 BYTES
	BC	15, COMB	BRANCH TO COMB

Line	Op	Address	Comment
GO	CLC	0(2, 9), 0(10)	COMPARE FILE A & FILE B
	BC	3, EQUAL	BRANCH IF EQUAL
	BC	4, ALO	
	MVC	0(2, 11), 0(10)	B IS LOW MOVE B TO C
	AH	10, H2	
	BC	15, INCC	INCREMENT FILE C
ALO	MVC	0(2, 11), 0(9)	MOVE FILE A TO FILE C
	AH	9, H2	INCREMENT FILE A
INCC	AH	11, H2	INCREMENT FILE C
	BC	15, GO3	
EQUAL	MVC	0(2, 11), 0(10)	MOVE FILE B TO FILE C
	AH	10, H2	INCREMENT FILE B
	MVC	0(2, 11), 0(9)	MOVE FILE A TO FILE C
	AH	9, H2	
	CH	9, PEND	
	BC	7, ALL	BRANCH TO SET SWITCH
	CH	10, PEND	
	BC	8, ALL	BRANCH TO SET SWITCH

Figure M-8

Merging--Double Step-Down

Label	Operation	Address	Comments
ROLL	MVC	0(2, 11), 0(9)	MOVE FILE A TO FILE C
	AH	11, #2	INCREMENT FILE C
	AH	9, #2	INCREMENT FILE A
	MVC	TEMP, 0(9)	MOVE FILE A TO TEMP
	AH	9, #2	INCREMENT FILE A
	CH	9, TEMP	COMPARE A TO A
	BC	4, ROLL	IF A#2 IS LOW, BRANCH TO ROLL
	CH	9, REND	IS IT THE END OF FILE A?
	BC	4, ROLL	IF NO - LOOP TO ROLL
ALL	MVC	SW, ONE	YES - SET SWITCH ONE
	BC	15, COMP	BRANCH TO COMP.
ROLL2	MVC	0(2, 11), 0(10)	MOVE FILE B TO FILE C
	AH	11, #2	INCREMENT FILE C
	AH	10, #2	INCREMENT FILE B
	MVC	TEMP2, 0(10)	MOVE AREA IN FILE B TO TEMP2
	AH	10, #2	INCREMENT FILE B
	CH	10, TEMP	COMPARE ? - ?
	BC	4, ROLL	IF B#2 LOW - BRANCH TO ROLL
	CH	10, REND	END OF FILE B
	BC	4, ROLL2	NO - BRANCH BACK TO LOOP ROLL2
ALL2	MVC	SW2, ONE	YES - SET SWITCH 2
	BC	15, COMP	BRANCH TO COMPARE

Label	Operation	Address	Comments
COMP2	CLC	SW3, SW	COMPARE END OF A & B
	BC	2, ROLL2	END OF A - BRANCH TO ROLL2
	BC	4, ROLL	END OF B - BRANCH TO ROLL
	BC	8, OUT	END OF BOTH - SO OUT
SWCH	MVC	SW3, ONE	SET SWITCH 3
	BC	15, GO4	BRANCH TO GO4
SWC0	MVC	SW4, ONE	SET SWITCH 4
	BC	15, COMP	BRANCH TO COMPARE
GO1L	CLC	SW3, SW4	COMPARE SWITCH 3 & 4
	BC	8, GO7	BOTH EQUAL - JUMP TO GO7
	BC	4, ROLL	SW3 LOW - BRANCH TO ROLL
	BC	2, ROLL2	SWITCH 3 HIGH - JUMP TO ROLL2
GO7	MVC	SW3, ZERO	ZERO SWITCH
	MVC	SW4, ZERO	ZERO SWITCH
	BC	15, GO3	

Figure M-8 (Part-2)

2.7 SEARCH

Search structures are concerned with the retrieval of data from data structures. The key plays an important part, because the key is the part of the record which is searched. Frequently, large data structures must be processed and tabulated.

A search structure consists of a search process together with a data structure method. Generally, it takes work on both of these, if either one is affected or improved; therefore it is quite important for the programmer to be cognizant of the system as a whole, not just the search, but the search and the structure also.

The selection of a search structure does in fact determine the search and insert and delete methods for a particular data base.

The speeds of these methods are also an important determining factor for selecting a search structure.

The linear search, as presented with the sequential file previously, examines each key in sequence and is terminated with a matching key, or when the last item in a file matches, a termination is effected.

Generally the number of items tested depends on the number of items in the complete data base. The number of items actually tested during a linear search is estimated at $N/2$.

2.7.1 LINEAR SEARCH

A search is made of the array TH of N integers for one that is equal to Q.

2.7.2 BINARY SEARCH

The binary search provides a fast search process through the

ordering of the data base, however the insert and delete process is relatively slow. The continual ordering and reordering of the data base each time a new item is added or taken out, is relatively slow.

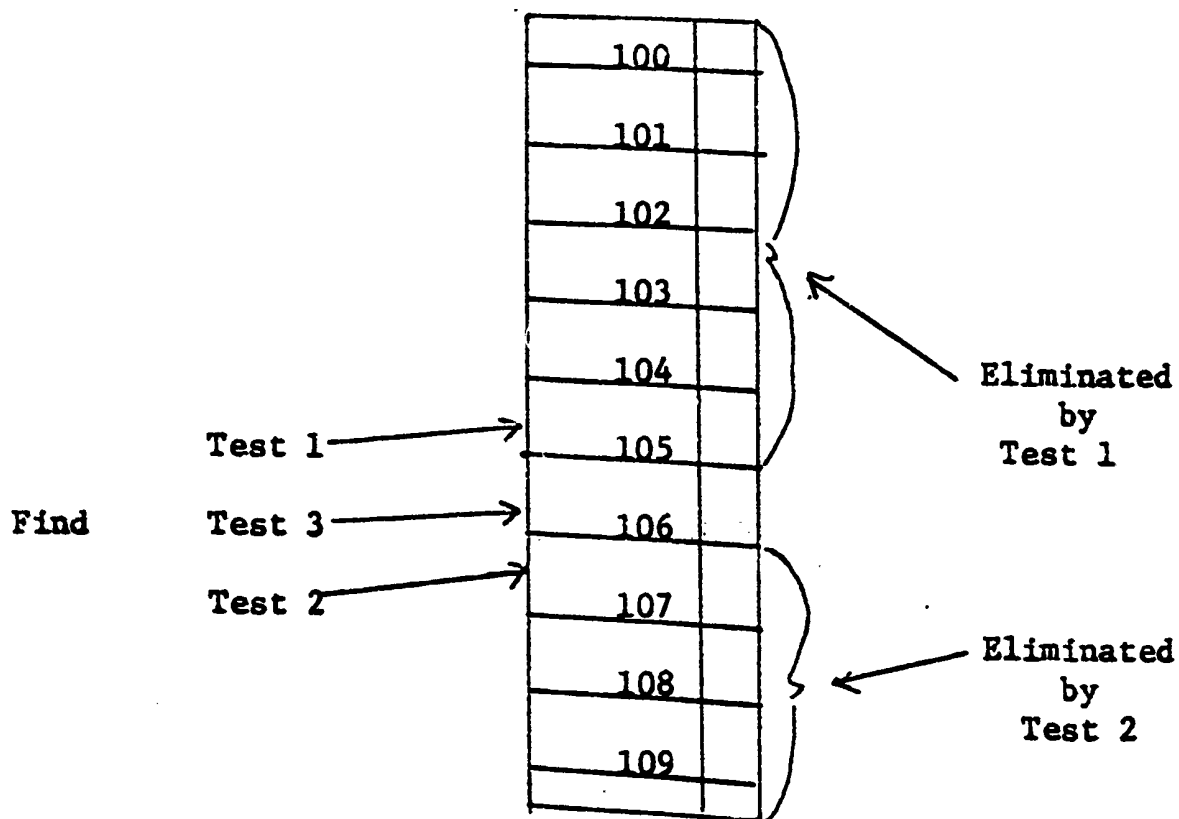


Figure BS-1

The search starts by entering the file or list near the middle. The word that divides the list in two is sometimes called a fence (a form of partitioning). By comparing a fence key with the key of the desired record, a decision is made whether to look further in the sublist above or below. If the sublist above is taken, the list is divided by 2, $(X/2)$, and another fence is created at this location. This process continues until either a list top, a list bottom, or matching record is determined.

Binary Search

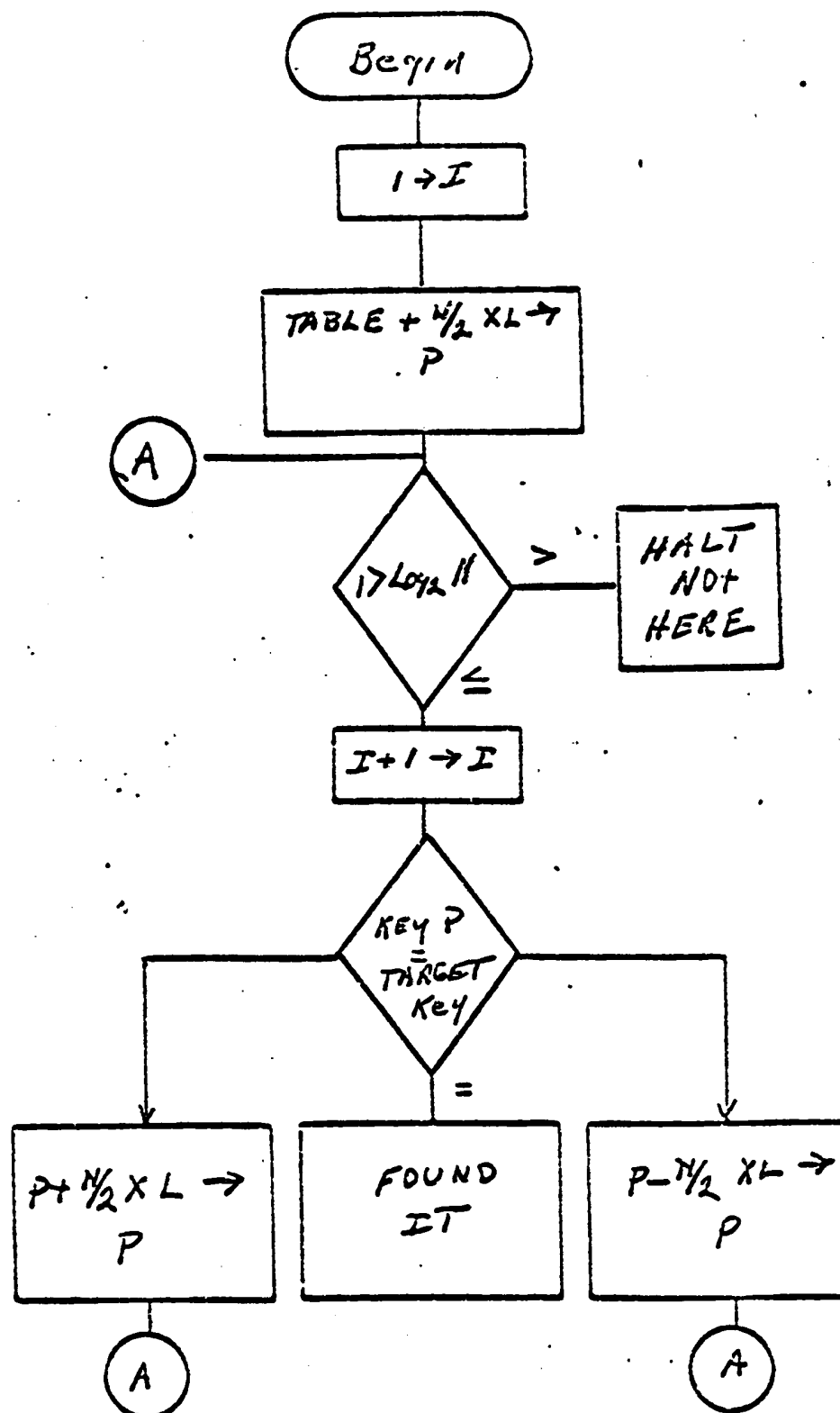


Table = First address of table

L = Length of entry in table

P = Pointer

I = Iteration

Figure BS-2

Table Search Program

FORTRAN STATEMENT

	DIMENSION TABLE (200)								
	READ 70 TABLE								
70	FORMAT								
	READ 80 AMONT								
80	FORMAT								
	J = 1								
	N = 1								
90	DO 100 I = N, 200								
	IF (AMONT - TABLE(I), 100, 200, 100								
100	CONTINUE								
	IF (J-1) 110, 110, 150								
110	PRINT 120, AMONT								
120	FORMAT								
150	STOP								
200	PRINT 230, J, I, AMONT								
230	FORMAT								
	J = J + 1								
	IF (I - 100) 250, 150, 150								
250	N = I + 2								
	GO TO 90								
	END								

Figure TS-1

2.8 DYNAMIC STORAGE ALLOCATION

The list processing languages have a common feature, which is, memory space for data structures does not have to be preassigned. The storage for each structure is allocated as it is needed and usually not sequentially. This is accomplished by the linking process.

In order to be able to reassign the use of memory cells during execution of a list processing program, a list processing language must provide for:

1. A storage of cells available for use.
2. Systems for obtaining "new" cells from, and returning unneeded cells to the store.

We will see that this also applies to virtual memory, which follows this section. Knuth mentions two methods for storage allocation but considers the First Fit Method the best.

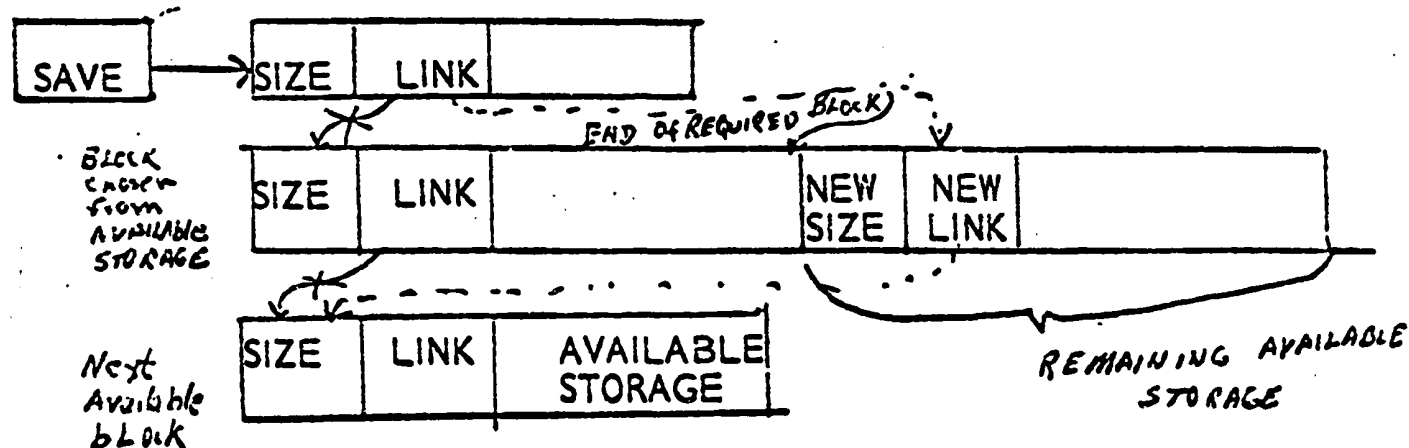


Figure DS-1

Dynamic Storage Allocation

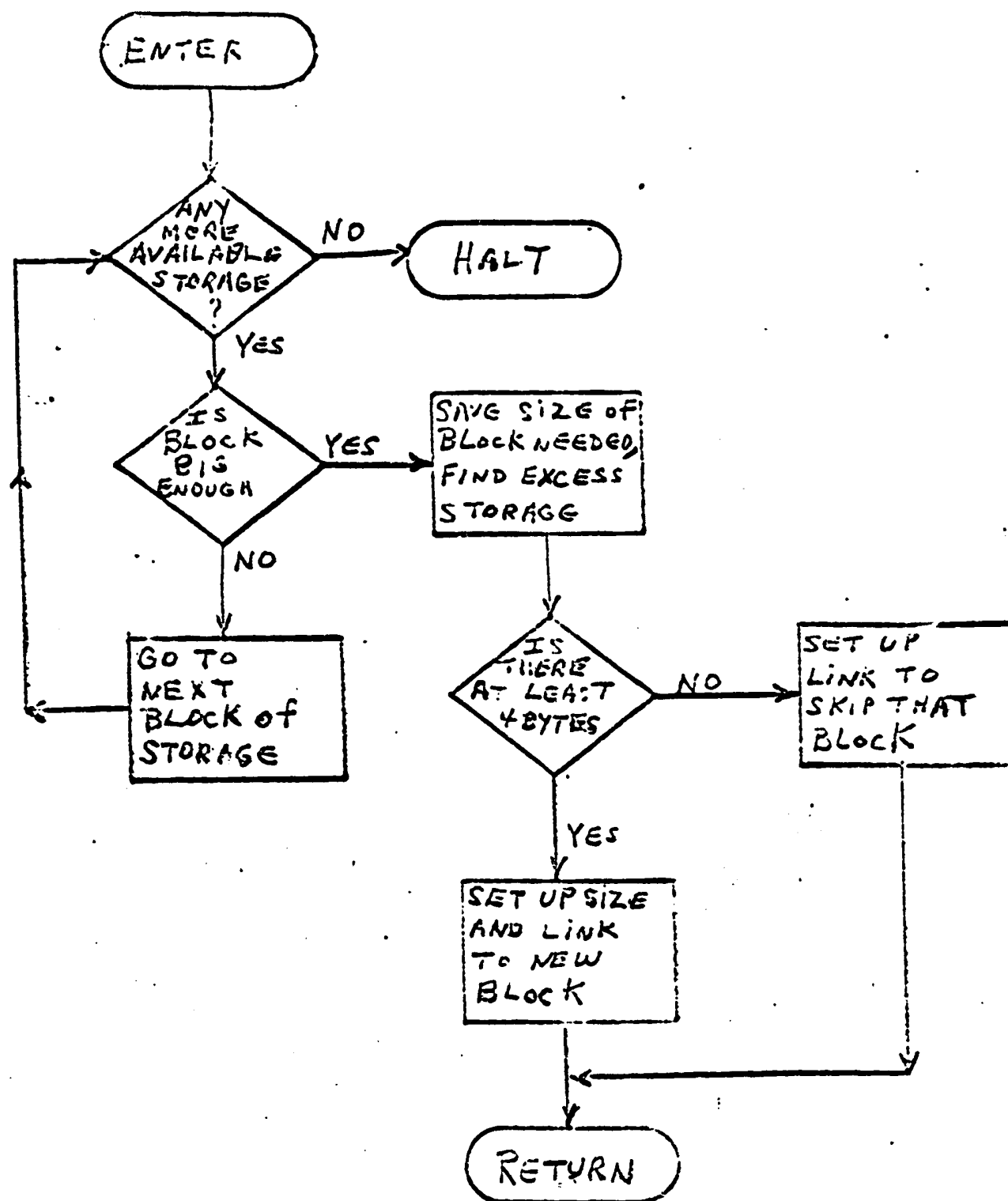


Figure DS-2

Dynamic Storage Allocation

PROGRAM		PROGRAM INSTRUCTIONS		STATEMENT		COMMON		SEQUENCE	
LINE NO.	OPERATION	OPERAND	OPERAND	STATEMENT	COMMON	SEQUENCE	SEQUENCE	SEQUENCE	SEQUENCE
	CHK	C/C	NULL, 0(L)	IS IT THE END OF AVAILABLE STOR.					
		B/C	8, NONE	YES, GO TO NONE AVAILABLE.					
		CH	9, 0(0, 8)	IS THE BLOCK BIG ENOUGH.					
		BC	13, OK	YES, GO TO NEXT PART OF PROC.					
		AH	8, H2	POINT TO LINK OF AVAILABLE STOR.					
		STH	8, SAVE	SAVE POINTER ADDRESS					
		LH	8, 0(0, 8)	POINT TO NEXT AVAIL. STORAGE					
		BC	15, CHK	CONTINUE TO CHECK STORAGE					
	*								
	OK	STA	9, SIZE	SAVE SIZE OF BLOCK REQUIRED					
		LH	9, 0(0, 8)	LOAD REG 9 WITH SIZE OF AVAIL. BLOCK					
		SH	9, SIZE	FIND AMOUNT OF EXCESS STOR IN BLOCK					
		CH	9, H4	IS THERE 4 BYTES LEFT?					
		BC	13, SKIP	NO, DON'T SAVE IT.					
	*								
	*								

PROGRAM		PROGRAM INSTRUCTIONS		STATEMENT		COMMON		SEQUENCE	
LINE NO.	OPERATION	OPERAND	OPERAND	STATEMENT	COMMON	SEQUENCE	SEQUENCE	SEQUENCE	SEQUENCE
	SR	10, 10		CLEAR REGISTER 10.					
	AR	10, 8		SAVE ADDRESS OF BLOCK IN RG 10.					
	AH	8, SIZE		FIND ADDRESS OF REMAINING BLOCK					
	STH	9, 0(0, 8)		STORE SIZE OF NEW BLOCK					
	LH	9, SAVE		POINT TO PREVIOUS AVAILABLE BLOCK					
	AH	9, H2		POINT TO LINK OF THAT BLOCK					
	STH	8, 0(0, 9)		STORE LINK TO NEW BLOCK					
	AH	10, H2		POINT TO LINK OF CHOSEN BLOCK					
	AH	8, H2		POINT TO LINK OF NEW BLOCK					
	MVC	0(2, 8), 0(10)		MOVE LINK INTO NEW BLOCK					
	BCR	15, 11		RETURN FROM SUBROUTINE					
	*								
	*								
	SKIP	SR	10, 10	CLEAR REGISTER 10.					
		AR	10, 8	SAVE ADDRESS OF BLOCK IN RG 10.					
		AH	8, H2	POINT TO LINK OF CHOSEN BLOCK					
		LH	9, SAVE	POINT TO PREVIOUS BLOCK.					
		AH	9, H2	POINT TO LINK OF THAT BLOCK					
		MVC	0(2, 9), 0(8)	LINK PREV BLOCK TO NEXT AVAIL. BLOC					
		BCR	15, 11	RETURN FROM SUBROUTINE					

Figure DS-2

2.9 VIRTUAL MEMORY

This discussion is based on a paper by Denning,¹ who has made a survey of virtual memory.

Earlier developments brought the state of the art to a point where computer costs would be too great to provide the large memories that were being suggested.

The proposal for the Atlas computer in 1961 set forth a one level store memory. It was known as Virtual Memory. The central idea is that "address" is entirely different from "physical location." The computer hardware and software automatically transfer information into memory at the precise time it is needed for processing. Also the hardware and software arrange for the program-generated addresses to be directed to memory locations that contain the information addressed. Virtual memory presented a great potential for overcoming part of the problems of storage allocation, because the memory use was based on system-observed actual use of space.

The mechanisms for affecting virtual memory now become quite important. The mechanisms referred to here are Segmentation and Paging.

Segmentation organizes address space (not memory) into variable size segments, while Paging organizes address space into fixed size pages of contiguous addresses.

The programmer is allowed to work with addresses which are different from memory, then the system provides the mechanism for translation of these program-generated addresses into the proper memory

¹Peter J. Denning, "Virtual Memory," *Computing Surveys*, II (September, 1970), 153-189.

location addresses. The programmer uses addresses called "names" or "virtual addresses".

This set of names is called a name space or address space, while the address used by memory is "location" or "memory address".

Since these two types of addresses are different, it requires a mapping system to compact the larger system address space into the smaller memory space.

Address Translation Mechanism

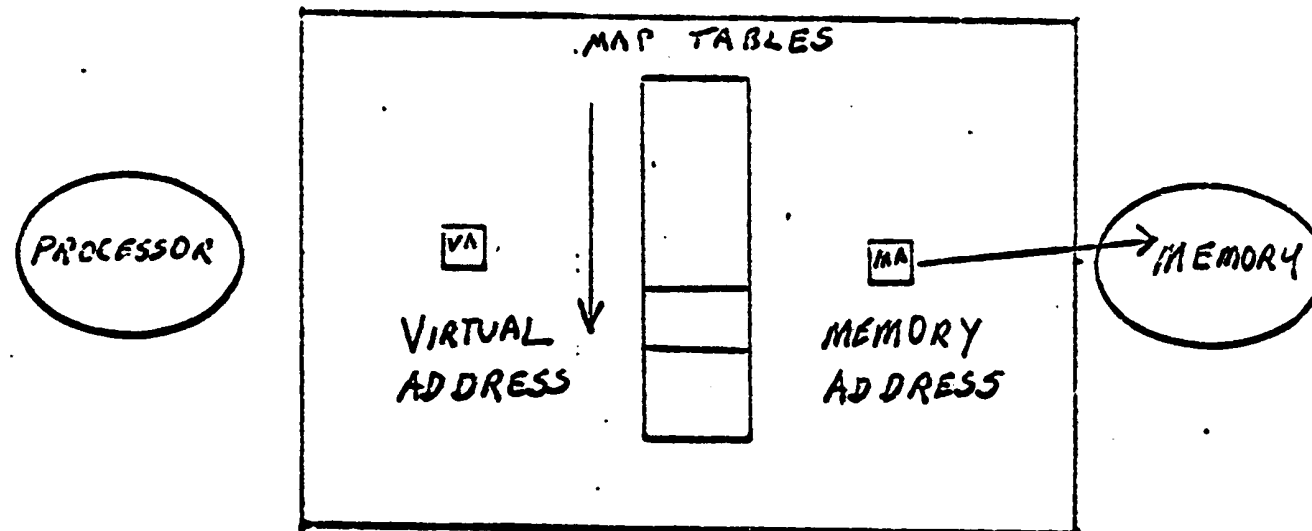


Figure VM-1

The addressing scheme allows the programmer to use a two component address technique which is (s,w). The "s" is the segment and the "w" is the word name which resides within the segment.

The segment is loaded into a contiguous area of memory at base address a. The letter b designates the number of locations s occupies.

Each entry segment is called a descriptor; the sth descriptor contains the base limit information (a,b) for segment s if s is present

in memory, and is blank if it is not present.

Segment Table

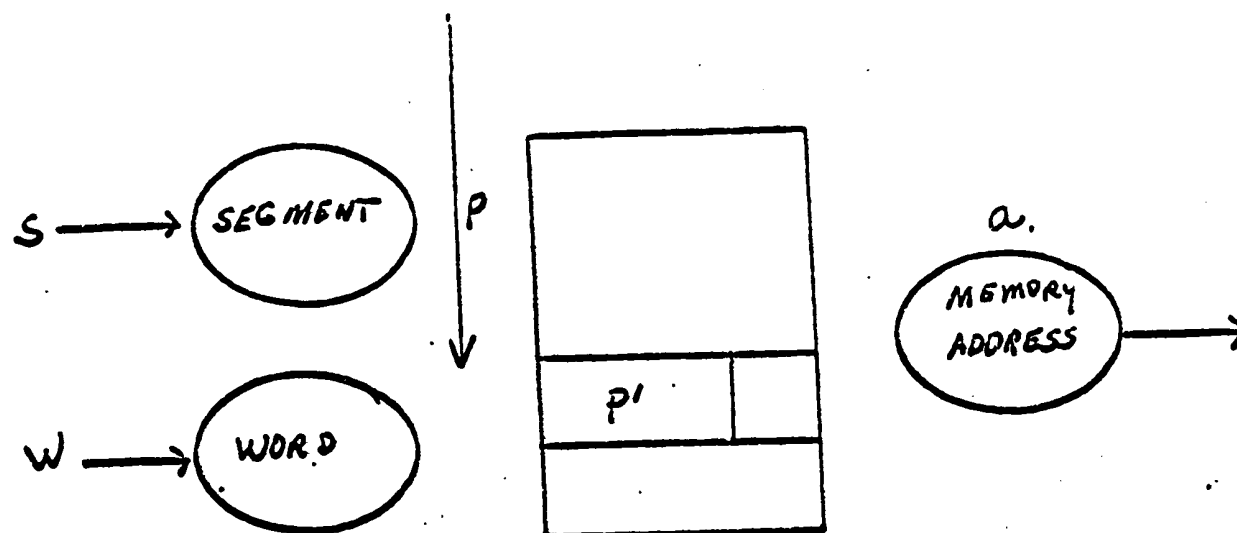


Figure VM-2

Paging divides memory into equal size blocks of locations.

"Page Frames" function as sides of residence for matching size blocks of virtual addresses. A Page serves dual functions:

1. Unit of information storage.
2. Transfer between main and auxiliary storage.

Page Frames are identified by a "frame address" (the location of first word of the page frame). The addresses for pages are written (p, w) where p is a page number and w is word number contained in page p .

Page Table

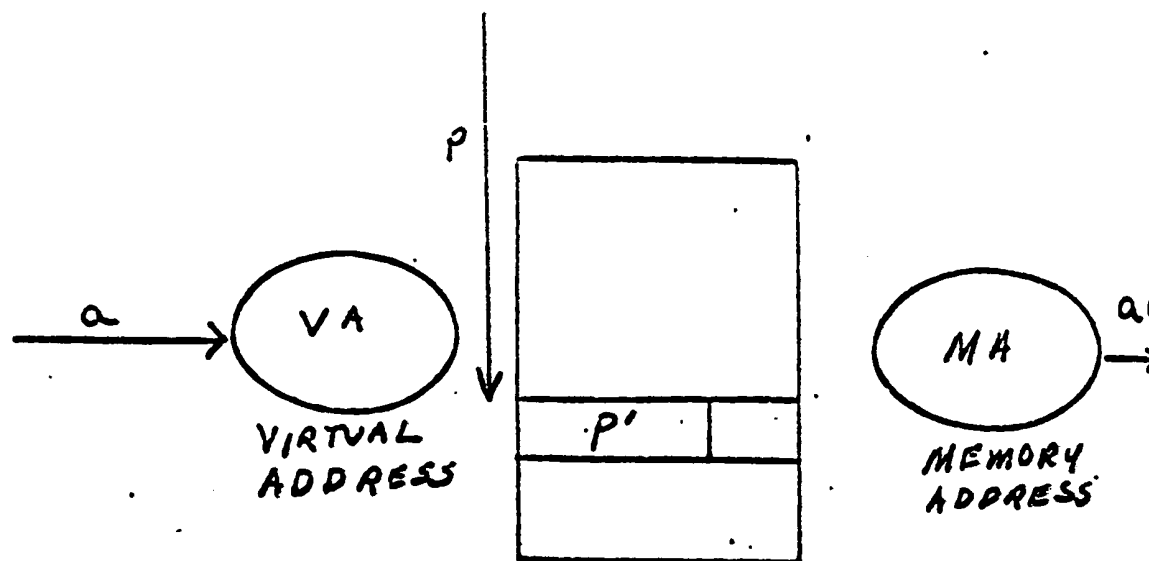


Figure VM-3

Page size was found to work best at forty-five words per page. Segmentation and Paging can be combined, if computer systems have a good selection of register to register operations.

2.3.3.2 RANDOM ORGANIZATION

The Random Organization data structure is based on the principle of retrieving and storing records on the basis of a predictable relationship between a key of the record and the address of the location the record is stored in storage media.

This random organization is used in information retrieval and symbol tables, which is to say, a directory or dictionary type process. One method of developing a directory or dictionary type process, is to use the alphabet as an array in which pointers link the directory to the various addresses which partition storage.

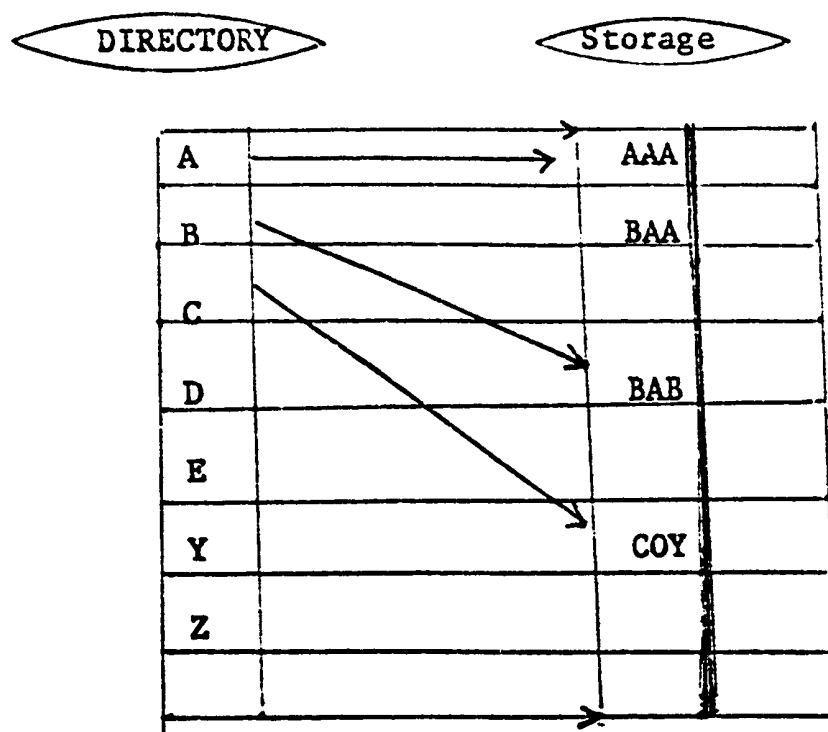


Figure R0-1

The relationship represented here is the binary code of the alphabet letters themselves. The subtables can be chained together. This random organization method has had various names, such as computed table, hash table, key transformation table, but they all achieve the predictable relationship between the key of the record and the location it is to be stored in or retrieved from. The alphabetic method above in Figure R0-1 is an example of the directory look up method. The key of the record was compared to the directory and compare was equal; the direct core address, and the key were brought together.

Another method of dividing the directory is to divide it into ten tables. This method is used frequently with the linear search. This can be done without any ordering of data because the size of the tables will be relatively small. The key is divided by ten, and the remainder (0-9) is used to compute the location to store the data in

one of the ten tables. This is sometimes called a hash table.

Another Random Organization method is to choose some of the bits from the middle of the square of the key. This means choose enough bits to be used as an index to determine the address of any item in the table. Since the square depends on all of the bits of the key, different keys will give rise to different hash addresses in the hash table.

When keys are multiword items such as NAME/ADDRESS/LIST LENGTH, it is possible to take the sum of the bits and transform them into an address or location. Also multiplication and division can be performed on parts of the key, but it is important that the calculation does not come out to zero a good percentage of the time.

Another way of presenting the random address method is to cut the key up into n-bit sections, where n is the number of bits needed for the hash address, and then to form the sum of all the sections. The low order n bits of the sum is used as the calculated address. This method can be used for both single-word keys and multi-word keys.

SECTION III

FORTRAN

III FORTRAN

The Fortran section of this book is slanted more toward the programming proficiency required for the DPMA test; however it also includes some Basic Assembler Language examples for comparison.

The first part of the section presents a little chart for the statements that exist in Fortran I, II, and IV. The following material is pertinent to the type required for the test. The remaining material is presented in the form of questions, with answers given.

It is assumed that the reader has a working knowledge of Fortran (at least to the degree that the basic statements are understood). If the reader does not have this facility as yet, he can still develop the proficiency to pass the DPMA test by analyzation of the small programs given herein.

III FORTRAN

1. Introduction
2. Summary of Statements that Exist in Fortran I, II, and IV
3. Expression Formation
4. Hierarchy of Operation
5. Nested Iterative Input and Output
6. Fortran Shortcuts
7. Buffering with Fortran
8. Syntax
9. Array Search
10. Sort Ascending and Descending
11. Installment Note Program

- 12. Statistics
 - 12.1 BAL Statistics Program--STAT
- 13. Selected Fortran Programs in the Form of Answers to Questions
 - 13.1 Compound Interest Rate
 - 13.2 Plotting
 - 13.3 Integration by Simpsons Rule
 - 13.4 Quadratic Equation
 - 13.5 Linear Equation
 - 13.6 Computed Go To
 - 13.7 The Derivative
 - 13.8 Evaluate a Polynomial
 - 13.9 Functions
- 14. Arithmetic Questions and Answers

2. SUMMARY OF STATEMENTS THAT EXIST IN FORTRAN I, II, AND IV

<u>STATEMENT</u>	<u>FORTRAN</u>	<u>I</u>	<u>II</u>	<u>IV</u>
ACCEPT n, list		X	X	
Arithmetic Statement $v = a$		X	X	X
Arithmetic Statement function name (x, y, z, ..) = h			X	X
ASSIGN i TO n				X
BACKSPACE i			X	X
BLOCK DATA				X
CALL name (h, i, j, ..)			X	X
CALL SWITCH (n, i,)				X
COMMON (h, i, j, ...)			X	X
COMPLEX				X
CONTINUE		X	X	X
DATA ijk.../ e, f, g,.../o, p, q, .../				X
DEFINE DISK (n,m)			X	
DIMENSION u(i), v(i),..		X	X	X
DO n i = j, k, m		X	X	X
DOUBLE PRECISION d, e, f,...				X
END		X	X	X
END FILE i			X	X
EQUIVALENCE (i, j, k, ...), (x, y, z, ...)			X	X
EXTERNAL e, f, g,...				
FETCH (i) d, e, f,...			X	
FIND (i)			X	
FORMAT (s, s, s,)		X	X	X

<u>STATEMENT</u>	<u>FORTRAN</u>	<u>I</u>	<u>II</u>	<u>IV</u>
FUNCTION name (i, j, k, ...)			X	X
GO TO n, (i, j, k, ...)				X
GO TO (J, K, L, ...), i		X	X	X
GO TO n		X	X	X
IF (x) i, j, k (arithmetic)		X	X	X
IF (sense switch n) i, j		X	X	
IF (x) s (Logical)				X
INTEGER a, b, c,				X
LOGICAL a, b, c, ...				X
Logical statement v = x				X
PAUSE		X	X	X
PRINT n, list		X	X	X
PRINT n, iterative list			X	X
PUNCH n, list		X	X	X
PUNCH n, iterative list			X	X
READ n, list		X	X	X
READ n, iterative list			X	X
READ INPUT TAPE i, n, iterative list			X	
READ TAPE i, LIST			X	
READ TAPE i, iterative list			X	
READ (i, n) list				X
READ (i, n) iterative list				X
READ (i) list				X
READ (i) iterative list				X
READ a, b, c,				X

<u>STATEMENT</u>	<u>FORTRAN</u>	<u>I</u>	<u>II</u>	<u>IV</u>
RECORD (i) a, b, c,				X
RETURN			X	X
REWIND i			X	X
STOP		X	X	X
SUBROUTINE name (lk m, n, ...)		X	X	
TYPE n iterative list			X	
WRITE OUTPUT TAPE i, n, list			X	
WRITE OUTPUT TAPE i, n, iterative list			X	
WRITE TAPE i, list			X	
WRITE TAPE i, iterative list		X		
WRITE (i, n) list			X	
WRITE (i, n) iterative list			X	
WRITE (i) list			X	
WRITE (i) iterative list			X	

3. EXPRESSION FORMATION

There are certain rules that will aid the programmer who does not work with Fortran constantly.

1. Variable operands must be previously defined.

Explanation: This is accomplished by reading them or computing them prior to their use in an arithmetic statement.

2. Operators may not occupy adjacent positions.

Explanation: They are separated by parentheses.

$$X = A * (-20)$$

3. The expression must not be made up of mixed mode.

Explanation: This can be either fixed or floating point but not both.

$20 * Y$ creates a mixture of fixed and floating point.

$20. * Y$ is correct since both are floating point.

4. A value may be assigned an exponent of different mode.

Explanation: This is one exception to the mixed mode rule.

5. Spacing can be varied.

Explanation: Items can be written together or spaced apart.

6. Operators may not be "assumed".

Explanation: This is particularly true in multiplication.

They must be written out.

$$10 * Y$$

4. HIERARCHY OF OPERATION

As with manual mathematics, the computer uses its own manner

and sequence.

1. Parentheses are considered first.
2. Exponentiation is carried out second.

Explanation: If there are more than one operand to be raised to a power, parentheses must be used.

Example:

$(gh)^{10}$ must be written $(g. * h.)^{**10}$

3. Multiplication and division are carried out in their order of appearance from left to right.

Example:

$A/3X$ is written $A/ (3. * X)$

4. Addition and subtraction are evaluated last, and occupy the same hierarchical level.

Parentheses

Parentheses are used to set off part of the arithmetic expression. The enclosed part is evaluated first.

$X = D (((C + B) / E^{** 1/2}))$

The sum of C and B are calculated first. Their sum is divided by E and raised to the one-half power. D would then be added to the sum.

5. NESTED INPUT/OUTPUT STATEMENTS VERSUS THE DO LOOP EQUIVALENTS

Fortran I uses the DO loop, while Fortran II uses the nested iterative statements. The pattern set by Fortran I is that of making the innermost loop do the number of iterations. This continues with Fortran II by using the innermost nesting for this purpose. This compares to what we learned in data structures.

The do loops are easier to debug however, since the logic is easier to follow. A comparison can be made with the following read, print and punch routines:

		FORTRAN STATEMENT						HP
		READ 32, K, L, ((B(I, J) I = 2, K), J = 2, L)						
		READ 56, K, L						
		DO 25 J = 1, L						
		DO 25 I = 1, K						
25		READ 58, A(I, J)						
		READ 22(A(I, J), Z(I, J), J = 1, 3), I = 1, 2)						
		DO 20 I = 1, 2						
		DO 20 J = 1, 3						
20		READ 52, A(I, J), Z(I, J)						

Figure FO-5

Nested Input Output Statements (con't.)

		PRINT 25	(W(K,J),	= 1,4),	H(K),	K = 1,4)														
		DO	40K =	1,4																
		DO	30J =	1,4																
30		PRINT 25,	W(K,J).																	
40		PRINT 25,	H(K)																	
		PUNCH 44	((NUM	I,J,K),	K = 1,2)	I = 1,20,2)														
		DO	200 I =	1,20,2																
		DO	200 J =	1,2																
		DO	200 K =	1,2																
200		PUNCH 99,	NUM	(I,J,K)																

Figure FO-5 (Part-2)

6. FORTRAN SHORTCUTS

Space may be saved when a loop has within it an expression, whose variables do not change during the sequence through a loop. The space saving comes from evaluating the expression outside the loop, and retaining the result until it is needed.

```
DO 60 I = 1,N
60 X (I) = Y*Z*W(I)/C
```

An easier method:

```
HOLD I Y * Z/C
DO 60 I 1,N
60 X(I) = HOLD I * W (I)
```

Repeated calculations can be improved by removing the redundant parts of the expressions.

```
X (C * D/A) * CO S(C * D/A).
```

These redundant expressions can be improved by using an area such as HOLD:

```
HOLDI - C * D/A
X = HOLDI * COS(HOLDI)
```

This makes it possible for the computer to make the computation of the expression one time only.

Polynomial Computation

The Fortran writing of:

```
X = C + D*Y + A* Y**2 + E*Y**3
```

If we analyze the requirements for calculation, we find that there are two exponentiations, three multiplications and three additions. The nested form of writing can be used to write the equation:

$$X = C + Y*(D+Y*(A+Y*E))$$

This method of presentation eliminates the multiplication and addition.

7. BUFFERING WITH FORTRAN

Input-output devices such as the teletype, card reader and punch work at speeds much slower than the computer. In the typical installation where computing time is quite costly, it is imperative that the computer be used for computing, with a minimal time for transmission to and from input-output devices.

Most computers now allow the computer to perform computations while data transmission is in process. The output data can be transmitted from memory to an intermediate buffer storage at high speed. This allows the computer to return to its computational tasks while data are being transmitted from the buffer to the output device, at the prescribed timing the output device requires. The process is similar for input data. This can be expanded for auxiliary storage such as disk and tape.

8. SYNTAX

Syntax is a study of language structure, not by a study of words themselves. A language syntax is a set of rules that dictate how the words, or basic elements, of the language are ordered to form meaningful phrases and statements.

Syntax for computer languages is not usually so simple as to have each character for one well-defined use letters representing themselves, symbols representing themselves, and symbols being used

2. Sort from largest to smallest.

Statement Number		FORTRAN STATEMENT
		DIMENSION Y(40)
		DO 100 I = 1, 40
		DO 100 J = I, 40
40		IF (Y(J) - Y(I)) 100, 100, 50
50		TEMP = Y(I)
		Y(I) = Y(J)
		Y(J) = TEMP
100		CONTINUE

Figure FO-10 (Part-2)

The outer DO loop, with one as index, ranks the elements one at a time (either largest to smallest or smallest to largest). Once an element is found, it is stored in Y (1). Then to get the next rank, the inner DO loop must examine Y (1+1) through Y (40). The inner DO loop searches the remaining elements each time to find the smallest (or largest).

12. STATISTICS

Mean, Variance, Standard Deviation

		FORTRAN STATEMENT															
		DIMENSION X(250)															
30		FORMAT															
40		FORMAT															
		SUMX = 0.0															
		SUMQ = 0.0															
		READ 3, N															
		DO 100 I = 1, N															
		READ 4, X(I)															
		SUMX = SUMX + X(I)															
		SUMQ = SUMQ + X(I)**2															
100		ZN = N															
		XBAR = SUMX / ZN															
		VAR = (ZN * SUMQ - SUMX**2) / (ZN * (ZN - 1.00))															
		STDEV = SQRT(VAR)															
		PRINT 9, N, SUMX, SUMQ, XBAR, VAR, STDEV															
9		FORMAT															
		END															

Figure FO-12

Two Approaches to the Standard Deviation

		FORTRAN STATEMENT															
		STD = ((SUMXS - SUMX**2 / AN) / (AN - 1.0))**0.5															
		STD = SQRT((SUMXS - SUMX**2 / AN) / (AN - 1.0))															

Figure FO-12 (Part-2)

Statistics Program

0000		STAT	START 0	00
0001			LISTING 0.0	002
014E			IMR 00350	002
014E	F410 OC5A OC5A	END	ZAP ENN,ZERO	002
0146	0212 0A07 0A06		MVC DIT(120),OUT-1	002
0144	F400 OC5A OC5A		ZAP SMX,ZERO	002
0170	F400 0CA1 OC5D		ZAP SMX,ZERO	002
017A	1A00		SR A,A	002
017A	F400 0A50 OC5D		ZAP INL(17(3,0),ZERO	002
017E	4A00 0A07		AM A,STEP	002
01A2	4000 0A06		CM A,TLIM	002
01AA	07C0 017A		BC 12,0-14	002
01AA	40A0 0A0A		RAS 10,ROCD READ HEADER CARDS	002
01AE	020F 0A0F 07EF		MVC MF01,IMP	002
010A	40A0 0A0A		BAS 10,ROCD	002
010R	020F 070F 07EF		MVC MED2,IMP	003
010E	40A0 0A0A		RAS 10,ROCD	003
01A2	020F 070F 07EF		MVC MED3,IMP	003
01AA	40A0 0A0A	READ	RAC 10,ROCD	003
01AC	0513 0C27 07EF		CLC NME(20),IMP BLANK CARD	003
01A2	47A0 027C		BC R,CALC LAST CARD READ	003
01AA	0A10 0C27 07EF		CLC NME(17),IMP NAME BLANK	003
01AC	47A0 0A0A		BC B,IMR	003
01C0	0507 0C27 0000	GO1	CLC NME(3),SCOR SCORE BLANK	003
01C0	4700 0A0A		BC B,IMR2	003
01CA	F222 0P00 0R00	GO2	PACK SCOR,SCOR	003
0100	0100 0A07 0C5A		MVM SCOR(211),ONE	004
010A	F410 OC5A OC5A		AP ENN,ONE COUNT N CARDS	004
010C	F402 OC5A 0R00		AP SMX,SCOR SUM OF XSUBI	004
01E2	F402 OC5D 0R00		ZAP SCR,SCOR	004
01E4	FC02 OC5D 0R00		MP SCR,SCOR SQUARE XSUBI	004
01EE	F400 0CF1 0C5D		AP SMX,SCR SUM OF SQUARES	004
0100	1A00		SR A,A	004
01F0	F422 0A00 0B50	TAB	CP SCOR,TBL(17(3,0)	004
01FC	47A0 020F		BC 10,DMN SCORE EQU,GT, TAB	004
0200	4A00 0CA2		AM 8,STEP	004
0204	47F0 01F0		BC 15,TAB	004
020A	40A0 0C04	DMN	STM 8,MARK MOVE TABLE DOWN	005
020C	4A00 0C04		LM 9,TLIM	005
0210	0213 9A3F 9B2B		MVC TBL(20,9),TBL(20,9)	005
021A	4A00 0CA2		SM 9,STEP	005
021A	40A0 0C04		CM 9,MARK	005
021E	4720 0210		BC 2,DMN-8	005
0222	0213 9A3F 07EF		MVC TBL(20,8),IMP PUT IT IN TABLE	005
0228	47F0 01A8		BC 15,READ LETS READ NEXT CD	005
			FIND THE MEAN, VARIANCE, AND THE STANDARD DEVIATION	
022C	F450 0C72 0C55	CALC	ZAP MEAN,ZERO	005
0232	D204 0C72 0C5A		MVC MEAN(5),SUMX	005
0238	D100 0C76 0C80		MVM MEAN(1),2M FIRST 4 BYTES HAS	005
023E	FD91 0C72 0C56		DP MEAN,ENN MEAN XXXX.XXC	006
0240	FC91 0C81 0C56		MP SMX,ENN N*SUM(XSUBI)SORD	006
0244	F864 000C 0C5A		ZAP WORK,SUMX (SUM X SUBI) SORD	006
0250	FC62 000C 0C5A		MP WORK,SUMX(2(3)	006
0256	F864 0C58 000E		ZAP SUMX,WORK(2(5)	006
025C	F804 0C81 0C58		SP SMX,SUMX	006

Figure SP-01

Statistics Program

0490	FA40 OC7A OC55	ZAP	MEU,ZEAL	014
049A	FA41 OC7A OUA8	SP	MED,IMII	014
049C	U20A O70E OJAA	MVC	MEDI=15(9),FREQ	014
049E	FA41 OC5H OJ9A	ZAP	SUR,UMU	014
0498	O705 OC77 OJ9A	LAPS	MVC	014
049E	O705 OC72 UC60	EU	MEAN,SOR=3	014
04AA	U20A OJ2E OC73	MVC	PLUT=14(5),MEAN=1 Y SCALE	014
04AB	O202 OJ21 C850	MVC	PLUT=13(1),TBL1=17(12) SCUM	017
04AD	O200 O020 96FF	MVC	PLUT=13(1),MED(19) Y LEGEND	017
04AA	IAAA	SR	A,A	017
049A	IAAA	SR	11,11	017
04BA	4AA0 OCA2	AM	U,STEP	017
04BE	4AA0 OCAH	LAPS	SM	017
04C2	F941 OC7A A803	CP	A,PAIR	10 IN RA
04CA	47C0 O62A	BC	MEU-SCUM=3(2,8)	017
04CC	4AB0 OC6C	LAPS	BC	017
04D0	4AA0 OC70	AM	12,XS	017
04D4	47C0 O4RE	CM	11,DEC	017
04DA	U277 OAA7 O020	CC	11,NTY	017
04DE	40A0 O5AA	MC	12,LAPS	XXXX FILLED IN
04E2	4AC0 OC62	MVC	OUT,PLUT	017
04E6	4AA0 OC6A	BAS	10,PRLN	014
04EA	O202 OJ21 C850	AM	12,STEP	018
04E0	O203 OJ2E OC27	AM	9,BIT	018
04F0	O200 O020 96FF	MVC	PLUT=1(3),TBL1=17(12)	018
04FC	O277 OAA7 OJ20	MVC	PLUT=14(4),NME	018
0502	40A0 O5AA	MVC	PLUT=13(1),MED(19)	018
050A	4AC0 OCA2	MVC	OUT,PLUT	018
050E	4AA0 OC6A	MVC	10,PRLN	018
050E	O202 OJ21 C850	AM	12,STEP	018
051A	O200 O020 96FF	AM	9,BIT	018
051A	U277 OAA7 OJ20	MVC	PLUT=1(3),TBL1=17(12)	018
0520	40A0 O5AA	MVC	PLUT=13(1),MED(19)	018
052A	4AC0 OC62	MVC	OUT,PLUT	019
052A	4AA0 OC6A	BAS	10,PRLN	019
052C	FA40 OC50 OC54	AM	12,STEP	019
0532	FA40 OC7A OC54	AM	9,BIT	019
053A	4720 O498	AP	SOR,ONE	DECA Y SCALE
053C	4260 OJ32	SP	MED,ONE	019
0540	O263 OJ33 O032	BC	2,LAPS	019
054A	4AC0 OC62	MVI	PLUT=1A,X*60	019
054A	O202 OJ21 C850	MVC	PLUT=19(100),PLUT=18	019
0550	O277 OAA7 OJ20	AM	12,STEP	019
055A	40A0 O5AA	MVC	PLUT=1(3),TBL1=17(12)	019
055A	4AC0 OCA2	MVC	OUT,PLUT	020
055E	O202 OJ21 C850	BAS	10,PRLN	020
056A	U277 OJ32 O063	AM	12,STEP	020
056A	O277 OAA7 OJ20	MVC	PLUT=1(3),TBL1=17(12)	020
0570	40A0 O5AA	MVC	PLUT=1A,LEGND	020
0574	4AC0 OC62	MVC	OUT,PLUT	020
057A	O202 OAA7 C850	BAS	10,PRLN	020
057E	40A0 O5AA	MVC	OUT=1(3),TBL1=17(12)	020
0582	O277 OJ32 O01F	BAS	10,PRLN	020
0588	O205 O062 OC93	MVC	PLUT(120),PLOT=1	020
058E	4AC0 OC62	MVC	PLOT=66(6),MED=6	021
		AM	12,STEP	021

Figure SP-Q4

Statistics Program

0402	0202 0071 C050	MVC	PL0T+113),T0L1+17(12)	021
0408	0277 0007 0020	MVC	OUT,PL0T	021
040F	0040 0400	BAS	10,PRLN	021
0412	04C0 0CA2	AM	12,STEP	021
0416	0202 0000 C050	MVC	OUT+113),T0L1+17(12)	021
041C	0100 0400	BAS	10,PRIN	021
0420	0000 01FF	HPR	X'000',0	021
0424	07F0 015E	BC	15,EX4	021
0428	0040 0007 0070	* PRINT AND READ SUBROUTINES FOLLOW		
043E	0710 0034	PRLN	X10 OUT(X'40',120	021
04C2	0700 0400	BC	1,ERRP	022
04C4	0440 04C4	BC	4,PRLN	022
04CA	0441 0034	TI0B	0,X'40'	022
04CE	0277 0007 0000	TI0B	ERRP,X'41'	022
04D4	07FA	MVC	OUT(120),OUT-1	022
04D6	0022 07EF 0030	RCR	15,10	022
04DC	0710 003C	RDC0	X10 INP(X'22'),00	022
04E0	0740 04D4	RC	1,ERR2	022
04E4	0420 04E4	BC	4,RDC0	022
04EA	0421 043C	TI0B	0,X'20'	022
04EC	07FA	TI0B	ERR2,X'21'	022
04EE	020F 0400 04FF	RCR	15,10	022
04F4	0040 0400	TOPS	MVC OUT+20(80),MED1	022
04F8	D20F 0600 070F	BAS	10,PRLN	022
04FE	0040 0400	MVC	OUT+20(80),MED2	022
0A02	D20F 0600 070F	BAS	10,PRLN	023
0A08	0040 0400	MVC	OUT+20(80),MED3	023
0A0C	0040 0400	BAS	10,PRLN	023
0A10	0277 0007 0C00	BAS	10,PRLN	023
0A16	0040 0400	MVC	OUT,MED4	023
0A1A	0040 0400	BAS	10,PRLN	023
0A1E	0209 0007 0C03	BAS	10,PRLN	023
0A24	0040 0400	MVC	OUT(6),MED4+0	023
0A28	07FA	BAS	10,PRLN	023
0A2C	0209 0033 0040	BCR	15,11	023
0A30	07F0 04CC	XS	MVC PLOT+19(10,11),PILL	023
0A34	0000 0100	BC	15,LAP4	024
0A38	07F0 0500	* ERROR SUBROUTINES FOLLOW		
0A3C	0000 0002	ERRP	HPR X'100',0	024
0A40	07F0 0500	BC	15,PRLN	024
0A44	0208 0809 0C30	ERR2	HPR X'002',0	024
0A4A	020F 0007 07EF	BC	15,RDC0	024
0A50	0040 0400	INER	MVC INP+22(12),MSG1	024
0A54	07F0 01C0	MVC	OUT(80),INP	024
0A58	020C 0813 0C07	BAS	10,PRLN	024
0A5E	020F 0007 07EF	BC	15,GO1	024
0A64	0040 0400	INA2	MVC INP+36(13),MSG2	024
0A6A	07F0 01CA	MVC	OUT(80),INP	024
0A6C	0208 0007 0013	BAS	10,PRLN	024
0A72	0040 0400	BC	15,GO2	024
0A76	0000 0003	NEG	MVC OUT(12),MSG3	025
0A7A	07F0 0304	BAS	10,PRLN	025
0A7E	0000 0004	HPR	X'003',0	025
		BC	15,NULL+12	025
		SKER	HPR X'004',0	025
				025

Figure SP-05

Statistics Program

0662	47F0 045C								
06A6	40								
06A7									
06AF									
074F									
079F									
07EF									
07EF									
0A00									
0A03									
0A3F									
0C27	4040 4040 4040 4040 4040 4040 4040 4040								
0C37	4040 4040								
0C3B	D5C1 04C3 4004 C9E2 E2C9 05C7								
0C47	E2C3 0400 0340 04C9 E2E2 C000 C7								
0C54	1C								
0C55	0C								
0C56									
0C58									
0C5D									
0C62	0014								
0C64	0304								
0C66									
0C68	0002								
0C6A	0001								
0C6C	0004								
0C6E	0004								
0C70	005A								
0C72									
0C78									
0C7D	2C								
0C7E	010C								
0C80	F0								
0C81									
0C81									
0C88									
0C8D									
0C8D	4040 0504 4040 E2C3 0409 C5E2 40								
0C94									
0C9C	4040 4040 4040 04C5 C105 40E2 C304 09C8								
0CAC									
0CB3	4040 4040 4004 C5C4 C9C1 0540 E2C8 06								
0CC2	09C5								
0CC4									
0CCB	4040 4040 05C1 09C9 C105 C3C5								
0CD7									
0CDE	4040 4040 4040 4040 E2E3 C105 C4C8 09C4								
0CEE	40C4 C5E5 C9C1 E3C9 0405								
0CF4									
0CF8	4040 4040 4040								
0D05	4020 2120 4820 20								
0D0C									
0D13	E2D8 09E3 4005 C0C7 4005 0440								
0D1F	40								

0D70									
0D98	0140								
0D9A	4070 2070 40C9								
0D40	E7E7 E7E7 F7F7 E7E7 E7E7								
0D44	C0D9 C5D4 E4C5 05C3 E8								
0D83									
0D83	F040 4040 4040 4040 4040 F1F0 4040 4040								
0D83	4040 4040 F2F0 4040 4040 4040 4040 F3F0								
0D83	4040 4040 4040 4040 F4F0 4040 4040 4040								
0D83	4040 F5F0 4040 4040 4040 4040 F6F0 4040								
0D83	4040 4040 4040 F7F0 4040 4040 4040 4040								
0E03	F8F0 4040 4040 4040 4040 F9F0 4040 4040								
0E13	4040 40F1 F0F0								
0E1E									

Figure SP-06

Statistics Program Print Out

BROWN CUR DOG SCHOOL
OBEDIENCE SCORES
ADVANCED TRAINING SCHOOL

NO. SCORES 20 MEAN SCORE 62.05 MEDIAN SCORE 60.00 VARIANCE 6.26 STANDARD DEVIATION 0.31

SCORES	DOG NAME	SCORE
99	SIR CHARLES	99
99	CHARGER	99
98	MFRMAN	98
98	FIFI	98
92	CUM	92
92	FIDO	92
88	SPOT	88
87	PUG	87
81	FLOPPY	81
64	READY	64
56	TROUBLES	56
55	SCAMP	55
52	BONSER	52
43	RUNT	43
36	BIG BOY	36
32	HAMILTON	32
26	DEMON	26
22	ROVER	22
11	LAZY BOY	11
10	GINGER	10

Statistics Program Print Out Number 2

BROWN CUR DOG SCHOOL
OBEDIENCE SCORES
ADVANCED TRAINING SCHOOL

NO. SCORES 20 MEAN SCORE 62.05 MEDIAN SCORE 60.00 VARIANCE 6.26 STANDARD DEVIATION 0.31

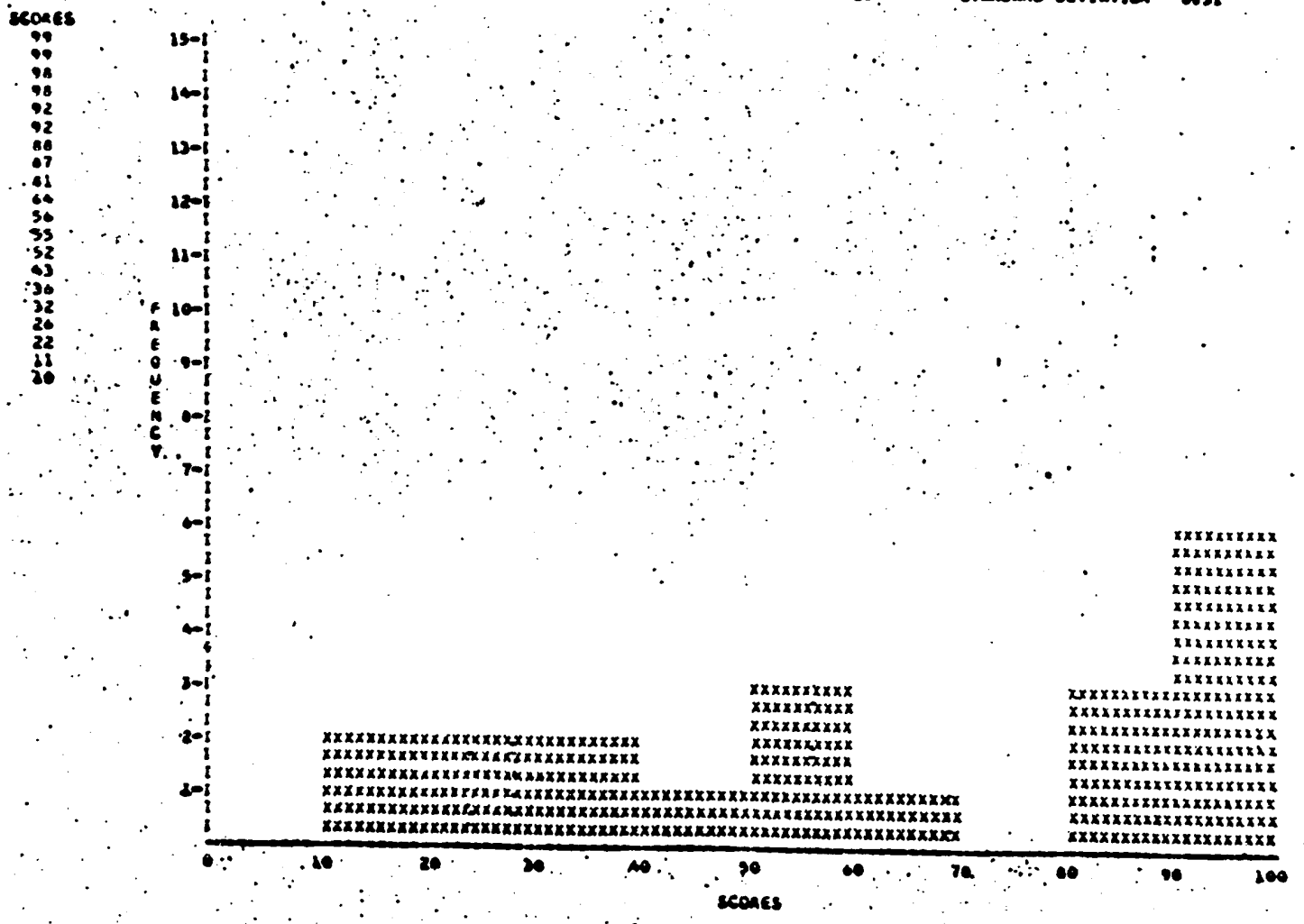


Figure SP-07

13. SELECTED FORTRAN PROGRAMS IN THE FORM OF ANSWERS TO QUESTIONS

Write the following Fortran programs:

13.1 Compound Interest Rate

13.2 Plotting

13.3 Integration by Simpson's Rule

13.4 Quadratic Equation

13.5 Linear Equation

13.6 Computed Go To

13.7 The Derivative

13.8 Evaluate a Polynomial

13.9 Functions

13.1 COMPOUND INTEREST RATE

LINE NO.	STATEMENT	FORMAL	PARA
10	READ (5, 20) BAL, XMON, RATE		
20	WRITE (6, 50) RATE		
30	FORMAT		
40	NR = XMON		
50	RATE = RATE / 12.0		
60	TOTAL = (RATE * ((1.0 + RATE) ** XMON) - 1.0) * BAL		
70	DO I = 1, NR		
80	INT = BAL * RATE		
90	PRIN = TOTAL - XINT		
100	BAL = BAL - PRIN		
110	WRITE (6, 70) I, XINT, PRIN, TOTAL, BAL		
120	FORMAT		
130	CONTINUE		
140	STOP		
150	END		

Figure FO-13.1

13.2 PLOTTING

		DATE	PROGRAM	PAGE
FORTRAN STATEMENT				
	DIMENSION I(100), PLOT(72)			
10	READ (1,1) N			
11	FORMAT			
	WRITE			
	DO 20 J = 1, N			
	READ (1,1) I(J)			
20	WRITE(1,13)			
	IDOT = 10752			
	IBLANK = 8192			
	DO 40 J = 1, N			
	DO 30 K = 1, 72			
	IF (I(J) - K) 21, 22, 21			
21	IPLOT(K) = IBLANK			
	GO TO 30			
22	IPLOT(K) = IDOT			
30	CONTINUE			
	WRITE			
40	FORMAT			
	PAUSE			
	GO TO 10			
	END			

Figure FO-13.2

13.3 INTEGRATION BY SIMPSON'S RULE

STATEMENT NUMBER		DATE		SYMBOLIC OPERATIONS		CYCLES		
				INSTRUCTIONS		MINUTES		
FORTRAN STATEMENT								
		DIMENSION						
5	READ (1, 10) N, H							
10	FORMAT							
	I = N+1							
	READ (1, 20) (Y(J), J=1, I)							
20	FORMAT							
	YE = 0.							
	DO 25 J = 2, N, 2							
25	YE = YE + Y(J)							
	YD = 0.							
	DO 26							
	M = N-1							
	DO 26 J = 3, M, 2							
26	YD = YD + Y(J)							
	R = Y(J) + Y(I)							
	S = R + 4. * YE + 2. * YD							
	AREA = S/3. * H							
	WRITE (1, 30) AREA							
30	FORMAT							
	PAUSE							
	GO TO 5							
	END							

Figure FO-13.3

13.4 QUADRATIC EQUATION

		FORTRAN STATEMENT									
5	READ (1,10) A, B, C										
10	FORMAT										
	$X = B^2 - 4 * A * C$										
	$DF(X) (0, 50, 20)$										
20	$Y = (-B - \text{SQRT}(X)) / (2. * A)$										
	$Z = (-B + \text{SQRT}(X)) / (2. * A)$										
30	WRITE (1,40) Y, Z										
40	FORMAT										
40	PAUSE										
	GO TO 5										
50	$Y = -B / (2. * A)$										
	$Z = Y$										
	GO TO 30										
60	GO TO 41										
	END										

Figure FO-13.4

13.5 LINEAR EQUATION

PROGRAM		NAME	PROGRAM	DATE	TIME	DATE	TIME	DATE	TIME
5	DO 20 J = 1, 7								
	DO 20 I = 1, 3								
	IF (I + J - 8) 10, 10, 20								
10	X = 4. * I + 3. * J								
	WRITE								
15	FORMAT								
20	CONTINUE								

Figure FO-13.5

13.6 COMPUTED GO TO

FORTRAN STATEMENT

10	READ (1,12) I, S								
11	FORMAT								
	GO TO (20, 30, 40) I								
20	T = S**2								
25	WRITE (1,26) T								
26	FORMAT								
	PAUSE								
	GOTO 20								
30	T = S**3								
	GO TO 25								
40	T = S**4								
	GO TO 25								
	END								

Figure FO-13.6

13.7 THE DERIVATIVE

PROGRAM		NAME	DATE	TIME	PAGE
FORTRAN STATEMENT					
	10	READ (1,20)H, (X(I), Y(I), I=1,7)			
	20	FORMAT			
		Z = (.75/H) * (Y(5) - Y(3) - .15/H * (Y(6) - Y(2))) +			
		1/(Y(7) - (Y(1)/(60. * H))			
		WRITE (1,30) Z			
	30	FORMAT			
		PAUSE			
		GO TO 10			
		END			

Figure FO-13.7

13.8 EVALUATE A POLYNOMIAL

PROGRAM		DATE	PROGRAMMER	CLASS	SECTION	DATE
FORTRAN STATEMENT						
	SUBROUTINE POLYX (A, X, N, POLY)					
	DIMENSION A(50)					
	POLY = A(1)					
	DO 300 = 2, N					
300	POLY = POLY + X**I * A(I+1)					
	RETURN					
	END					

Figure FO-13.8

14. ARITHMETIC QUESTIONS AND ANSWERS

The answers immediately follow the questions for the first six numbers.

Examples

$$1. \quad V = C + E^{**}(1./N)*(R/2)-4$$

$$V = c + e \quad 1/n. \quad r/2-4$$

$$2. \quad Y = (a + b) \quad 1/6$$

$$Y = (A + B)^{**}(1/6)$$

$$3. \quad r = \frac{a(a + b)}{b^2 - c}$$

$$R = A*(A + B)/(B^{**2}-C)$$

$$4. \quad V = 4/3 \pi r^3$$

$$V = (4./3)*3.14*(R^{**3})$$

$$5. \quad y = \frac{A + B}{C}$$

$$Y = (A + B)/C$$

$$6. \quad x/y^{p-1}$$

$$X/Y^{**P-1}$$

Questions

1. $(3 \cdot B)^{-1/3}$
2. $(2 \cdot D \cdot C)^{(1/2)}$
3. $(3 \cdot A \cdot B)^{(X-3)}$
4. $\frac{1/8 (4 + A)}{3-B}$
5. $3 (A + B)$
6. $5A/2B$
7. $4A + 2$
8. $6(2/A)B$
9. $(6 + 2/A) B$
10. $6 - 2/A - B$
11. $(4B/2C)$
12. $((5.-A)/(3.-B))*C$
13. $B/2+C$
14. $6 + A/2 - B$
15. $6 + A/3B$
16. $6 + 2/A$
17. $6. + 2*A$
18. $A - C + 4$
19. $y = \frac{a^2 + b}{b^2 - 2ab}$
20. $B*2**L-4. *P/C$
21. $3 \cdot \frac{\text{sales} \cdot \text{ocost}}{\text{Xinv} - \text{ucost}} \quad 1/2$
22. $\frac{(y + a)/r}{f + 1} \cdot p^{1/3} + \underline{v \cdot 2}$

$$23. \quad h + (d/e)^2 \cdot \frac{f \cdot g \cdot r}{s + t} \quad 1/3$$

$$24. \quad y - 3mn^2 - r/s^3$$

Answers

$$1. \quad (3*A*B)**(-1/3)$$

$$2. \quad (2DC)^{1/2}$$

$$3. \quad (3AB)^{x-3}$$

$$4. \quad (1./8.)*((4. + A)/(2.-B))$$

$$5. \quad 3.*(A + B)$$

$$6. \quad .5*(2. + B)$$

$$7. \quad 4. * A + 2.$$

$$8. \quad 6.=* 2./A*B$$

$$9. \quad (6 + 2./A)*B$$

$$10. \quad 6. - 2./A - B$$

$$11. \quad (4.*B/((2.*C))**P$$

$$12. \quad \frac{5 + A}{3 - B} \quad C$$

$$13. \quad B/(2.*C)$$

$$14. \quad (6. + A)/(2. - B)$$

$$15. \quad (6. + A)/(2.*B)$$

$$16. \quad 6. + 2./A$$

$$17. \quad 6 + 2B$$

$$18. \quad A - C + 4$$

$$19. \quad y - (A**2 + B)/(B**2 - 2*A*B)$$

20. $b^2 - 4p/c$
21. $(3*SALES * OCOST)/(XINV-UCOST)**0.5$
22. $((Y + A)/4)/(F + 1) * (P**(1./3.) + (V*Z)/(B*W))$
23. $(h + (D/E)**2*((F*G*R)/(S + T))**(1./3.))$
24. $Y = 3*M*N**2 - R/S**3$

Activities

1. Write a Fortran program pertaining to the compound interest rate.
2. Write an installment note program.
3. Prepare an integration model using Simpson's Rule.
4. Write a program and draw a flowchart for the quadratic equation.
5. Write a program to solve linear equations.
6. Prepare a program utilizing the computed GO TO statement.
7. Prepare a program to compute the derivative.
8. Write a program to evaluate a polynomial.
9. Write a subroutine subprogram to invert an array.

Use three arguments:

1. The given array
 2. The array into which the inverted array will be placed.
 3. An integer quantity that tells how long the arrays are. The array will be called TURBACK.
10. Write a function subprogram, which includes a dimension statement, to find the algebraically smallest quantity in

an array with a maximum of five hundred values. The function is to be called LEAST, and it will have two arguments:

1. The dummy array name.
2. An integer variable telling how many numbers are in the array.

Fortran Language Statements

Equivalence Statement

Form:

EQUIVALENCE (D, A, C(3))

1. The variables D, A, and C(3) will be stored in the same location in memory, however they can't be stored together at the same location at the same time.
2. This allows a saving of space, since after one variable is used, the next variable then occupies the space etc.

Subroutine Statement

Form:

SUBROUTINE NAME (A₁, A₂ A₃...)

1. The Dimension statement is non-executable.
2. It tells the compiler the amount of storage to set aside for each variable.
3. Any number of variables, separated by commas, may be written in the statement.

Dimension Statement**Form:**

```
DIMENSION B (40, 40)
```

1. The Dimension statement is non-executable.

Form:

```
READ 13 IMNAM, RATE
13   FORMAT F6.1, F6.3)
```

1. The format statement is non-executable.
2. It specifies input.
3. It specifies output.
4. The F specification transmits only floating point numbers to and from internal storage.

Go To Statement**Form:**

```
GO TO (100, 101, 105, 120),L
```

1. Control is transferred to the statement numbers when the value of the indexing register is equal to each one.

Go To Statement (unconditional branch)**Form:**

```
GO TO 15
```

An unconditional branch is made to statement 15.

Continue Statement**Form:**

```
10 CONTINUE
```

X Specification Statement

Form:

```
                READ 105, SRATE  
105             FORMAT (33X F5.0)
```

1. It describes the number of positions to be ignored when a card is read.
2. Output causes blank characters to be inserted into the positions indicated. (There will be thirty-three blanks in this example.)
3. The numerical value of SRATE is punched into the thirty-fourth columns.

Read Statement

Form:

```
                READ 110, RATE  
110             FORMAT (E12.1)
```

1. It is nonexecutable.
2. It depicts the form in which the data will appear on input.
3. It shows how the data will be converted for storage in memory.
4. It also specifies, for output, the manner in which the data will be converted from memory and how it will appear in the output device for printing or punching.
5. The number is converted to internal floating point, (ENAM).

F Specification Statement

1. It tells the compiler the amount of storage to set aside for each variable.
2. Any number of variables, separated by commas, may be written in the statement.
3. The example presented here has two dimensions and will set aside sixteen hundred locations (forty times forty).

H Specification Statement**Form**

```

                READ 150
150            FORMAT (15H H GO TO TOWN)

```

1. The alphameric information punched in the first fifteen columns of a card replace the fifteen characters that are in temporary storage. These fifteen characters in temporary storage were previously written in statement one hundred fifty.

I Specification Statement**Form:**

```

                READ 105, SRATE
105            FORMAT 33X F5.0

```

1. It describes the number of positions to be ignored when a card is read.
2. Output causes blank characters to be inserted into the positions.

3. It is the last statement of a DO loop, when the DO loop would cause a transfer type statement of the wrong kind.

DO Statement

Form:

```

DO 8 = 1, 10, 3
A(1) = 8
A(1,1) = 1,1
A(1,2) = 1,2
8 PRINT 9, A(1), A(1,1), A(1,2)
9 FORMAT (3 F8.0)

```

1. The DO statement tells the computer to execute repeatedly the statements which follow up to, and including the statement, with the statement number given (it is eight in this case).
2. The third number to the right of the equal sign is the amount of the increment for the index register.
3. After the statements have been executed ten times (n), control passes to the statement number given here as eight.

If Statement

Form:

```
IF(INT-100) 20, 2, 4
```

The expression within the parentheses is evolved:

1. If the value is negative, control is transferred to the second control statement.
2. If the value is zero, control is transferred to the second

control statement.

3. If the value is positive, control is transferred to the third statement.

Arithmetic Statements

Form:

$$b = d$$

1. The *b* is a variable, which can be subscripted.
2. The *d* is arithmetic in type, which is used to form a sequence of constants, operation symbols and variables.
3. The sequence connotes a series of calculations. The expression on the right hand side of the equal sign is completed, then the resulting numerical value is stored in the location assigned to the variable on the left side of the equal sign.

Call Statement

Form:

CALL TAX (BIC)

1. Tax is the name of a subroutine subprogram, and BIC is an expression of fixed or floating-point constants or variables.
2. The CALL statement calls the subroutines and does one of the following:
 - a. It receives results back from the subprogram, or
 - b. Sends data to the subprogram through an argument list.

Print Statement**Form:**

```
PRINT 10 PROD, YTOT
10    FORMAT (F12.2, F10.2)
```

1. The two totals in the Format statement are printed.

Pause Statement**Form:**

```
PAUSE 32
```

1. A halt is taken when the program reaches the pause statement.
2. Continue by pressing the start.

Slash and Repetition Statements**Form:**

```
FORMAT (4/ E 10.2// (3F10.7))
```

1. The first slash causes the typewriter to start a new line.
2. Then skip a line.
3. Print R numbers per line until all numbers in the list are printed.

End Statement**Form:**

```
END
```

1. It informs the compiler the translation is completed.

Type Statement**Form:**

TYPE 100 TAX, ITEM

100 FORMAT (F10.0, F9.0)

1. The two variables are retrieved, converted and printed in the first nineteen positions.

Return Statement

Form:

1. Return indicates a subprogram has been completed and causes an unconditional branch back to the main program.

Stop Statement

Form:

STOP 999

1. It causes a halt.

Read Statement

Form:

READ 10, EARN

10 FORMAT (F10.2)

1. When READ is reached in the program a sufficient number of cards are read to satisfy the list of variables.
2. The format statement describes the data in detail.
3. The data is stored in memory as the floating point EARN.

Punch Statement

Form:

PUNCH 100, STAX, EARN

100 FORMAT (F10.0, 16)

1. This statement is similar to the read statement.

Function Statement

Form:

FUNCTION LOW (P, Q)

1. One argument must be included in the function statement.
2. The function name must consist of six alphameric characters, and the first must be alphabetical.
3. The function statement operates when the name of the function is encountered in the arithmetic statement.
4. Listed arguments move data to the function routine, and a single quantity is returned to the calling program.

SECTION IV

RPG

IV RPG

The purpose of this section is to provide models and information which will aid in the preparation for the DPMA test. The Basic Assembler program presented here is used as a comparison for the invoicing program written in Report Program Generator (RPG).

The reader is advised to review all the pertinent forms for RPG since they are considered important for the test. The latter part of the section provides the necessary information to perform the Swanson Type programs with RPG, and presents the assignment that could make this text last a year or longer.

Object Program Logic

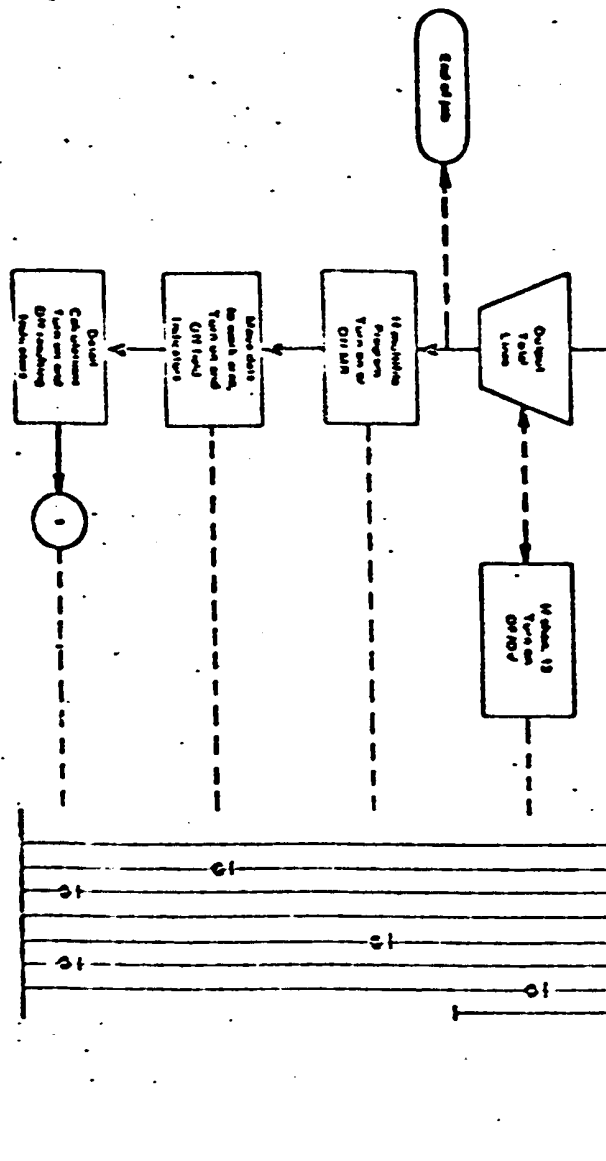


Figure RPG-0

PRE-BILLING CALCULATION WITH INVENTORY CONTROL

This example illustrates one of numerous approaches to an order-processing/inventory control job. The application has been arbitrarily slanted to a distribution business (perhaps a mail-order house) with customer orders to be filled from warehouse stock. An attempt has been made to be reasonably realistic in the application, including the complexities of such a multipurpose operation.

BASIC ASSUMPTIONS

1. A card has been keypunched:

- a. For each item line on a customer order--Card 9, no X in col. 11
- b. For each item line on a customer return--Card, X in col. 11
- c. For each item line on a stock receipt (or purchase-order cards are used as stock receipt cards)--Card 5
- d. For each stock adjustment--Card 6: No X in col. 11 to reduce on hand, X in col. 11
- e. For each item on a stock purchase order--Card 7, no X in col 11 when ordered, X in col. 11 if order is cancelled or reduced
- f. For a new stock item or a change in price, description, warehouse location, etc. (Obsolete master cards are removed manually or, at least, separately from this operation.)

2. An Inventory Master Card file exists, with one card per item carried in stock. Changes to the file are made manually, or in some other data processing operation (i.e., addition and deletion of items, changes in price, warehouse location, etc).

3. It is desired to process customer order-item cards against inventory records before attempting to fill the orders in the warehouse. At the same time, the inventory records will be updated and an up-to-date inventory report prepared. The customer-order cards are thereafter ready for invoicing. (The cards could be sorted by warehouse location prior to invoicing.) A copy of the invoice, or the cards themselves, serve as order-picking medium, i.e., either sequential or bulk picking is employed. If orders are processed once daily on this basis, the inventory records are always up to date.
4. If the quantity on-hand is insufficient to satisfy the quantity in the customer-order card, no partial quantity will be applied for that item. The item order:
 - a. Will be marked "cancelled" if no stock is on order; or
 - b. Will be marked for back order if not previously back-ordered, and provided stock is on order; or
 - c. Will be marked "cancelled" if previously back-ordered.
5. Where previously back-ordered item cards are re-entered, they are to receive priority for available inventory.
6. Some items have a lower unit selling price when at least the specified criterion quantity is ordered by the customer.
7. Stock adjustments are made without attempt at modifying the unit cost of the item.
8.
 - a. Besides price extension, gross profit is to be included in the item detail cards for a subsequent report by merchandise class and division, and by Stock No. (The first digit of Stock No. represents merchandising division, the second the

classification within division.)

- b. Value of inventory on hand (average cost basis) is to be continually available.
- c. Available quantity (on-hand plus on-order) less than an established minimum to be signalled.

PROCEDURAL DETAILS

1. Safeguards

- a. Certain control totals will be carried, partially as audit trails. Control totals are presumed to have been established for the various kinds of transaction cards, so that new on-order totals can be proved out.
- b. Customer-order detail cards that are being cancelled will be identified. If such a card is re-entered, it is selected out, and calculation for it is bypassed.
- c. Matched old master cards (for which new ones are created) will be identified, and selected to a separate stacker. If such a card is accidentally re-entered, the entire stock-number group is selected to separate stacker, and calculation is bypassed.
- d. The entire stock-number group (except the first card) is selected to a separate stacker, processing is bypassed, and the system halts after the second card, whenever there is more than one master card for a group.
- e. The entire stock-number group is selected, and calculations are bypassed, when a master card with a negative on-hand quantity has been read. When a negative on-hand quantity

is created as a result of calculation, the cards from the point of error are selected, and calculations are bypassed.

- f. Whenever the blank trailer card is missing or mispositioned within the group, all cards in the group from the point of the error detection are selected, the system halts, and further calculation is bypassed for the group.
- g. Unmatched transaction cards, including the trailing blank card, are selected, and calculations are bypassed.
- h. If the on-order quantity turns negative, the system halts. The inventory report also indicates this condition.
- i. For known error conditions that affect the new inventory values, the data is omitted.

2. Any merchandise receipts, stock adjustments, and customer returns precede order-item details, so that the customer orders are correctly applied to the latest on-hand status. Stock purchase-order cards are also placed ahead of customer order details, because it was decided not to back-order items for which no stock is on order. Former back-order cards precede other order-item cards to get first chance at on-hand goods.

3. The cards are assumed to be in ascending sequence by Stock No. Inventory master cards are to be in the primary feed of the MFCM-- preceded by a single card to read in today's date. All other cards will be placed in the secondary feed.

A previous operation has placed a blank card at the end of each Stock No. group of secondary file cards. These blank cards will become the new (updated) inventory master cards for stock numbers for which

there are transactions. (These blank cards were merged in on the MFCM of the Model 20, or they could have been merged on a collator.)

4. Stacker Selection

- a. The date header card is directed to stacker 1; any other stacker would do equally well.
- b. All old inventory master cards with stock numbers, for which there are transaction cards in the secondary file, are directed to stacker 1 (the normal stacker-chosen to contain obsoleted cards), because a new inventory master card will be punched and placed in stacker 2.

Each unmatched old inventory master is selected to stacker 2, because no new master is punched in such case.

Stacker 2 ultimately contains the complete up-to-date inventory master file (except for known error-condition cards) consisting of new cards where transactions occurred and old masters where no transactions applied.

- c. Stacker 3, receives the customer order-item cards, ready for warehouse picking (if cancelled and BO (back-orders) are sorted out), or to be sorted on order and account numbers for invoicing.
- d. Stacker 4 has been assigned to unmatched transaction cards (secondary file), and to all other detected error-condition cards.
- e. Stacker 5 has been assigned to stock orders, receipts, adjustments, and merchandise returns. These may also be left together with the other transaction cards by directing them

to stacker 3 instead; they could easily be segregated by sorting cols. 1 and 11.

Card Layouts

IBM INTERNATIONAL BUSINESS MACHINES CORPORATION Form 224 6199 6
 Printed in U. S. A.

MULTIPLE-CARD LAYOUT FORM

Company _____ by _____ Date _____ Job No. _____

STOCK NO.	QTY.	UNIT PRICE	DESCRIPTION	WHSE. LOC.	QTY. SOLD LAST YEAR	QTY. SOLD THIS YEAR TO DATE	QTY. ON ORDER	MIN. STOCK QTY.	AVG. UNIT COST	INVENTORY VALUE, QTY. ON HAND & AVG. UNIT COST	DATE LAST TRANS.
STOCK NO.	QTY.	UNIT PRICE	DESCRIPTION	WHSE. LOC.	GROSS PROFIT, PRICE (QTY. - COST)	CURT. ORDER NO.	ACCOUNT NO.			DATE	
STOCK NO.	QTY.	UNIT COST								DATE	
STOCK NO.	QTY.										DATE
STOCK NO.	QTY.										DATE
STOCK NO.	QTY.										DATE
STOCK NO.	QTY.										DATE

Figure PO

Diagram of Card Flow

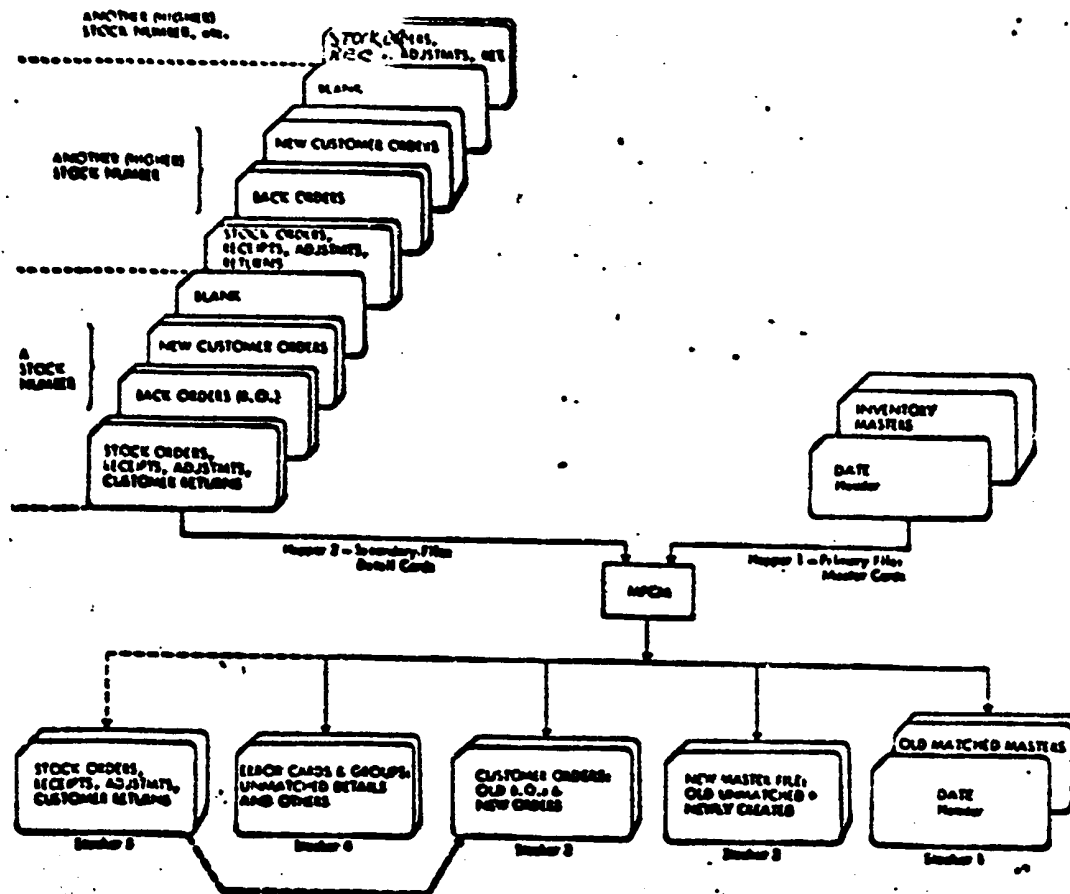


Figure P00

Study of card layouts and the systems flow as seen in Figure P-1 will clarify the details of the operation. The report has been laid out to fit within the 120-position print span of all IBM 2203 and 1403 Printers attachable to model 20. Explanation of specifications sheets follows.

IBM 407, 408, 409, 1403, 1404, 1443, and 2203

6 Lines Per Inch

Printer Spacing Chart

Figure P000

FILE DESCRIPTION SPECIFICATIONS (Figure P-1)

The file inventory master cards is named OLDMASTER, and associated with the primary hopper of the MFCM. It is defined as a combined file (C in col. 15) so that stacker selection may be performed via output specifications, and to allow punching of a code for "obsolete" at output time into those old masters that are replaced as a result of new transactions.

The detail transaction cards are assigned to the file named TRSACTN, and associated with the secondary hopper of the MFCM. Stacker selection is dependent on calculation operations; therefore--and because output is required to some customer-order item cards-- TRSACTN is a combined file.

The input files are in ascending sequence (A in col. 18). A

sequence is required, and must be uniform for the input files, when matching of records is called for in two or more files. If col. 17 is blank, or contains E, for all input files, the LR indicator does not turn on until all input files are exhausted.

The printer is associated with an output file named REPORT.

File Description Specifications

File Name	File Type	File Organization	Length of Record		Record Address Type	Device	Symbolic Device	Name of Label File	Label File for Data
			Block Length	Record Length					
01 OLDMASTR	P	A				MFCM1			
01 TRSACTN	CS	A				MFCM2			
01 REPORT	O					PRINTER			

Figure P-1

INPUT SPECIFICATIONS (Figure P-2)

Because the file OLDMASTR is specified ahead of the TRSACTN file, it is therefore the primary file; i.e., matching cards from the OLDMASTR file are processed ahead of their matching TRSACTN file cards.

Line	File Name	Record Identification Codes							Field Location		Field Name	Field Indicators				
		Position	Length	Code	Position	Length	Code	Position	Length	Code		Start	End	Code		
00	OLPMAETRM	01	1	C												
00								2	7	STKNO	L	M				97
00								8	11	ONHAND						99
00								12	16	PRICEA						
00								17	21	PRICEB						
00								22	22	CRITQY						38
00								23	24	UNMEAS						
00								25	39	DESCR						
00								40	42	WHSLOC						
00								43	47	LASTYR						
00								47	47	NEVITM						
00								48	52	YRDAY						
00								53	56	ONORDR						
00								57	57	LEADT						
00								58	61	MINQTY						
16								62	66	AVGCST						
17								67	76	INVAL						
18								76	80	TRNSDA						

Figure P-2

Input Specifications

Line	File Name	Record Identification Codes							Field Location		Field Name	Field Indicators				
		Position	Length	Code	Position	Length	Code	Position	Length	Code		Start	End	Code		
00		04	09	1	C											
00								76	80	DATE						
00	TRACIN TR	11	1	C5												
00	OR	12	1	C6												
00	OR	13	1	C7												
00	OR	15	1	D9												
00								1	1	DOCARD						21
00								2	7	STKNO	L	M				98
00								8	11	QTY						22
00								12	16	UNGCST						
00								76	80	ACCNT						
00		04	09	1	C5											

Figure P-3

INVENTORY MASTER CARDS--OLDMASTR File (Page 02)

The old Inventory Master cards are identified by 0 in col. 1, and assigned indicator 01. Since they are the only card type in the file, apart from the initial single Date card, an alphabetic code is specified in Sequence (cols. 15-16). If any other undefined card types (besides the Date or Master card) appear in the file, the system halts.

Lines 02-18 contain the normal specifications for reading those fields from the old Inventory Master cards that may be needed for processing of the application. Fields defined as numeric are used in calculations, edit operations, or numeric compare. Points of special interest are:

1. Stock No. is defined as numeric to allow formatting in the output by edit word, and to simplify detection of an obsolete master card (see 4 below).
2. The files are matched and sequence checked on Stock No. (M1 in cols. 61-62 for Stock No.).
3. The L1 indicator is turned on for the first card of each stock-number group (L1 in Control Level for Stock No.).
L1 is not used in this program for total printing or punching, it is used solely to inherent connection with matching of files, and L1 is not needed merely because M1 is assigned.
4. Whenever an old Master Card is replaced by a new one, to reflect transactions, the old card is overpunched with an 11 -punch in col. 7 at output time (Fig. P7, line 06), to mark it as obsolete. If such a card is accidentally re-entered next time, indicator 97 on-the 11-punch causes the

Stock No. to read in as negative (a matching-field sequence error does not, however, arise because all zone punches are eliminated from the matching-fields operations of a numeric field).

5. Indicator 99 turns on if the Quantity On-Hand is negative in the old Inventory Master card. Such a card should never appear, because subsequent specifications (Fig. P4) cause output to be bypassed if On-Hand turns negative.
6. Indicator 20 turns on if the Criterion Quantity field is zero. The zero code indicates that only Unit Price A applies, and that the Price-B field is to remain blank both in the report and in a new Inventory Master card.
7. Col. 47 appears twice among the input fields--the first time as part of a normal numeric field: the second time with another name and as a single-column alphameric field. If col. 47 is X-punched (11 -punch), Quantity Sold Last Year does not apply because the item is new this year. The word NEW is then to appear in the report, and the field is only to contain an X-punch in a new Inventory Master card. But a numeric field that is blank or zero with an X-overpunch in the units position will set on an indicator for Zero or Blank--not for Minus. Therefore, the column that contains the X-punch for "new" is separately defined as alphameric. It can then be tested for a minus done by a TESTZ calculation specification.
8. Stacker assignment is not known until calculations are per-

formed. It must therefore be specified at output time.

DATE CARD OLDMASTER FILE (Fig. P3)

The single Date card at the front of the file is identified by an X-punch in col. 1 and assigned indicator 09. The date is stored in a field given the name DATE. It is defined as numeric to allow editing.

No matching is specified for this card. It is therefore processed first. The date card is to enter the normal stacker for the MFCM primary hopper and therefore need not have stacker selection specified. However, when no output operation is to be performed on a combined-file card type, and the desired stacker number is known at input time, a stacker-selection specification, even for the normal stacker, should be given in the input specifications; this maximized I/C overlap. (For a single card in an entire file, this is of course insignificant). The File Name need not be repeated where no others intervened.

TRANSACTION CARDS (Except Blank Trailer) - TRSACTN FILE (Fig. p4)

The four types are identified and assigned separate indicators. The customer-order or merchandise-return item card is checked for digit rather than character 9 because back orders have an X-overpunch in col. 1.

Stacker selection is dependent on calculations, and is therefore assigned in the output specifications. In the case of card type 9 (indicator 15), output to the card is also required; this precludes stacker selection in the input specifications.

Points to note:

1. Indicator 21 is turned on for order-item cards that were previously back-ordered: 11/9 in the low-order, or sole, position of a numeric field indicates a negative value. (Back-order cards are so designated at output time-- Fig. P7, line 17.) The field BOCARD is not used in the program; it is assigned only so that a Field Indicator may be set. Alternatively, a separate card-type Resulting Indicator could have been assigned via an OR line.
2. The same name is assigned to Stock No. here as for the OLDMASTR file, to conserve core storage space. No harm is done because there is no situation in this program where the distinction needs to be preserved.
3. When an order-item cannot be filled, and is not to be back-ordered, col. 7 of the card is overpunched with an 11 -punch (Fig. P7, line 18) to designate "cancelled." If such a card is inadvertently reentered, indicator 98 turns on because the 11 -overpunch causes Stock No. to be read as negative.
4. Indicator 22 distinguishes between order-item and merchandise-return cards--both card-type Resulting Indicator 15.
5. The field UNCOST applies only to Receipt cards. No harm is done reading it also from card types with Resulting Indicators 12, 13, and 15, because utilization in the calculation specifications is confined to card type 11

(Fig. P5, line 06). If it were necessary to restrict the input of this field to Receipt cards, the indicator number (11) would be entered in Field-Record Relation (cols. 63-64).

BLANK TRAILER CARD-TRSACTION FILE (Fig. P3)

The trailer cards--destined to become new Inventory Master cards-- are identified by absence of a punch in col. 1, and are assigned indicator 19.

The blank trailer card at the end of each stock-number group in the TRSACTION file is not matched (no entry in Matching Fields) against the OLDMASSTR file; therefore, it is processed immediately after the card it follows in the same file, before the Inventory Master card for the next Stock No.

CALCULATION SPECIFICATIONS (Fig. R4)

In order to minimize the need for conditioning indicators (Indicators, cols. 9-17), branching (GOTO) over entire sections has been employed to bypass a series of inapplicable calculation specifications.

Where practical, specifications lines are discussed sequentially. In some areas, however, it is preferable, for clarity, to relate non-consecutive lines. Note: In several instances, result fields are defined as smaller than the theoretically possible maximum. We assumed that knowledge of the particular business indicated that these field sizes are adequate for the actual figures that could occur.

Calculations Specifications

Line	Code	Indicator	Factor 1	Operation	Factor 2	Result Field	Field Length	Resulting Indicators				Comments
								High	Low	Zero	Compare	
09				SETON								IDENT NEXT CARD
09				GOTO	END							DATE BYPAS CALC
11				SETOF								NEW GRP-CLR ERR
11		LN9	LN92	SETON								AND V. NXT LINE
11		9	LN92	SETON								MISSING BLANK
11				SETOF								CLR 1-CYCLE INC
11		11		SETOF								SEE LINE 14
11		01	STKNO	COMP	OLDNO							DUPL MASTR HALT
11		01		MOVE	STKNO	OLDNO	69					PREPAR DUP TEST
11		M1		SETON								DUPLC MASTR INC
11		99		SETON								NEGATV ON HAND
11		NMR	01N19	SETON								UNMATCHD TRANSA
11		07		SETON								OBSOLETE MASTR
11		01		SETON								BLANK MISPOSTING
11		00		GOTO	END							CANCLD ITM-ORDR
16		01		SETON								IND FOLWG MASTR
17		19		SETON								IND FOLWG BLANK
18		90		GOTO	END							GRP ERROR CONDT
19		THIS COMPLETES GENERAL CHECKS FOR PRINCIPAL ERROR CONDITIONS										
19		19		GOTO	NEWMSY							SKIP TO RELEVNT

Figure P4

Line	Code	Indicator	Factor 1	Operation	Factor 2	Result Field	Field Length	Resulting Indicators				Comments
								High	Low	Zero	Compare	
01				MOVE	AVGCST	OLDCST	52					SAVE OLD AVGCST
01		MR		GOTO	END							OLD MASTR COMPL
01		NMR		GOTO	NEWMSY							SKIP TO RELEVNT
15		N15	N13	ONORDR	ADD	QTY	ONORDR					UPDATE ON-ORDER
15		13		AVGCST	MULT	QTY	CSTEXT	82				VALUE OF ADJSTS
15		13		UNCOBT	MULT	QTY	CSTEXT					VALUE OF RECPTS
15		N15	N13	INVVAL	SUB	CSTEXT	INVVAL					UPDT VAL-RCT2AJ
15		N15	N13	ONHAND	SUB	QTY	ONHAND					UPDT ON-RCT2AOT
15		90		GOTO	END							ERROR CONDITION
15		17	26	INVVAL	DIV	ONHAND	AVGCST					REVISE AVGCST
15		N15		GOTO	END							CROS 3/6/7 FINI
23				ONHAND	SUB	QTY	ONHAND					APPLY ORDR/RTN
23		23		ONHAND	ADD	QTY	ONHAND					RESTOR -QTY/ORD
23		N23	N21	ONORDR	COMP	0						ESTAB. BO. INDCTR
23		23		GOTO	END							ITEM UNAVAILABLE

Figure P5

Calculations Specifications (con't.)

Line	Indicators	Factor 1	Operation	Factor 2	Result Field	Field Length	Recording Indicators			Comments
							Imp	Sum	Trans	
01										
02	93	CRITOT	SUB	QTY I	CRTEXT	1			27	PRICE CRITERION
03	93	CRITAT	ADD	QTY	CRTEXT	1			27	PRICE CRITERION
04			Z-ADD	PRICER	UNPRIC	52				TRANSFER PRICER
05	93		Z-ADD	PRICER	UNPRIC	52				SUPERSEDE BY B
06		QTY	MULT	UNPRIC	EXTPRI	62				PRICE EXTENSION
07		QTY	MULT	AVGCST	CRTEXT	1				PREP GRG PROF
08		EXTPRI	SUB	CRTEXT	CRSPRO	52				CALC PROF/LOSS
09		YATDAT	ADD	QTY	YATDAT					UPDATE SALES/YR
10			GOTO	END						ORDR/RTRN FIN
11		NEWMYT	TAG							START NEW MASTER
12			TESTE		NEWLYM				30	NEW THIS YEAR
13		ONHAND	ADD	ONORDR	AVAIL	60				AVAILABLE
14		MINQTY	COMP	AVAIL					31	MIN EXCOS AVAIL
15		AVGCST	COMP	AVGCST					32	COST TREND
16		ONHAND	MULT	AVGCST	INVVAL					UPDATE INVVAL
17		TOTOND	ADD	ONHAND	TOTOND	60				TOTAL ON-HAND
18		TOTORD	ADD	ONORDR	TOTORD	60				TOTAL ON-ORDER
19		TOTAL	ADD	INVVAL	TOTAL	122				TOTAL INV VALUE
20		END	TAG							BEGIN DTL OUTPT

Figure P6

Date Card (Card-Type Resulting Indicator 09)
(Fig. P4, lines 01 and 02)

No calculation operations are performed on this card. Indicator 93 is turned on (line 01) solely for use in a subsequent check on proper card-type sequence (line 05). The entries in line 02 cause branching to the end of the calculation specifications (page 06, line 20), so that No. 9 need not be specified in Indicators in subsequent lines.

Error Control - (Fig. P4, lines 03-18)

Calculation specifications are employed to test for certain error conditions. Where an error is recognized that affects only the

individual card, calculations are bypassed for that card, and the card will be selected by output specifications to stacker 4. Where the effect pervades the entire stock-number group, all calculations for the group are bypassed from the point of error recognition, and those cards will be selected to stacker 4. For certain error situations, the system is also halted.

Indicator 90 is set on for all of the major error conditions tested for, and is used to specify the bypassing of calculations and the selection (see output specifications) of the group to stacker 4. Specifications line 03 clears indicator 90 at the beginning of each control group (i.e., stock-number group), so that the error actions do not carry through to the next group.

Indicator 91 is turned on in line 17 if indicator 19 (blank card) is on. Next program cycle, indicators 90 and H2 are turned on if indicator 91 is still on when the instructions in line 14 are reached by the program. However, if indicator L1 (first card of control group) is on when the instructions in line 07 are reached, indicator 91 is turned off.

Thus, an error is signalled (90 and H2 are turned on) if there is no control break (L1) following a blank card (91 turned on by 19). Trailer card present but not at end of group.

Duplicate Master or Sequence Step-Down

In line 08, the stock number in the old Inventory Master card is compared algebraically with that of the previous old Master card. If the number is the same (duplicate master) or lower, H1 is turned on to halt the system after the card has been processed. In line 09,

the next master card is processed.

In line 10, indicator 90 is turned on if H1 was turned on in line 08, so that all processing for the remainder of the group will be bypassed, and the cards selected to stacker 4.

Note: Because the matching fields assigned in the input specifications were defined as numeric (line 03 of page 02, and line 08 of Fig. P3), the sequence check performed as a result of the M1 specification ignores sign. For that reason, the H1 indicator is also turned on for a negative comparison result--otherwise a duplicate is not detected if one card is positive and one negative in the stock-number field. However, indicators 97 and 98 also signal a negative stock number, but without a halt.

Obsolete Old Inventory Master Card

Indicator 97 turns on if the Stock No. in the old Master card is negative, signalling reentry into the operation of a previously obsoleted card. In line 13, indicator 90 is turned on if that situation exists.

Negative on-Hand in Old Inventory Master Card

Indicator 99 turns on if the On-Hand field is negative at input time of the old Master card. In line 11, indicator 90 is set on for that condition.

Cancelled Order-Item Card

Indicator 98 turns on when a transaction card with a negative stock-number field is read. This signals reentry of a previously cancelled order-item card.

Indicator 98 is used to specify bypassing of calculations for that card only (see line 15), and its selection to stacker 4, but the remainder of the group is processed normally because it is not otherwise affected.

Unmatched Transaction Cards

The specifications in line 12 cause indicator 90 to be turned on for unmatched cards (NMR), other than Inventory Master cards (N01), and other than blank trailer cards (N19) which are always unmatched.

Bypassing Calculations for the Error Group

In line 18 the program branches to END (line 20 on Fig. P6) when indicator 90 is on. This makes detail output the next operation, omitting all calculations below line 17 on Fig. P4.

Line 19 illustrates use of a Comments card (* in col. 7). It will be printed during generation as punched, but otherwise it does not enter the generation process. (It is checked for proper position, based on cols. 1-6.)

Bypassing Detail-Card Operations on Master Cards (Fig. P4, line 20 and Fig. P5, lines 01-03)

Line 20 of page 04 provides program skipping past all the specifications lines that do not apply to the new Inventory Master card (i.e., the blank trailer card). This minimizes the need for N19 specifications in Indicators in subsequent lines.

In line 01 of page 05 the average Unit Cost from the old Inventory Master card is saved for later determination of cost trend when compared with new merchandise costs.

In line 02, all calculations are terminated for old Inventory Master cards that will be replaced by new ones (i.e., there are matching transaction cards).

In line 03, the program skips--for old Master cards that are to be retained (i.e., there are no transactions)--to the same point at which calculations are resumed for new Inventory Master cards. This permits uniform preparation of report data for both situations.

Merchandise-Receipt, Stock-Adjustment, and
Stock-Order Cards (lines 04-11 of Fig. P5)

In line 04, the On-Order quantity is revised to reflect Merchandise Receipts, new Purchase Stock Orders, and cancellation of Stock Orders. Cards 5 and 7 are so coded in col. 11 that addition provides the proper algebraic operation (see Figure P00). The system is halted if the operation results in a negative On-Order quantity. (Indicator 90 is not turned on, because such an error was not deemed of sufficient significance to require bypassing of the remainder of the group.)

Line 05 provides for extending the cost of a stock adjustment, based on last-known unit cost, so that the value of the inventory may be adjusted (in line 07). A new work field (CSTEST) is set up for the product.

Line 06 provides for the same operation as line 05, but using the specific unit of cost at which new merchandise was received.

In line 07, the extended cost of an Adjustment or Merchandise Receipt is algebraically subtracted from the total Inventory Value of the stock item. The signs in cards 5 and 6 are appropriately coded

(see Card Layout).

In line 08, the On-Hand quantity is updated to reflect Receipts and Adjustments. Indicator 90 is turned on if On-Hand has become negative; further calculations are then by-passed for that stock-number group (by entry in line 09), and the cards from this point on are selected to stacker 4 (output specifications).

In line 10, a new Average Unit Cost is established during processing of Receipt cards, because each of these cards contain unit cost. (In lines 06 and 07 we adjusted the Inventory Value to reflect the cost of the new Receipt proportionately.) The quotient is half-adjusted.

Division by zero is not permitted, nor meaningful. Indicator 26 (turned on in line 08 if On-Hand was greater than zero) is therefore a conditioning indicator.

Line 11 causes termination of calculations for cards 5, 6 and 7.

Order-Item and Merchandise-Return cards cannot cause On-Hand to turn negative. If On-Hand was already negative, entries in lines 08 and 09 caused branching to END. Therefore, indicator 23 turns only for a customer order-item card containing a quantity larger than the positive or zero On-Hand quantity.

Lines 13-15 are executed only to handle the insufficient-stock situation (i.e., indicator 23 is on). In accordance with our Basic Assumptions:

1. No order-item will be partially filled.
2. No item card will be back-ordered if it was previously back-ordered.

3. No item will be back-ordered unless merchandise is on order.

In line 13, the quantity is added back to On-Hand to restore the prior status.

In line 14, indicator 24 is turned on if Quantity On-Order is greater than zero (COMP operation), provided the card was not previously back-ordered (N21--see Fig. P3, line 07). Indicators 23 and 24 determine, in the output specifications, whether the card is to be identified as Back-Ordered or Cancelled (Fig. P7, lines 17 and 18).

Indicator 24 is turned off each cycle (see Fig. P4, line 06) before this point is reached, because line 14 is not executed each time. Incorrect card identification in col. 1 would otherwise be punched when non-back-order cards follow a back-order card.

Line 15 causes branching to the end of the calculation specifications for order-item cards that could not be filled. The specifications in lines 02-10 of Fig. P6 will not be executed for these cases.

On Fig. P6, lines 02 and 03, respectively, set on indicator 27 if the customer-order or merchandise-return quantity is equal to or greater than the Criterion Quantity that qualifies for Price B.

We are only interested in the Resulting Indicator, not the actual result quantity. However, an arithmetic operation requires a result field. In order not to waste core storage space, a field only temporarily needed elsewhere (Fig. P5), but now available, has been utilized. A numeric Compare operation is always algebraic, therefore a more complex routine would have had to be substituted for the ADD operation in line 03 (where QTY is negative) if COMP were to be used instead.

Line 04 places Price into a new field, UNPRIC, which will be used for the unit-price factor in the selling-price extension.

In line 05, Price B is substituted for Price A in the UNPRIC field--but only provided the quantity in the Order-Item or Merchandise-Return card satisfied the criterion (lines 02 and 03) and provided Criterion Quantity was not) (N20--see Fig. P2, line 06): zero in col. 22 indicates that Price A applies in all cases.

In line 06, the quantity in the item card is multiplied by the unit price previously selected (lines 02-05). The new field, EXTPRI, will be negative for a Merchandise-Return card, because quantity is negative.

In line 07, cost of the item sale or return is calculated, using the Average Unit Cost as updated during processing of any stock Receipt cards (Fig. P5, line 10). Again, the same work field CSTEMT, is utilized, because the product is not needed beyond line 08.

In line 08, gross profit is calculated for each item card. For merchandise returns, the sign is automatically reversed: -EXTPRI- (-) CSTEMT = -GRSPRO (unless selling price is less than Average Unit Cost).

In line 09, Quantity Sold This Year is updated for this item card. Returns reduce the value, because quantity in these cards is negative. Because it is possible for returns early in a year to exceed sales, provision is made for a negative total (Fig. P11, line 17, edit word).

Line 10 terminates calculations for card 9.

New Inventory Totals (lines 11-19, Fig. P6)

This section contains the specifications for completing the data needed (1) to punch the new inventory Master cards for stock numbers with transactions, and (2) to print the Inventory Status Report for all stock numbers.

No conditioning indicators are required because the program has been instructed, in earlier lines, to branch past this section-- to END (line 20)--for all card types except blank trailer cards or unmatched (i.e., no transaction) old Master cards.

Output Specifications

Line	Function	Code	Output Indicators		Field Name	Field Length	Field Position	Constant or Edit Word	Starting Line	Ending Line
			1	2						
11	OLD MASTRD	04	00	01						
		OR2	NMR	01						
		OR	MR	01						
			DS							
			NMRNG0							
			MRNG0							
12	TRSACTN	04	00	01	NG0					
		OR4	00	15						
		OR5	11							
		OR5	12							
		OR5	13							
13		D3	NG0	NG0	N22					
		AND	15							
		OR5	NG0	NG0	N22					
		AND	15							
17			24							
18			N24	N23						
19			N21		DESCR	39				

Figure P7

Pre-Billing - With Inventory Control

No	Filename	Seq	Do	Output Indicators		Field Name	Seq	No	Content or Edit Word	Starting Invt. Position
				And	And					
01						UNPRTC		16		
02						EXTPRI		22		
03						UNMEAS		24		
04						WNSLOC		42		
05						CRSPRO		47		
06						STKNO		116	0 - - -	
07						QTY		118	CR	
08						WNSLOC		130		
09						ACCNT		200		
10			02		N96 19					
11						STKNO		7	0	
12						ONHAND		11		
13						PRICEA		16		
14						PRICEB		21		
15						CRITOT		22		
16						UNMEAS		24		
17						DESCR		39		
18						WNSLOC		42		
19										

Figure P8

Output Specifications

No	Filename	Seq	Do	Output Indicators		Field Name	Seq	No	Content or Edit Word	Starting Invt. Position
				And	And					
01						LASTYR		47		
02						NEWYR		47		
03						YRTDAT		52		
04						ONORDR		56		
05						LEADT		57		
06						MINQTY		61		
07						AYGCST		66		
08						INVVAL		74		
09								75	+	
10								75	-	
11						DATE		80		
12						DESCR		116		
13						STKNO		116	0 - - -	
14						ONHAND		116		
15						ONORDR		123	0 -	
16						MINQTY		129		
17						LEADT		130		
18						YRTDAT		139		
19						DATE		150	0 -	

Figure P9

Pre-Billing with Inventory Control

Line	File Name	Code	Output Indicators		Field Name	Line	Constant or Edit Word	Starting Page Number
			1	2				
01	REPORT	0	1	1				
02		OR						
03						51	DAILY INVENT	
04						70	ORY STATUS	
05						85	REPORT	
06						115	PAGE	
07							PAGE 2	
08								
09								
10						59	QUANTITY	
11						64	QTY. SOLD	
12						91	SELLING AVG.	
13						129	VALUE DATE A	
14								
15								
16						22	STOCK # DESCRIPTION	
17						46	UM ON ON L AVAIL	
18						64	MIN. LAST YR. TO	
19						95	PRICE UNIT U/	
20						129	OF VMS LAST	

Figure P10

Output Specifications

Line	File Name	Code	Output Indicators		Field Name	Line	Constant or Edit Word	Starting Page Number
			1	2				
01								
02								
03								
04						67	HAND ORDR T(OMTOD)	
05						75	YEAR DATE STNCRD	
06						94	REDUCED COST D	
07						129	INVENTORY LOC TRANSACT M	
08								
09								
10								
11								
12								
13								
14								
15								
16								
17								
18								
19								
20								
21								
22								
23								
24						26	DESCR	
25						27	UNMEAS	
26						32	ONHAND	
27						39	ONORDR	
28						46	AVAIL	
29						51	MINQTY	
30						59	LASTYR	
31						59	NEW	
32						67	YRDTAT	
33						8	STKNO	

Figure P11

Pre-Billing with Inventory Control

Line	File Name	Output Indicators	Field Name	Length	Constant or Edit Word	Starting Position
01			PRICEA	75		
02		N20	PRICEB	03		
03			CRITQTY	05		
04			AVGCST	23		
05		32		04	+	
06		39		04	-	
07			INVVAL	195	\$0.	
08			WHSLOC	110		
09		N51	DATE	119	/ /	
10		31		128	'	
11		7 2 9 1 4 8				
12				31	FINAL TOTAL	OH
13			TOTOND	32		
14				41	00	
15			TOTORD	53		
16				29	INV VAL	
17			TOTVAL	105	\$	
18						
19		01	TRNSDA	119	/ /	

Line 16 is not needed when there are no transactions; but there is no harm in executing it. Although there is no change in Average Unit Cost when there are no Merchandise-Receipt cards in the group, line 15 (in conjunction with line 01, Fig. P5) provides a uniform method of determining cost trend that sets the indicators appropriately, regardless of whether there has been a Receipt.

Line 11 is the destination point to which the program branched from page 04, line 20 (blank trailer card) or page 05, line 03 (unmatched old Inventory Master card).

Line 12 provides for determining whether the item is new this year (X-punch in col. 47 of old Master card--see line 11, Fig. P2). Indicator 30 will be used in output specifications to control punching into cols. 43-47, and printing in print positions 51-59.

In line 13, the available quantity (On-Hand + On-Order) is calculated for the report.

In line 14, indicator 31 is turned on if the available quantity is less than the minimum specified in the old Master card, so that this condition can be signalled by a symbol in the report (print position 120).

In line 16, the updated Inventory Value is calculated after all transactions have been processed.

Lines 17-19 contain the specifications for summing quantity On-Hand, Quantity On-Order, and Inventory Value for report grand totals.

The first two serve only as audit trails and control totals, to balance out former totals with control totals for Receipts, Adjustments, Stock Orders, Merchandise Returns, Back-Orders, and Order-Item cards. The Inventory Value total is also important as far as a management figure.

Line 20 represents merely the destination point to which the program branched from a number of previous lines when calculations were complete. It is followed by detail-time output.

Output Format Specifications (Fig. P7)

All output is at detail time (D or H in col. 15) except for grand totals, based on LR indicator which terminates the job after total-output time.

Old-Inventory Master Cards--Oldmastr File (Fig. P7, lines 01-06)

Lines 01-03 specify different stacker selection for card types

in an OR relationship. Cards with major errors (indicator 90 - see Fig. P4) are selected to stacker 4; the remainder (i.e., the bulk) are selected to stacker 2 if unmatched (NMR), or stacker 1 if matched (MR). Stacker 1--the normal stacker--need not be specified.

Thus, when a new master card will be created because there were transactions, the matched old Master is directed to pocket 1; if no new Master is created, it is directed to stacker 2, which will also receive the newly punched Masters for groups with transactions.

Indicator 01 is needed:

1. To prevent old Master cards of the following stock-number groups being passed through output operations, without being read, in the program cycles during which matched secondary cards are being processed (MR on).
2. To prevent an old Master card being passed through output operations, without being read, during the detail-time output preceding the reading of the first card at 1p time (MR is then off; thus, NMR would apply).
3. To prevent performance of this output for the Date card (during whose processing NMR applies). The Date card was specified as requiring input only, by the stacker selection designated for it in the input specifications.

Line 04 specifies that obsolete old Inventory Master cards (which are replaced by the trailer card of the matched transaction-card group) are to receive an 11-punch in col. 7. This is the safeguard against accidental reentry of these cards next time (see indicator 97: Fig. P2, line 02 and Fig. P4, line 13).

Note: Indicators in File-Identification lines of card types in an OR relationship are tested in sequence. If indicator 90 is on, line 01 is applied and therefore, N90 is not needed in the next two lines. However, in line 04, N90 is necessary, because each Field-Description is considered separately for all card types in an OR relationship.

**TRANSACTION CARDS: RECEIPTS, ADJUSTMENTS, STOCK
ORDERS, AND ERRORS - TRSACTN FILE (Fig. P7, Lines 07-12)**

Cards of groups with major recognized errors are selected to stacker 4. A previously cancelled order-item card that was inadvertently reentered (indicator 98--see page 03, line 08) is also selected to stacker 4. Fifteen is specified in line 08 (with indicator 98) so that additional cards following a cancelled order-item card are not also selected. Indicator 98, once on, is not reset until the next transaction card other than a blank is read.

Receipt, Stock-Adjustment, and Stock-Order cards are selected to stacker 5, they could instead be directed to stacker 3 with the order-item cards and subsequently sorted apart on card No. (col. 1).

**ORDER-ITEM AND MERCHANDISE-RETURN CARDS-TRSACTN FILE
(Fig. P7, Lines 13-19 and Fig. P8, Lines 01-09)**

By the entries in lines 13-16, Merchandise-Return cards (indicator 22 on--see Fig. P3, line 09) are directed to stacker 5, whereas order-item cards are selected to stacker 3. The Returns cards could also, of course, be selected to stacker 3, and subsequently sorted apart by the X-overpunch in col. 11. The file name need not be repeated in line 13.

Line 17 provides for an 11-overpunch in col. 1 on order-item cards being back-ordered. (See Fig. P5, lines 12 and 14, for indicators.)

Line 18 specifies an 11-overpunch in col. 7 for order-items to be cancelled. (See page 05, lines 12 and 14, for indicators.)

Line 19 on page 07 and lines 01-05 on page 08 provide for punching of the pertinent data. Descriptions (line 19) are not punched the first time these cards are processed. The other fields (lines 01-05) are not punched into cards now being back-ordered or cancelled (indicator 23 on); these fields will be punched into the back-ordered cards when they are reprocessed, if the order-item is then filled.

Lines 06-09 provide for document printing (interpreting) on the Order-Item and Merchandise-Return cards, on an MFCM Model A1.

Warehouse location (line 08) is printed only if the item was filled, because the goods could be at a different location when new merchandise is received and the back-orders are filled.

The other three items are interpreted the first time the card is processed (to facilitate card handling), and are therefore not printed again on previously back-ordered cards.

Stock No., Quantity, and Warehouse Location are printed by print head 1; Account No. is printed by print head 2.

Stock No. (line 06) is edited with hyphens between digit positions two and three, and between the fifth and sixth (the presumed self-check digit). The third hyphen in the edit word is the status portion and identifies a cancelled card. All leading zeros, except the first, are preserved.

Zero Suppress is used to eliminate leading zeros in Account No. (line 09).

Note: These cards hereafter contain all the information needed to:

1. Run invoices.
2. Serve as warehouse picking tickets.
3. Run sales, cost-of-sales, and gross profit reports by stock number and merchandise class.

PUNCHING NEW MASTER CARDS (BLANK TRAILER CARDS)--
TRSACTION FILE (Fig. P8, Lines 10-19 and Fig. P9, Lines 01-11)

The pertinent fixed data from the old Master card and the updated variable information are specified for punching as per the card layout (Fig. P0).

The pertinent fixed data from the old Master card and the updated variable information are specified for punching as per the card layout (Fig. P0).

If Criterion Quantity was 0 (indicator 20 on--see Fig. P2, line 06), the field for Price (line 15) is left blank. If the item is new this year (indicator 30 is on--see Fig. P6, line 12), the single-position alphameric NEWITM field (consisting of an 11) is punched into col. 47, line 02. If the item existed last year, the five-digit numeric field LASTYR is punched into cols. 43-47 (line 01). If cost trend is up (indicator 32 is on), a plus (12/6/8) is punched into col. 75 (line 09); if it is down (indicator 33 is on), a minus (11) is punched (line 10, fig. P9); if there was no change in merchandise cost (indicators 32 and 33 are off), col. 75 is left

however, 10 is only on at detail time; therefore, detail output time is the simplest way to handle the operation.)

The heading consists of constants, with one exception, the output field PAGE is specified. This is the only field name (as contrasted with constants) that can provide other than blank or zeros before the first card has been read (i.e., when LP is on). Page No. 1 will be printed in the first heading line of the first page (it is not possible to start with any other value before a card has been read); it will be incremented by 1 before printing on the first line of each succeeding page. Zero Suppress is specified to eliminate leading zeros.

Lines 08-13 contain the specifications for the first print line of column headings, 3 lines beneath the report heading. The form is single-spaced after printing.

The column headings, too, are to appear on each page (first and overflow pages). The file name need not be repeated.

Lines 14-20 contain the specifications for the second print-line of column headings. A single space follows printing.

Lines 01-06 on Fig. P11 take care of the third print line of column headings. After printing, the form is advanced to the next channel -2 punch.

PRINTING THE ITEM LINES-REPORT (Fig. P11, Lines 07-18 and Fig. P12, Lines 01-10)

Lines 07 and 08 specify that the data is to be printed at detail-time (D in col. 15) while processing either:

1. A former blank trailer card (indicator 19 on) that does

not belong to a recognized error group (N90--see Fig. P4; and Fig. 05, line 08); or

2. An unmatched (NMR) old Inventory Master card (indicator 01 on) that does not belong to a recognized error group (N90).

Thus one line will be printed per stock number, showing the original old Inventory Master card data for items without new transactions (NMR), and the updated information where transaction cards exist.

Points to note:

1. In lines 11, 12, 13, 15, and 17 on page 11, the edit word is designed so that one 0 is printed when the quantity is completely zero, and a minus sign is printed for negative values in fields that can be negative. In lines 01, 02, 04 and 07 on Fig. P12, the edit word provides for printing of .00 when the amount field is all-zero. The edit word in line 07 of Fig. P12 also provides for a floating dollar sign.

2. The maximum number of leading zeros (i.e., all but one) is preserved for the Stock No., in line 18 on page 11, and hyphens separate merchandise class from the remainder of the number, and the principal number from the self-check digit.

3. The dates-lines 09 and 091 on Fig. P12 are edited to be printed with slashes between Month, Day, and Year. There is no point in placing a 0 in the edit word; the date can at most have one leading zero (months 01-09), and its suppression cannot be prevented by an edit-word entry.

4. Line 091 on Fig. P12 illustrates insertion of a specifications

line that had been forgotten initially, by assigning it a line number sequentially between two pre-printed numbers.

5. Lines 15 and 16 on page 11 cause the Quantity Sold Last Year to be printed (in print positions 54-59) if indicator 30 (see page 6, line 12) is off, but the word NEW to be printed instead (in print positions 57-59) if indicator 30 is on (i.e., new item this year).

6. Lines 05 and 06 on Fig. P12 provide for printing a + symbol if the cost trend is up, a - symbol if it is down, and leaving the print position blank if there has been no change in cost since the previous report. (See page 06, line 15, for setting of indicators 32 and 33.)

7. Lines 09 and 091 on Fig. P12 determines whether today's date (DATE) from the Date card or the date (TRNSDA) from the old Inventory Master card is to be printed. If there were transactions i.e., the report data is not printed while a Master card is being processed--NO1), DATE is selected; if there were no transactions (i.e., the report is based on data in the old Inventory Master card-indicator 01), TRNSDA is selected.

8. Line 10 on Fig. P12 provides for printing an asterisk in print position 120 when Quantity Available (i.e., On-Hand + On-Order) is less than the Minimum Stock Quantity (see Fig. P6, lines 13 and 14, for setting of indicator 31).

PRINTING THE GRAND TOTALS - REPORT FILE
(Fig. P12, Lines 21-17)

The line is printed at total-output time, (T in col. 15), after the last data card has been processed (LR indicator on). It must be

at total-output time, because the job is thereafter terminated if the LR indicator is on. The form is upspaced 2 lines before printing, providing 3 blank lines between the last detail line and the grand totals.

The form is advanced to channel 1 afterwards. Constants describing the fields are printed preceding the values. A fixed dollar sign is used in the edit word for Inventory Values.

INVOICING

This report utilizes the order-item cards processed in the previous program example (Pre-Billing with Inventory Control), in conjunction with sold-to and ship-to name-and-address cards. The same mail-order company is assumed, with modifications to illustrate more features.

The example is deliberately kept fairly simple, its main purpose being to provide an illustration of:

1. Printing sold-to name and address side by side, each of three lines, and each from a single card
2. Predetermined total line
3. Summary punching. The summary cards can be used for:
 - a. Accounts receivable
 - b. Sales report by customer
 - c. Sales report by salesman
4. Card-type sequence check by sequence entry (cols. 15-16, input specifications)
5. Table lookup

ASSUMPTIONS

1. The item cards from the preceding example serve as detail cards (customer order-item cards), excluding Merchandise-return cards with 11 -overpunch in col. 11). They are assumed to have been sorted by Warehouse Location and Account No. after the Pre-Billing operation.

2. The heading and detail cards have been previously match-merged so that there are no missing masters or legitimate missing details. (This match-merging could have been done on a collator, or on an MFCM.)

The card with today's date and the starting invoice number (less 1) is placed ahead of this group of cards. The deck is placed in the primary hopper of the MFCM.

3. Name and address are confined to three lines from a single card. The presence of a ship-to card is optional. When it is present, it precedes the sold-to card; when there is no ship-to card, the sold-to name and address are to be printed in both positions.

Placing the optional ship-to card ahead improves throughput and printing of name and address can proceed during processing of the sold-to card. If the sold-to card were placed first, printing of name and address could not be commenced, when there is no ship-to card, until the first detail card is being processed; only then can the program know that no further Name-Address card (namely, ship-to card) must be awaited.

4. The blank cards, which are to become summary cards, are a separate file in the secondary hopper of the MFCM.

5. Arbitrarily, the MFCM is used for the two files; other

Card Flow

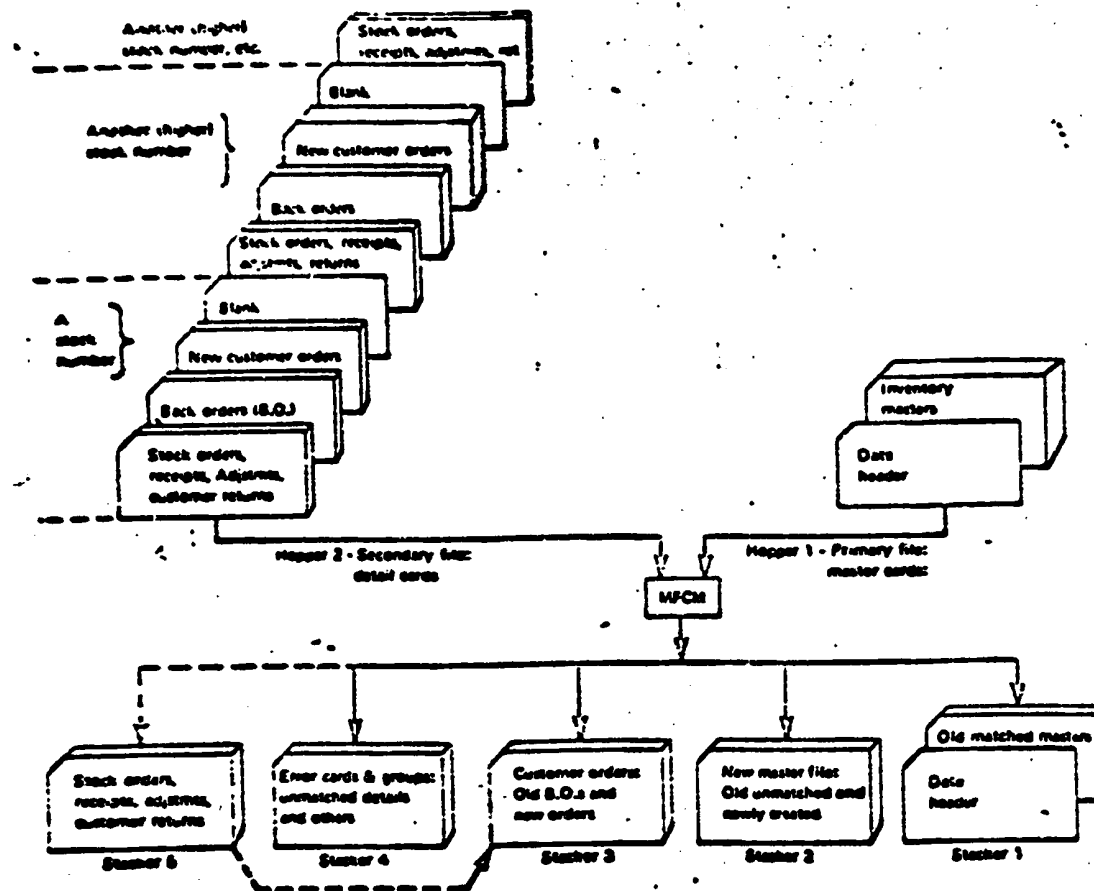


Figure 100

Invoice Layout

INTERNATIONAL BUSINESS MACHINES CORPORATION
PRINTER SPACING CHART
IBM 407, 408, 409, 1403, 1404, 1443, and 2203

6 Lines Per Inch

Print Spans:

- IBM 1403 Models 1 & 4
- IBM 407, 408, 409, and 1403
- IBM 1403 Models 2, 3, 5, 6, 7, 8, 9, 10, 11, 12, 13, 14, 15, 16, 17, 18, 19, 20, 21, 22, 23, 24, 25, 26, 27, 28, 29, 30, 31, 32, 33, 34, 35, 36, 37, 38, 39, 40, 41, 42, 43, 44, 45, 46, 47, 48, 49, 50, 51, 52, 53, 54, 55, 56, 57, 58, 59, 60, 61, 62, 63, 64, 65, 66, 67, 68, 69, 70, 71, 72, 73, 74, 75, 76, 77, 78, 79, 80, 81, 82, 83, 84, 85, 86, 87, 88, 89, 90, 91, 92, 93, 94, 95, 96, 97, 98, 99, 100
- IBM 1443 Models 1, 11, and 12

1	2	3	4	5	6	7	8	9	10
11	12	13	14	15	16	17	18	19	20
21	22	23	24	25	26	27	28	29	30
31	32	33	34	35	36	37	38	39	40
41	42	43	44	45	46	47	48	49	50
51	52	53	54	55	56	57	58	59	60
61	62	63	64	65	66	67	68	69	70
71	72	73	74	75	76	77	78	79	80
81	82	83	84	85	86	87	88	89	90
91	92	93	94	95	96	97	98	99	100

Handwritten text on the chart includes:

- Customer Name: [Illegible]
- Address: [Illegible]
- City, State, ZIP: [Illegible]
- Account No. [Illegible]
- Date: [Illegible]
- Quantity: [Illegible]
- Unit Price: [Illegible]
- Gross Amount: [Illegible]

Vertical text on the left side: "This is a copy of a file on the [Illegible] system."

Figure 1000

Model 20 I/O devices can be used if the File Description Specifications are changed.

6. Stacker Selection has been arbitrarily determined thus:

Date and Invoice-No. heading card-stacker-stacker 1;

Name and address cards-stacker 2;

Detail cards-stacker 2;

Summary cards-stacker 3.

7. A discount percentage is applied to the invoice total based on a customer-type code in the sold-to card. For this, table lookup is employed.

8. Certain identifying data is repeated on overflow pages. Invoice totals are to be printed at a predetermined point on the page.

Fig. 1000 shows input and output formats. Constant headings are not printed by the program, because use of a pre-printed invoice form is usual.

In the explanations that follow for the application example, most of the obvious points will be omitted, as the reader is by this time familiar with them.

FILE DESCRIPTION SPECIFICATIONS (Fig. 11)

The input file, named INPCARDS, is associated with the primary hopper of the MFCM. It consists of one card containing the day's date and the starting invoice number (less 1) and, for each customer account number represented, contains--in this order:

One Ship-to Name-and-Address card (optional)

One Sold-to Name-and-Address card;

At least one Order-Item detail card.

A file of blank cards (named SUMCARDS), which will become the Invoice Summary cards, is to be placed in the secondary hopper of the MFCM.

The printer has been assigned the file named INVOICE.

File Description Specification

File Name	File Type	Stacker		Depth of Inventory		Device	Symbolic Device	Name of Label Set	Label Call for BAK	Comments
		Stack Length	Control Length	Length of Stock	Length of Address Field					
INPCARDS						MFCM1				
SUMCARDS						MFCM2				
INVOICE						PRINTER				

Figure 11

INPUT SPECIFICATIONS (Figure 12)

There is only one input file, named INPCARDS, constituted of four card types.

Date/Invoice-No. Card (Fig. 12, Lines 01-03)

Sequence (cols. 15-16) is alphabetic, because the card appears only once, and does not fall into a sequence within each account-number group.

Stacker selection need not be specified, because 1 is the normal stacker for the primary hopper of the MFCM.

No card-type Resulting Indicator is needed; the card is never referenced, and all calculations are conditioned by indicators of the

appropriate cards.

Ship-To Card

If present, this card is to precede all others of the group; therefore, it is sequence number 01 (cols. 15-16). If present, only one is permitted; therefore, 1 is specified in col. 17. Its presence is optional; therefore, and 0 in col. 18.

Control Level 1 is assigned to customer account number-both (1) to perform end-of-invoice routines, and (2) to guard against cards out of sequence, or missing Sold-To card (see calculation specifications).

Stacker 1 is the normal stacker, and need not be designated.

Sold-To Card (Fig. 12, Lines 09-16)

Exactly one card (1 in col. 17) of this type must be present (no 0 in col. 18), and it follows the Ship-To card, but lower than for the detail cards (Fig. 13, line 01).

Different field names are used for name and address in this card: the name-and-address data from the Ship-To card (if any) is to be printed alongside that from the Sold-To card.

The same field name is used for AC NTNO in all cards because the data should be the same from all cards within the group and therefore need not be saved from card to card. If it is not the same, a control break will occur (L1 is assigned to Account No.).

Indicator 20 is utilized to recognize the first detail card of each invoice (see page 06, lines 12 and 13).

Stacker 1 need not be specified.

function table TABPRC is stored and becomes available as a calculation factor and as output-field data.

The tables are defined in File Extension Specifications--see Fig. 15.

In line 06, indicator H1 is turned on to stop the system after this card--if no Discount-Code match was achieved.

Lines 06-12 provide for the following calculations during total time following the last detail card of each invoice:

1. The invoice gross total is multiplied by the table-supplied percent of discount to establish the discount amount (line 07). (Note that half-adjustment is used, and 4 decimal positions are dropped. There are 2 decimals in INVGRS and 4 in TABPRC, since percentages less than 100 expressed as ratios fall to the right of the decimal point).
2. The discount amount is subtracted from the gross invoice amount to produce the net invoice amount (line 08).
3. The three invoice amount totals (gross, discount, net) are accumulated in three other fields, to provide grand totals (lines 09-11).

The operations in lines 07-11 are executed only when the Discount Code matched an entry in the argument table (indicator H2, and halt the system after this card, if the first card of a control group is not a Name-and-Address card (i.e., neither a Ship-To nor a Sold-To card).

Note: Since the test is made at total time (L1 in cols. 7-8), the first group will not be checked: total time is bypassed on the first card with Control Level specifications. (The test could have

been programmed for detail time instead; but our approach offers the opportunity to remind the reader of the initial total-time bypass.)

Calculation Specifications

Code	Op	Indicators			Factor 1	Operation	Factor 2	Source Field	Add Length	Punching Indicators			Comments
		1	2	3						4	5	6	
00	c	L1	B2			MOVE	NAMED	NAMED					SHIP-TO-SOLD-TO
00	c	L1	B2			MOVE	ADDRSD	ADDRSD					SHIP-TO-SOLD-TO
00	c	L1	B2			MOVE	CTYSTD	CTYSTD					SHIP-TO-SOLD-TO
00	c	B3			INVGRS	ADD	GRSAMP	INVGRS	72				INVC GRS TOTAL
00	c	B2			DSECTO	LOOKUP	TABCOD	TABPRC			23		NO DISCT PRCT
00	c	B2	N13			SETON					41		NO D-CODE MATCH
00	c	L1	23		INVGRS	MULT	TABPRC	AMTDSC	62				CALC DISCT AMT
00	c	L1	23		INVGRS	SUB	AMTDSC	INVNET	72				CALC INVC NET
00	c	L1	23		TOTGRS	ADD	INVGRS	TOTGRS	82				CALC TOTL GRS
00	c	L1	23		TOTDSC	ADD	AMTDSC	TOTDSC	72				CALC TOTL DISCT
00	c	L1	23		TOTNET	ADD	INVNET	TOTNET	82				CALC TOTL NET
00	c	INB1	NB2			SETON					42		NO NAME/ADD CAR
00	c												
00	c												
00	c												
00	c	L1			INVCNO	ADD		INVCNO					INCRMT INVCNO
00	c												
00	c												
00	c												

Figure 14

FILE EXTENSION SPECIFICATIONS (Figure 15)

Two tables are used in this application--one as an argument table (TABCOD) and the other as a function table (TABPRC). For convenience, the two tables are punched alternately in the same card, but this has nothing to do with the manner in which they are employed (argument or function). The table cards (in this instance, a single card) must be loaded at program-generation time.

There are only 14 codes, and all fit in one card; therefore, both the number of table entries per card and per table are the same. The code is a single character (thus, 1 in col. 42), and the percentage is 4 digits long (format XX.XX%). Since the term "percent" means "per hundred," the decimal point must be moved two positions further to the

left when multiplying by a percentage; thus, the field contains 4 decimal positions (not 2).

File Extension Specifications

Line	Table Name	From File Name		Table Name	Number of Table Entries Per Record		Table Name	Length of Table		Comments
		From File Name	To File Name		Number of Table Entries Per Table	Length of Table Entry		Table Name	Length of Table Entry	
01				TASCDP	14	14		4	4	
02										
03										
04										
05										
06										
07										
08										
09										
10										
11										

Figure 15

OUTPUT FORMAT SPECIFICATIONS (Figure 16)

(Refer to Card and Report Layouts.)

Note that all detail output is specified ahead of all total output.

Detail Printing on the Invoice-File (Fig. 16 and Fig. 17, Line 01-11)

This output is performed at detail time (D or H in col. 15).

INVOICE was associated with the printer (see Fig. 12, line 03).

Lines 01-04 on Fig. 16 control the printing of the first print line of the page. The Sold-To name (NAMED-Name sold), read card 2 assigned card-type Resulting indicator 02, is printed in positions 11-29; the Ship-To name (NAMEP-Name ship), read from card 01 (indicator 01) is printed in positions 58-76. Both are printed in the same line on the invoice form.

The printing at the beginning of each Account-No. group takes place as the Sold-To card is being processed (indicator 02 on); at that

time, both the ship-to and sold-to information is available, and can be printed concurrently (if L1, instead of O2 were specified, only data from the first card of the group would be available).

The names are also repeated at the top of overflow pages, at overflow-output time (indicator OF). NLI is specified, so that the old names are not printed at overflow time at the top of one new page-- followed by the new names on the next page from card type O2--when the overflow point and the end of a group coincide.

In the calculation specifications (Fig. 14, line 01), the Sold-To name was moved into the Ship-To name field if there was no Ship-To card; therefore, both names are the same in that case.

Note: If the Ship-To area on the invoice is to be left blank when there is no Ship-To card (rather than repeating the Sold-To card information), lines 01-03 on page 04 (calculation specifications) would be omitted. However, a B must then be placed in col. 39 (Blank After) of lines 04, 07, and 10 of page 06; otherwise, whenever there is no Ship-To card in a group, the data from the last preceding Ship-To card remains in storage, and will be printed.

Lines 05-07 and 08-10 provide the equivalent functions for the second and third lines of the addresses. However, the street addresses and city/state are not repeated on overflow pages.

No entry is required in cols. 17-22 of line 08, because spacing to the miscellaneous-data print line is specified in line 11. The 0 in col. 18 is entered only for compatibility with other RPG's (any entry would satisfy that requirement).

Lines 11-12 (and, as explained below, line 13) control the

conditions under which the miscellaneous data is printed above the first detail line on the invoice.

The form is skipped to the next channel-2 punch before the miscellaneous-data line is printed, and to the next channel-3 punch (first detail line) thereafter.

Note: Instead of utilizing Skip/Before in specification line 11 to reach the miscellaneous-data print line (the simplest way to program this), Skip/After-which is usually more efficient in terms of throughput-can be used in the name-and-address specifications lines. It requires several entries, however, because all three name-and-address lines are printed at the start of a new customer group, but only the name line is printed on overflow pages. The entries in cols. 17-22 (forms control) of specifications (lines 01, 02, 08, and 11) should then read:

Line 01-~~bb~~ 01 02

Line 02-~~b3~~ 01 ~~bb~~

Line 08-~~bb~~ ~~bb~~ 02

Line 11-bb bb 03

The miscellaneous-data line is printed after the name and address for a new group, and ahead of the first detail line. It is also printed in the same position on overflow pages (when overflow does not coincide with the end of a group); but some of the fields are not printed (NOF) on overflow pages.

Because Customer Order No. (ORDRNO) is not available until the first detail item card has been read, the miscellaneous-data line must be printed after the first detail item card has been read, yet

above the regular detail data. Therefore, it is printed during processing of a detail card (indicator 03 in specifications line 12), yet before the print line for the regular detail data (see Fig. 17, lines 01-11). It is to be printed only before the first detail line (apart from overflow identification specified in line 11); therefore, the first detail card of a group must be identified. We chose to accomplish this as follows: The data for the field DSCTCO is supplied by the Sold-To card, where it is never blank. When the first detail card is processed, the DSCTCO field, therefore, contains data. (One of the possible Discount Codes in this example is 0--see Discount Table in Figure 71--but 0 is treated as non-blank in an alphameric field. DSCTCO was defined as alphameric--see Fig. 12, line 15). Indicator 20 is on only when DSCTCO is blank (see page 02, line 15); it is therefore off when the first detail card is processed.

Specifying N20 with 03 in line 12 permits the output to be performed for the first detail card, because indicator 20 is off. As the data from the DSCTCO field is transferred to the output-storage area, the Blank-After (B in col. 39) instruction causes the field to be cleared, and indicator 20 to turn on. The output controlled by the specifications in line 12 will thus never be performed again until another Sold-To card has preceded a detail card because indicator 20 remains on until data is read into the DSCTCO field again. (The entries in line 11 provide for the output at overflow time.) The field DSCTCO was chosen because its data is not needed again in the remainder of the operations for a group.

Note: An alternate approach would be:

Change all L1 specifications to L2.

Then, specify Control Level L1 for ORDRNO (Fig. 13, line 10).

In place of N20 on Fig. 16, line 12, specify L1.

The B in line 13, Fig. 16 is then not needed; nor is indicator 20 in line 15, Fig. 12 required.

This technique might be employed if the contents of all pertinent fields had to be preserved for summary punching.

Specifications lines 13-19 specify the data to be printed in the miscellaneous-data print line.

Although the field DSCTCO is not suppressed for overflow lines (no NOF entry), nothing will be printed from it, because it is blank at that time (see above).

Lines 01-11, fig. 17 contain the specifications for printing of the item detail lines.

The ampersand symbols in the edit word for WHSLOC provide blank spaces on the invoice between the three digits.

If the order item was not filled (i.e., it was back-ordered or cancelled), the Unit Price (UNTPRI) field was left blank (in the previous operation), and indicator 24 is on (see Fig. 13, line 05). Outputs of Unit Price (UNTPRI) and Gross Amount (GRSAMI) are suppressed (N24) when these fields do not apply (i.e., they are blank, with UNTPRI used as the criterion to set indicator 24). Although the fields are blank at input, blank numeric fields are converted to zeros, and .00 would be printed if the output is not suppressed.

The QTY in line 06 pertains to Quantity Ordered; in line 07, it represents Quantity Shipped, although the data is taken from the

same field. The quantity in line 07 is therefore allowed to print only if the order item was filled (N24-UNTPRI field not blank)--it was part of the assumptions in the preceding application example that no partial fills would be made: either stock was sufficient to satisfy the quantity ordered, or the order item was not filled at all (it was then back-ordered or cancelled).

B.O. is printed in the Quantity-Shipped area on the invoice (see specifications line 08) if the order item was back-ordered and not cancelled: indicator 21 is on if the card is identified in col. 1 as a back-order card (see Fig. 13, line 02); indicator 22 is on if the order item was cancelled (see page 03, line 03). All three indicators (24, 21, N22) are needed to establish an active back order, because the item might have been previously back-ordered, and filled or cancelled in the most recent pre-billing pass (see preceding application example).

CANC is printed in the Quantity-Shipped area on the invoice (see specifications line 09) if the item was cancelled (indicator 22--see Fig. 13, line 03).

SUMMARY PUNCHING-SIMCARDS FILE (Fig. 17, Lines 12-20 and Fig. 18, Lines 01-05)

This output is performed at total time (T in col. 15), at the end of an Account-No. control group (L1 in output indicators, line 12), when all totals accumulated from the cards of the group are available.

The file name SUMCARDS was associated with an output file in the secondary hopper of the MFCM (see Fig. 11, line 02). The cards are directed to stacker 3.

Lines 13-20 on page 07 contain punch, rather than interpret,

instructions, because col. 41 is blank or 0.

Lines 01-05 on page 08 contain interpreting instructions for selected fields--they are interpreting, rather than punching, specifications because col. 41 contains a print-head number (i.e., is not blank or 0).

Note 1: The interpreting feature is available only on the MFCM Model A1.

Note 2: Punching of the summary card was specified between detail and total printing to optimize throughput; generally, alternating forms printing and card punching tends to increase throughput.

TOTAL PRINTING ON THE INVOICE-INVOICE FILE
(Fig. 18, Lines 06-16)

The form is first advanced to a predetermined total line (04 in cols. 19-20, specifications line 06). Three lines of totals are then printed at total time (T in col. 15) when the L1 indicator is on (i.e., after each Account-No. group). The form is double spaced between the total lines. In specifications line 11, no entry is needed in col. 18, because forms advance before the grand-total line is specified in line 13. Zero is entered only for compatibility with other RPG's (for that purpose, any digit 0-3 is satisfactory).

Output for the second and third total lines (see specifications lines 08 and 11) is also subject to indicator 23 being on. This suppresses the discount and net amount lines when no match on Discount Code is achieved between the code in the Sold-To card and those in the argument table. While calculation of these amounts was suppressed in such case--see page 04, lines 07 and 08--.00 (not blank) would be

printed for the two amount fields (because of the format of the edit words) if output were not suppressed, and a percentage figure from an earlier LOKUP operation would be printed from TABPRC.

Whenever the total in specification line 07 is transferred to the output-storage area, the field is cleared to zero (B in col. 39) to be ready for accumulation of the total for the next group. Note that the Blank-After instruction could not be entered on page 07 (SUMCARDS) otherwise, the field would be zero before output for printing.

In line 09, note the location of the decimal point in the edit word: in the file-extension specifications TABPRC is defined as consisting of four decimal places so that decimal alignment is correct when calculating the percentage amount. When printing the figure, however, it is to appear as a percentage again--the printing of a decimal point (like any other constant) has no connection with the location of the decimal point for arithmetic operations, as specified in the field definition.

Lines 13-16 control the printing of the grand totals at the end of the report (LR indicator on). The form is advanced to a new invoice page, and all three final totals are printed on the first line.

Output Specifications

Line	Filename	Code	Output Indicators		Field Name	No. of Characters	Constant or Edit Word	Starting Line Position
			And	And				
01	INVOICE	01	01	01				
02					NAMEP	23		
03					NAMEP	76		
04					ADDRSP	54		
05					ADDRSP	77		
06					CTYSTP	86		
07					CTYSTP	85		
08			02	02				
09			03	03				
10			04	04				
11			05	05				
12			06	06				
13			07	07				
14			08	08				
15			09	09				
16			10	10				
17			11	11				
18			12	12				
19			13	13				
20			14	14				

Figure P6

Invoicing

Line	Filename	Code	Output Indicators		Field Name	No. of Characters	Constant or Edit Word	Starting Line Position
			And	And				
01								
02					INSLDC	9		
03					STENO	26		
04					DESER	57		
05					UNHEAS	91		
06					QTY	46		
07					QTY	61		
08						61	B.O.	
09						51	CAVE	
10						69		
11						71		
12	SURCARDS							
13								
14								
15								
16								
17								
18								
19								
20								

Figure P7

Output Specifications

Line	File Name	Code	Output		Field Name	Field Length	Field Position	Constant or Edit Word	Starting Page Position
			Field	Field					
01					ACATNO	05			
02					SALSN	10			
03					INVR	02			
04					INVT	02			
05					DATE	03			
06	INVOICE	T	204	11					
07					INVR	02			
08			11	23					
09					FABPC	02			
10					ANTDSC	02			
11			11	23					
12					INVT	02			
13			01	11					
14					TOTGR	05			
15					TOTDSC	05			
16					TOTNET	05			

Figure P9

CUSTOMER INVOICING

The purpose of this program is to compare the RPG invoicing program to a BAL invoicing program by means of models. The reader, who is familiar with the invoice can compare the methods, while the reader, who is not familiar with the invoice, can benefit two ways, one from the programming and the other from the invoice.

Cards

The invoice summary card contains the first billing number to be used, and the date of the billing. It has a card code of zero.

The customer address card is comparable to an accounts receivable subsidiary ledger in the file, but being processed only if a purchase was made by the customer who has purchased an item between billing dates. One card has the customer number, order number, and

order date. This is used for each customer who has purchased an item between billing dates. One card is made of each item purchased. This card has a card code of zero.

The detail item card contains a customer number, the item number, name of item, quantity of the item, the unit price of each item and the discount allowance to the customer. Later in the program, the net amount and gross amount will be punched and printed on the card.

The invoice summary card is produced at the end of the customer billing. One is produced for each customer that has ordered an item during the billing period. It will contain the customer number, the invoice number, the invoice date, and the total net amount of all items purchased. It has a card code of four.

At the end of all the billing, a new invoice card is produced, which contains only the starting invoice number. The date is punched in by the operator at the date the new invoicing is to be started. This card has a card code of zero. The last card on the deck has a cc of 5, and it is a dummy to facilitate ending the program.

Sequencing

Cards must first be sorted to be put into sequence. The customer card, the order card and the detail cards must be sorted in sequence by customer number.

Sorting

Sort on col. 1 first. This will leave three stacks of cards with a card code of 1, 2 and 3. Then remove cards with card code 2 and sort on cols. 17-16. This is the date. In case one customer ordered

more than one item during the billing period, then sort on col. 15-14. This sorts the month. The year in cols. 89-18 may also be sorted if necessary.

Now place all three stacks together. If the cards are apart to begin with only the order cards will need to be sorted by the date of the month and then the month.

All cards are then sorted on cols. 7-2, the customer number. This should place the stack ready for use. Add the invoice summary and date card to the top of the deck in front of it. Then place the two dummy cards on the back of the deck. These two cards are used to put all the cards through the machine and print and punch the final cards. The final print out will be "end of invoicing."

Every order card should have an address card behind it, followed by at least one detail card and perhaps a second order card from the same person. This person will not have an address card, but the card must be followed by at least one detail card. There may be several address cards in a row, but these are from people who have not purchased anything. They will be skipped.

Invoicing

Place all the data cards in the primary hopper.

Register 13 is being used as a base register, and the first card read is the inventory date card. It gives the starting invoice number and the date of billing. This is dumped into stacker #5. This information is moved from area 5 to area 1 (74 bytes). Then the card code only is blanked out. The inventory number, located in number 2 is packed into number 12.

The second card is read. It could be an address card or it could be an order card. If it is an order card, the information from AR5 is moved into AR3. We then compare the card code, which is F1. It is an F1 so we branch and store to LB1.

At LB1, we will store the information such as the customer number, the order number, the month, the day and the year. The rest of the area is now cleared. This card drops into stacker number 3. BAS 9 takes us back to where we came from, and then we locate a customer address card.

At LB3, the next card is read, and this information is stored into AR2. The customer number is compared. If it does compare with the order card number, the mode of payment, located on the customer address card, is moved into number 9.

The print area is cleared (work area AR4). The name is loaded into AR4 and printed. After the next two steps are skipped, a space is taken and street address, customer, inventory number and the date are loaded into AR4. These are moved to print where they are printed and followed by a space.

We now blank out AR4-47(23) to clear the work area. By doing this we can use the same area for many things. City and State are moved to PRIN and printed, then a space is taken. The order card contents are now printed out (they are stored in LB1). Number 14, which is the customer number, is moved to AR4 and printed.

This will print out as "YOUR ORDER CARD", with customer number and dated. This is followed by a space.

At LB8, the next card is read and stored into AR3. The following

step is a no-op statement the first time through.

LB6--Compare the card code to see if it is a detail card. If it is a detail card, we go to LB4.

LB4--Process the detail item card. First check to see that there are no punches in columns 8 and 27. If no punches are located, load the quantity into number 7 and load price into 6, then multiply these two together. Now we add number 7 to number 5. They will continue to work as an accumulator to produce the total Gross Income. The Mask is moved to AR4 and the gross amount is edited. This information is moved to BRIN and printed.

The new amount is determined by packing discount into number 8 and multiplying the gross amount by the discount. The 01 instruction rounds off this amount. We may lose a penny or gain a penny by using this method, but it averages out over a period of time.

The discount is subtracted from the gross amount and the mask is moved into AR4 followed by the edit. The rest of AR4 (43 bytes), is now blanked out. The net plus gross amount are moved into AR6. This allows us to punch the information into the card.

CONT--Uses print head 2 and at AGAI (XIO AR6-7), we print on the card head the net amount and gross amount. The area is now blanked out. We check for channel 12, if yes, we go back to LB8. This means that we have now read the first invoice card, the first address card and the first detail card. Let's assume that there are no more detail cards for this person, and proceed to go back through the file again. This puts us back to where we would normally read a detail card, LB8.

LB8--A card is read and put into AR3. Now we go past the no-op

step and check to see if it is a detail card. If it is not, it must be an address card. If so, go back to LB8.

Read another card and store the information into AR3. We go past the no-op step and compare to see if it is a detail card (F3). If it is not, we pass through the next step. Is it an address card? If not we assume it is an order card. Therefore we take the BAS9 to LB2. At LB2, we store the order number, the month, the day, year and clear part of this area again. The card falls into stacker 3 and a BCR is taken back to where we came from. A compare is made to determine whether this was a detail card. If not, we proceed to LB5, and print the invoice summary information which will be the ending statement for this customer.

A mask is moved into AR4, and the gross amount is edited and printed. After printing, a space is taken, and the percent of discount is entered.

The discount (no. 10) is packed into number 8. The percent is multiplied times the amount to produce the net amount. A round is taken and the mask is moved into PRIN and edited, after which, we print and space and then blank the rest of AR4.

We take the gross amount and subtract the net amount from it, and move the mask to AR4. Then we edit the net amount into AR4. The net area is cleared by subtracting it from itself. This information is printed and spaced.

Read a card and store this information into AE2. Compare the customer number (let's say they do compare). If they did not compare, we would go back to LB3 and cycle until we did find a number that does

compare, or we would go back and cycle until we came to the customer number that has a higher number than the one located in LB2. If this happens, we branch to SEQR, which is an out of sequence routine, which halts the program.

LB3--Read a card from the order file. At LB3, the information from this card is put into AR3 and we go past the no-op. We now compare the card to see if it is a detail card. If it is, we branch to LB4.

LB4--A compare is made of columns 8-27 to be sure there are no punches. However, this time there is a punch in the card. Branch off to PCHF, which is a read format error routine. This card is checked again to make sure it is a detail card. (We assume it is.) Switch with the MVI instruction EOFQ-1 and move 00 into this condition code. The next instruction moves a blank into AR4 followed by a clear with a MVC instruction.

The print error test (number 13) is moved into AR4+50 for 24 bytes. This error message will print out on the paper as a "punch error" (in the item card). Then blank out the text-error message and place this card into stacker number 2.

Read the next card and place the information into AR3, check the customer number to see if it is a new customer. If it is not a new customer, we go back to CB9 and read another card. We continue this until all of a customer's cards are used up. This is indicated by the appearance of a different customer number.

If a new number appears, we move down to the next MVI instruction and change the card code that changed before. This time it is changed to 15. Then we unconditionally branch back to LB3-4; the

instruction is a BAS9 to LB1. Store the customer number, order number, day and year, then clear the area. This card will go to stacker number 3. The BCR is taken back to LB3.

At LB3, locate a customer's address card to match this order number. If the numbers do match, print the invoice header information, order card contents.

Read a card from the order file, and go to LB4 for the processing of the detail item card. Check for punches, if no, go on down and print the information, and punch and write on the card. The TIOB instruction checks for the overflow which is "43". Check for channel 12. If we were there, go to OVFL, which is the channel 12 condition routine. The MV1 15 into LB8-11 moves the unconditional branch condition into LB8+11. We then branch unconditionally back to LB11 - 4. This instruction is a RC 15 to LB8.

At LB8, we read a card from the order file. The information is moved into AR3. Then BC15 to LB11. The zero condition code has just been changed to a 15 by the MV1 instruction in the overflow area (OVFL).

At LB11, we immediately make LB8+11 a zero condition code again. Compare customer numbers, if not equal, branch to LB12 and at LB12 (BAS9) to LB2.

LB2--Store the order number, month, day and year, and then clear the area. Put this card into stacker 3, and branch back to the statement following the BAS 9 where we just came from (LB12).

The next instruction is a CL1 statement, whereby, we compare to see if this is a detain card. If it is not a detain card, we branch to LB5.

LB5 is an area where the ending routine is effected for this customer. We print the invoice summary information, produce the summary card, increase the invoice number, and now branch back to LB3.

The dummy card is read at LB3, which is a 5 in column 1, and in columns 7-27 is punched "end of invoicing". This information is stored into AR2. We then compare the customer numbers. The customer numbers do not compare because they are all blanks. We drop through the next two routines. Blank out the mode of payment because this area is blank.

The net amount is to be saved so it can be punched into the invoice summary card. The area is cleared and mode is moved. Amount of days, percent of discount and the days net are moved into mode.

The last line is printed out as Mode of Payment 20 days 2% discount OK 30 days Net. Naturally the figures will change with each customer. Check for channel 12, if not, go down and print and skip to column 1 and blank out more of AR4 to clear the area.

Move the card number, which is NO. into AR2 number (24). The customer number, inventory number, the date and the net amount (2), and put AR2 into AR7.

We take the BAS 8 and branch and store to punch.

This action produces a new invoice summary card, which has a card code of 4. A return is now made from the punch routine.

Move AR12, which is the invoice number, into CPL1. Move area 13, which is an inventory date into CPL 2, clear part of AR2, and move customer number into CPL4 and BAS 8 to HEAD, where we write on the card.

The card will contain customer number, invoice number, invoice date and net amount. This card will then be placed into stacker 5.

To increase the invoice number, we take number 19, which is the number 1, and add it to 12, which is the invoice number. Move in the mask and edit the number, and go back to LB3.

At LB3, we are going to try to locate the customer address card number to match the order number, that we have stored in LB2.

The header information is printed by first printing the name, but this time it is not the same, it is the "END OF INVOICING". This is printed out after which we come to the CL1 we have been passing up all the time during the program. We check to see if it is a card code of 5. We assume it is so we go to EOFR, which is the end of the file routine.

EOFR--Move the new invoice number, which is at number 2 to AR4+35. Reset the switch at PNCH-5. This cuts down the amount of information that is going to be printed on the card from 16 down to 8. Now we take the next MVI instruction which moves a 00 into PUER3. This changes the hopper code from 4 to 5. This card will now drop into stacker number 5.

We then move a card indicator code into AR4+34. Move AR9 into AR7, which is the invoice number. We then branch to PNCH and pick out this last card. Only the invoice summary number will be punched. The date will have to be punched in by the operator at the next billing period.

After number 1 is selected to print this same information, the final action is the halt.

Invoicing--Basic Assembler

0156		INVO	START 340		
0156	0000		BASR 13.0		
0156			USING K.13		
			START OF PROGRAM		
015A	4040 037C		BAS R,GET		MOV INVO DATE CR
015B	4420 0005		CID 5,X'20'		SP1, STA, 5
015C	0240 047C 0471		MVC AR1(4),AR5		MOV INPR WK AREA
015D	4240 047C		MVI AR1,X'40'		NRK INR (LND) LINE
015E	4736 05CF 0470		PACK N(12),M(2)		PACK INVO NO
015F	4040 037C	LMUN	MAS R,GET		READ INR CARD
0172	0244 04CA 0471		MVC AR3(7),AR5		MOV INPR WK AREA
017A	95F1 04CA		CLI AR3,X'P1'		IS IT INR CARD 2
017C	4420 0005		CLI 3,X'20'		AND CARD S/S 3
017D	4720 0014		MC 2,LMUN		NO INR CARD
017E	4040 037C		BAS 4,LM1		STORE INR CARD
			LOCATE CUSTOMER ADDRESS CARD		
018A	4040 037C	LM3	BAS R,GET		FINR ADD CRD
018C	4244 04HC 0471		MVC AR2(7),AR5		MOV INPR WK AREA
018E	0505 0475 0480		CLC N(1),AR2+1		COMPARE CUST NO
0194	4720 0032		MC X'2',LM3		NO GET ANOTHER
019C	4740 0324		MC X'4',SETR		ADD CRD NOT SEU
01A0	0207 05CA 04CF		MVC N(4)(1),AR2+67		MOV INPR OF PAY
			PRINT INVOICE HEADER INFORMATION		
01A6	0244 0407 0427		MVC AR4,PR(1)		CLERK WORK AREA
01AC	0213 04E3 0493		MVC AR4+12(20),AR2+7		PRINT NAME
01B2	0244 0427 0407		MVC PR(1),AR4		MOVE TO PRINT
01B4	4040 03A4		BAS R,PUT		PRINT
01B6	95F5 04HC		CLI AR2,X'F5'		LAST CARD CR
01B8	4740 0342		MC R,EIFM		GO TO FINISH
01BA	4040 03EC		MAS R,SP43		SPACE 2
01BC	0213 04E3 04A7		MVC AR4+12(20),AR2+27		PRINT STREET
01BE	0216 0506 0475		MVC AR4+6(123),M(1)		CUSTOMER INVOICE DATE
01C4	0244 0427 0407		MVC PR(1),AR4		MOVE TO PRINT
01C6	4040 03A4		MAS R,PUT		PRINT
01C8	4040 03FE		MAS R,SP42		SPACE 1
01E2	0216 0506 0505		MVC AR4+6(123),AR4+64		MOVE BLANK
01E4	0213 04E3 0480		MVC AR4+12(20),AR2+67		PRINT CITY+STATE
01E6	0244 0427 0407		MVC PR(1),AR4		MOVE TO PRINT
01F4	4040 03A4		MAS R,PUT		PRINT
01F6	4040 03FE		MAS R,SP42		SPACE 2
			PRINT ORDER CARD CONTENTS		
01FC	0223 04E3 05ER	LR7	MVC AR4+12(3A),M(16)		CUSTOMER DATE
0202	0244 0427 0407		MVC PR(1),AR4		MOVE TO PRINT
0204	4040 03A4		MAS R,PUT		PRINT
020C	4040 03FE		BAS R,SP42		SPACE 1
			READ A CARD FROM ORDER FILE		
0210	4040 037C	LMH	MAS R,GET		NO NEXT CARD
0214	0240 04CA 0471		MVC AR3(7),AR5		MOV INPR WK AREA
021A	4700 0242		MC X'0',LM11		OVERFLOW SWITCH
021E	95F5 04CA	LMH	CLI AR3,X'P3'		IS IT ITEM CARD
0222	4740 021E		MC R,LM6		ITEM CARD
0224	95F1 04CA		CLI AR3,X'P1'		IS IT ORDER CARD
0226	4720 004A		MC 2,LMH		NO NO NEXT CRD
022E	4040 037C	LM12	MAS 4,LM2		STORE INR CARD
			PREPARE INVOICE FOR PRINTING ORDER CARD		

Figure INVO-1

Invoicing--Basic Assembler

0232	0505 0675 0607	CLC	M0,AN3+1	2 AND OVER
0234	6700 0000	BC	7,105	Y+S
0236	6000 0300	0AS	6,SPA2	SPACE 1
0240	6700 0000	MC	15,107	GO TO INVOICE NET
* PRINT INVOICE SUMMARY (OPTIONAL)				
0246	0200 0515 0410	LMS	MVC	AM4002(12),AN2
0248	0000 0515 0505		MI	AM4002(12),M15
0250	6000 0300		MAS	M,SPA2
0252	0244 0627 0407		MVC	M1,AN6
0254	6000 0300		MAS	M,M1
0256	6000 0300		MAS	M,SPA2
0258	0201 0500 0500		MVC	AM4053(2),M110
0260	0271 0501 0501		MASK	MIM,M110
0262	0075 0501 0505		MP	MIM,M15(A)
0276	4000 0507		UI	MIM+X*OP
0278	0200 0515 0610		MVC	AM4002(12),MSK2
0279	0000 0515 0507		EU	AM4002(12),MIM+1
0284	0244 0627 0407		MVC	M1,AN6
0286	6000 0300		MAS	M,M1
0288	6000 0300		MAS	M,SPA2
0292	0201 0500 0500		MVC	AM4053(2),AM4052
0294	0055 0505 0502		SP	M15(6),MIM+1(A)
0296	0200 0515 0610		MVC	AM4002(12),MSK2
0298	0000 0515 0505		EU	AM4002(12),M15
0299	0055 0505 0505		SP	M05(A),M05(A)
02A0	0244 0627 0407		MVC	M1,AN6
02A2	6000 0300		MAS	M,M1
02A4	6000 0300		MAS	M,SPA2
02A6	0200 0600 0510		MVC	AM2+24(11),AM4053
02A8	0200 0515 0515		MVC	AM4053(11),AM4052
02AA	0230 0607 0700		MVC	AM4(54),M10E
02AB	0201 0600 0500		MVC	AM4027(2),M15
02AC	0201 0604 0500		MVC	AM4034(2),MIM+2
02AD	0201 0507 0500		MVC	AM4040(2),MIM+6
02AE	0244 0627 0407		MVC	AM4,AN6
02AF	4000 0500		TIIM	JUMP,X*43
02B0	6700 0100		MC	15,MIM
02B2	6000 0300		MAS	M,SKIP
02B4	6000 0300		MAS	M,M1
02B6	6000 0300		MAS	M,SKIP
02B8	0200 0600 0600		MVC	AM4034(16),AM4033
* INCREASE INVOICE NUMBER				
0302	0217 0600 0676		MVC	AM2(24),M1
0304	0222 0600 0600		MVC	AM7(35),AM2
0306	6000 0610		MAS	M,MIM
0312	0200 0521 0605		MVC	CPL1,AM12
0314	0205 0501 0600		MVC	CPL2,AM13
0316	0200 0605 0600		MVC	AM204(15),AM204
0320	0221 0501 0600		MVC	CPL4,AM11
0322	6000 0630		MAS	M,MIM
0324	0205 0675 0607		MVC	M1,AN3+1
* INCREASE INVOICE NUMBER				
0336	0030 0500 0600		AP	M12(4),M114(1)
0338	0207 0670 0500		MVC	AM1(A),MSK1
0340	0007 0670 0500		EU	AM1(A),M112
				ADD 1 TO INV NO.
				MIME IN MASK
				END COST MI

Figure INVO-2

Invoicing--Basic Assembler

Address	Instruction	Operation	Comments
0346	47F0	0032	MC X'F',L83 GET ADD. CRD
034A	0205	0475 04C7	LM1 MVC NO1,AR3+1 STORE CONTENTS OF INDEX CARD
0350	0205	05F9 04CD	LB2 MVC NO15,AR3+7 CUSTOMER NO
0356	0201	0607 0403	MVC NO16,AR3+13 ORDER NO
035C	0201	060A 0405	MVC NO17,AR3+15 MONTH
0362	0201	0600 0407	MVC NO18,AR3+17 DAY
036A	0201	0407 0409	MVC AR3+17(2),AR3+19 YEAR
036E	9H20	0003	CIU 3,X'20 CLEAR AREA
0372	07F9		BCR X'F',9 S/S 3
0374	0511	04CD 06FB	LB4 CLC AR3+7(18),AR8 PRINT ORDER CARD
037A	4770	02C4	BC 7,PCMF CK NO PUNCHES
037E	F253	05AA 0506	PACK NO7,AR4+47(4) YES-PUNCHES
0384	F224	05C3 050F	PACK NO6,AR4+56(5) QUANTITY
038A	FC52	058R 05C3	MP NO7,NO6 UNIT PRICE
0390	FA55	0585 058B	AP NO5(6),NO7 GROSS AMT
0396	0209	0517 061C	MVC AR4+64(10),MSK3 ADD TOTAL GROSS
039C	DE09	0517 058C	ED AR4+64(10),NO7+1 PRINT ITEM CARD
03A2	0249	0627 0407	MVC PRIN,AR4 EDIT GROSS
03A8	4080	03A8	BAS 8,PUT MOVE TO PRINT
03AC	F271	0591 05CC	PACK NO8,NO10 PRINT
03B2	FC75	0591 058B	MP NO8,NO7 NET AMOUNT
03B8	9A0F	0597	OI NO8+6,X'0F MULT % X GROSS
03BC	FB55	058B 0592	SP NO7,NO8+1(6) ROUND OFF NET
03C2	0209	050E 061C	MVC AR4+55(10),MSK3 SUB. DIS-GROSS
03C8	DE09	050E 058C	ED AR4+55(10),NO7+1 PUNCH ITEM CARD
03CE	022A	04DF 04DE	MVC AR4+8(43),AR4+7 EDIT NET
03D4	0218	068F 0508	MVC AR6,AR10 BLANK AREA
03DA	0024	068F 0019	XIO AR6(X'24),25 MOVE NET+GROSS
03E0	4740	0284	HC 4,REPT PUNCH ITEM CARD
03E4	9822	0002	CIU 2,X'22 PUNCH BUSY
03E8	9823	0010	CIU 16,X'23 SEL. STAK 2
03EC	4740	0292	BC 4,CONT SEL PRT HD 2
03F0	0020	06C6 0012	SC 4,AGAI PRT HD BUSY
03F6	4740	029A	MVC AR6+7(X'20),18 PRINT ON CARD
03FA	0211	050F 050E	TIUB UVFL,X'43 PRINT HEAD BUSY
0400	9A43	0374	BC X'F',LB8 BLANK AREA
0404	47F0	008A	OVERFLOW CONDITION TEST CH 12
0408	9200	00C5	LB11 MVI LB8+11,X'00 RESET SWITCH
040C	0505	0475 04C7	CLC NO1,AR3+1 NEW CUSTOMER
0412	4770	008A	BC X'7',LB12 YES
0416	47F0	00A6	HC X'F',LB7 ORDER CARD
041A	95F3	04C6	PCMF CLI AR3,X'F3 IS IT ITEM CARD
041E	4770	0000	BC 7,LR6+8 NO IT IS NOT
0422	9200	0310	MVI FUF1+1,X'00 SET RFORM INDIC
0426	9240	0407	MVI AR4,X'40 BLANK CARD CODE
042A	0230	040R 0407	MVC AR4+1(49),AR4 HLANKS AREA
0430	0217	0509 0503	MVC AR4+50(24),NO13 PRINT ERROR TEXT
0436	0249	0627 0407	MVC PRIN,AR4 MOVE TO PRINT
043C	9A43	02EE	TIUB MUP,X'43 TEST CH 12
0440	47F0	02F2	BC 15,CHNG NOT CH 12
0444	4080	030A	MUP BAS 8,SKIP YES GO TO CH 1

Figure INVO-3

Invoicing--Basic Assembler

044A	4040 0329	CMNG	MAS	H,PUT	PRINT PRINM
044C	0217 0504 0504		MVC	AR4+50(24),AR4+49	HLA,KS,ARFA
044E	4420 0002	L49	CIU	2,X'20'	SPL STAK 2
044A	4040 037C		MAS	R,GET	READ PWF CARD
044A	0210 04CA 0471		MVC	AR3,AR5	MOVE TO WK ARFA
0440	0505 0475 04C7		CLC	HL,AR3+1	NEW CUSTIDMFK
044A	4740 02FC		MC	R,LM4	YES
044B	42F0 0310		MVI	FILE+1,X'F0'	RESET RFORM ERR
046E	47F0 002E		MC	X'F',LM3-4	START NEW UPER
0472	47F0 00EE		MC	X'F',LB5	NO RFORM ERROR
047A	47F0 0032		BC	X'F',LM3	
047A	0214 0505 0545		MVC	AR4+44(24),M04	PRINT ERROR TEXT
0440	0201 04ED 04EC		MVC	AR4+22(2),AR4+21	OUT OF SEQUENCE
044B	0249 0427 0407		MVC	PRIN,AR4	MOVE TO PRINT
044C	4040 03AA		BAS	R,PUT	PRINT
0440	9900 0066		HPR	100.0	HALT
0494	47F0 033A		MC	X'F',0-4	END-STOP
049A	0206 04FA 0470		MVC	AR4+35(7),M02	A NEW INVOICE NO
049E	9208 0415		MVC	PUCH+5,X'0A'	RESET SWITCH
04A2	9200 0431		MVI	PUER-3,X'00'	RESET STA SEL
044B	92F0 04F9		MVI	AR4+34,X'F0'	CARD INDICATOR
04AA	0207 0404 04F9		MVC	AR7(H),AR9	INVOICE NO.
0480	4040 0410		BAS	R,PUCH	PUNCH
0484	9423 0020		CIU	32,X'23'	USE PRINT HEAD 1
048A	0020 0608 0008		XIU	AR7(X'20'),8	PRINT CARD
048E	4740 0362		BC	4,RPET	PRINT HEAD BUSY
04C2	9900 0FFF		HPR	X'FFF',0	HALT
04CA	47F0 036C		BC	X'F',0-4	END
04CA	92F0 00C5		MVI	LB8+11,X'F0'	SET UV-FLUM SN
04CE	47F0 02AE		BC	15,LB11-4	
04D2	0022 0471 004E		XIU	AR5(X'22'),78	
04D4	9A24 0342		TIUB	-EUIK,X'24'	
04DC	4740 0396		BC	8,BSYR	
04E0	4740 037C		MC	4,GET	
04E4	9900 0001		HPR	X'001',0	
04EA	47F0 037C		BC	15,GET	
04EC	9A20 0396		TIUB	*,X'20'	
04FO	9A21 03A0		TIUB	EKKK,X'21'	
04F4	07FA		BCK	15,H	
04F6	9900 0011		HPR	X'011',0	
04FA	47F0 037C		BC	15,GET	
04FE	0040 0627 004A		XIU	PRIN(X'40'),74	
0504	4740 03AE		BC	8,MSYP	
0504	4740 03A9		MC	4,PUT	
050C	9900 0002		HPR	X'002',0	
0510	47F0 03AA		BC	15,PUT	
0514	9A40 03BE		TIUB	*,X'40'	
051A	9A41 03CA		TIUB	ERKP,X'41'	
051C	47F0 0302		BC	15,EHAS	

Figure INVO-04

Invoicing--Basic Assembler

0520	9400 0027	ERKP	MPC	X'022',0	
0524	47F0 03AM		MC	15,PUT	
0524	0244 0627 0A20	ERAS	MVC	PRIN(74),PRIN-1	
052E	07FM		BCK	15,R	
0530	9A45 0001		SPACE	MINUTES	
0530	07RM	SKIP	CIU	1,X'45'	
053A	4740 031A		BCK	R,R	
053A	9900 0007		RC	4,SKIP	
053E	47F0 030A		MPC	X'007',0	
0542	9A44 0002		BC	15,SKIP	
054A	07RM	SPA3	CIU	2,X'44'	
054A	4740 03EC		MPC	R,R	
054C	9900 0003		BC	4,SPA3	
0550	47F0 03EC		MPC	X'003',0	
0554	9A44 0001		BC	15,SPA3	
055A	07RM	SPA2	CIU	1,X'44'	
055A	4740 03FE		BCK	R,R	
055E	9900 0004		RC	4,SPA2	
0562	47F0 03FE		MPC	X'004',0	
			MC	15,SPA2	
0566	0027 060A 0023		PUNCH	ROUTINE	
056C	4780 0426		XIO	AK7(X'27'),35	
0570	4740 0410		BC	8,PUNY	
0574	9900 0005		RC	4,PUCH	
057A	47F0 0410		MPC	X'005',0	
057C	9A20 0426		BC	15,PUCH	
0580	9A21 0434	PUNY	TIUB	8,X'20'	
0584	9B22 0004		TIUB	PUER,X'21'	
058A	07FM		CIU	4,X'22'	
058A	9900 0006		BCR	15,R	
05AE	47F0 0410	PUER	MPC	X'006',0	
			RC	15,PUCH	
0592	9B23 0020		WRITE	ROUTINE	
059A	4780 044C	HEAD	CIU	32,X'23'	
059A	4740 043C		BC	8,WRCO	
059E	9900 0077		RC	4,HEAD	
05A2	0020 0609 0022		MPC	X'077',0	
05A8	4780 0462	WRCD	XIU	AK7+1(X'20'),34	
05AC	4740 044C		BC	8,WRSY	
0580	9900 0007		BC	4,WRCO	
0584	47F0 044C		MPC	X'007',0	
058A	9A22 0462		RC	15,WRCO	
058C	9A25 046C	WRSY	TIUM	8,X'22'	
05C0	07FM		TIUM	ERWA,X'25'	
05C2	9900 0004		BCR	15,R	
05C6	47F0 044C	ERWA	MPC	X'004',0	
			BC	15,WRCO	
05CA	F4		DEFINITION	OF AREAS AND CONSTANTS	
05CB		NU	UC	X'F4'	INDICATOR INV
0501	40	NU1	US	CL6	CUST.NU.
0502			UC	X'40'	
0503		AR1	OS	CL1	WORK AREA
05DA		NU2	US	CL7	INV. NO
050C			OS	CL2	
		NU3	OS	CL6	DATE OF INVOICE

Figure INVO-05

Invoicing--Basic Assembler

0152		AN7	US	CL5P		WINK AREA
0153		AN3	US	CL17		WINK AREA
0154		AN4	US	CL74		WINK AREA
0155		AN9	EUU	AN4+34		
0156		AN10	EUU	AN4+44		
0157		AN11	EUU	AN7+1		
0158		AN12	EUU	AN7+4		
0159		AN13	EUU	AN2+14		
0160		CPL1	US	CL7		CHU. PMT AREA L1
0161			UNG	**29		
0162	C1C4 C4D9 C5E2 E240 C3C1 D9C4 40D6 E4E3	NU4	DC	C ADDRESS CARD UNIT		
0163	40D6 C640 E2C5 D1E4 C5D5 C3C5		DC	C UP SEQUENCE		
0164		CPL2	US	**30		CHU PMT AREA L2
0165			UNG			
0166	0000 0000 0000	NU5	DC	XLA**		GROSS AMOUNT
0167		NU7	US	CL6		MULTIPLICATION
0168		MSK1	DC	CL6		MULTIPLICATION
0169	4020 2020 2021 2020	CPL4	US	CL34		CHU PMT AREA L4
0170		NU6	US	CL3		UNIT PRICE
0171		NU9	US	CL4		MADE UP PAYMENT
0172		NU10	DC	CL2		DIGSUNT
0173	00	NU11	DC	X'00'		INDICATOR
0174		NU12	US	CL4		INVOICE NO
0175		NO13	DC	C PUNCH ERROR IN		
0176	017E4 05C3 C440 C5D9 D90A D940 C9D5 40		DC	C ITEM CARD		
0177	C9E3 C5D4 40C3 C1D9 C4	NU14	DC	C YOUR ORDER NO.		
0178	ER06 E4D4 60D6 D9C4 C5D9 40D5 D640	NU15	US	CL4		ORDER NO.
0179			DC	C DATED		
0180	4040 C1C1 E3C5 C440	NU16	DC	CL2		MONTH
0181	01		DC	C'/'		
0182		NU17	US	CL2		DAY
0183			DC	C'/'		
0184		NU18	DC	CL2		YEAR
0185	1F	NU19	DC	X'1F'		PACKED DECIMAL 1
0186	4020 2020 2020 2021 2020 2020	MSK2	DC	X'40202020202021202020'		
0187	4020 2020 2021 2020 2020	MSK3	DC	X'402020202021202020'		
0188	40		DC	X'40'		
0189		PRIN	US	CL74		
0190		AR5	US	CL78		
0191		AR6	US	CL25		
0192		AR7	US	CL35		
0193		ARA	US	CL1A		
0194	4040 4040 4040 D40A C4C5 40D6 CA40 D7C1	NU0E	DC	C MADE UP PA		
0195	ER06 C5D5 E340 4040 40C4 C1E4 E240 C1E3		DC	C YMENT DAYS AT		
0196	4040 4040 6CC4 C9E2 C3D6 E4D5 E340 D4D2		DC	C DISCOUNT UK		
0197	4040 40C4 C1E4 E240 D5C5 E3		DC	C DAYS NET		
0198		END		INVO		

Figure INVO-06

Invoicing Output

M. R. MURLEY
 125A I ST. 542375 1111115 042371
 BEND, OREGON
 YOUR ORDER NO. 325454 DATED 05/12/71
 532853 BOTTLE WATER 12 00125 1500
 YOUR ORDER NO. 532875 DATED 05/28/71
 5352756 BALL BEARINGS 24 00523 12552
 14052
 03 0421
 13631
 MODE OF PAYMENT 20 DAYS AT 03% DISCOUNT OR 30 DAYS NET

Figure INVO-7

RPG DISK OPERATION

This section will aid in the preparation of the work required for the Swanson study, since a disk operation or sorts is assumed. The main topics are presented on the next page, and are presented as File Description Specifications. The two main divisions of the section are:

- a. Sequential Disk Files
- b. Indexed Sequential Files

SEQUENTIAL DISK FILES

Pre-sorted records are read from another external file (or disk) and written in sequence within the extents on the disk as fixed length records, either blocked or unblocked. The Job Control program is presented the actual location of the file on the disk at the time the file

File Description Specifications

Line	Filename	File Type	File Organization	File Format	Length of Key Field or of Base of Address Field	Base of Address Field	Number of Records	Number of Bytes	Number of Bytes	File Address
01	CARDIN	IP	F	80	80					
02	SEADSK	OS	F	279	90					
<p>RECFM: DISK:FBYSP2NS</p> <p>XTENT CONTROL STATEMENT</p> <p>DISK REQUIRES STANDARD LABELS</p> <p>DEVICE IS 2311 DISK DRIVE</p> <p>DISK FILE FORMAT IS FIXED</p>										

Input Specifications

Line	Filename	Record Identification Codes				Field Location		Field Name	Control Key at 116	Starting Position
		Position	Length	Position	Length	From	To			
01	CARDIN	011	01	1	01					
02		027	02	1	02	11	40	A		
03		031	03	1	03	11	50	B		
04						11	73	C		
05						41	SEVANT			

Figure S-1. RPG Coding
Sequential File - Loading the File, Part 1 of 2

This RPG program extracts fields from types of cards and writes a sequential field of ninety character records per block. A numeric field is packed on a disk that was unpacked on cards. A program constant "M" is placed in the first byte of a record. The "M" is a constant which is not contained in data cards.

Output-Format Specifications

Line	Filename	File No.	Output Indicators		Field Name	Field Length	Field Position	Content or Edit Word	Starting Position
			Start	End					
00	SEQDSK	0		03					
01					A	31			
02					B	81			
03					C	86			
04					AMT	90P			

Figure S-1 (continued). RPG Coding Sequential File - Loading the File, Part 2 of 1

SEQUENTIAL FILE - SEQUENTIAL RETRIEVAL

File Description Specifications

Line	Filename	File Type	File Description	File Length	File Position	File Organization	File Processing	File Address
00	SEQDSK	1PE	F 216 945	(55)				(2)
01						SEQUENTIAL ORGANIZATION	ALWAYS BLANK FOR SEQUENTIAL PROCESSING OF SEQUENTIAL FILE.	
02						NOT APPLICABLE		
03						SEQUENTIAL PROCESSING		

Figure S-2 Sequential File - Sequential Retrieval

Sequential Retrieval is accomplished by letting column 28, column 31 and column 32 remain blank, but the number of extents as indicated by column 68-69 must agree with the number assigned when the file was written.

SEQUENTIAL FILE - SEQUENTIAL RETRIEVAL
WITH MATCHING RECORD LOGIC

File Description Specifications

File	Name	File Type		File Description		File Location		File Attributes		File Control	
		File Type	File Description	File Location	File Attributes	File Control	File Control	File Control	File Control		
00	SEODSK	01	01	01	01	01	01	01	01	01	01
00	CARDIN	02	02	02	02	02	02	02	02	02	02

Input Specifications

File	Name	Record Identification Codes			Field Location		Field Name	Field Indicator	Sorting
		Position	Position	Position	From	To			
00	SEODSK	AA	01	CS	2	8	CUST01	M1	SELECTIVE PROCESS BY MATCHING CAR AGAINST DISK FILE MATCHING FIELDS
00					7	9	DATAP		
00	CARDIN	BB	02	CS	2	8	CUST01	M1	
00					7	9	DATAP		

Figure S-3. RPG Coding Sequential File
Sequential Retrieval with Matching Record Logic

This permits selective processing of from two through nine input files. The files must be in the same collating sequence, but may be on any input device. The matching fields are assigned on the Input Specifications in columns 61 and 62 as matching fields M1 through M9. The files are processed first by selecting them in the order they are entered on the Input Specifications. The internal matching record indicator controls the processing of matching fields. The MR is tested on the calculation specification and processing is modified by its setting.

SEQUENTIAL FILE - RANDOM RETRIEVAL
WITH ADDRROUT OPTION

File Description Specifications

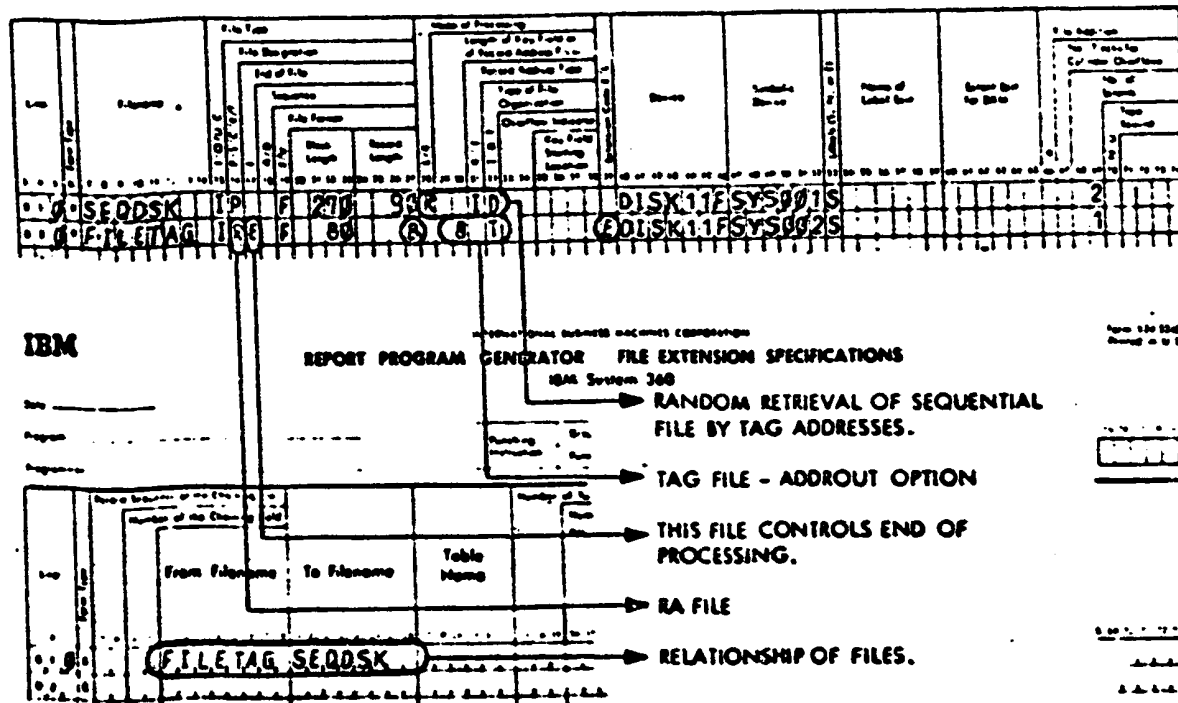


Figure S-4. Sequential File
Random Retrieval with ADDRROUT Option

SEQUENTIAL FILE - UPDATING THE FILE

File Description Specifications

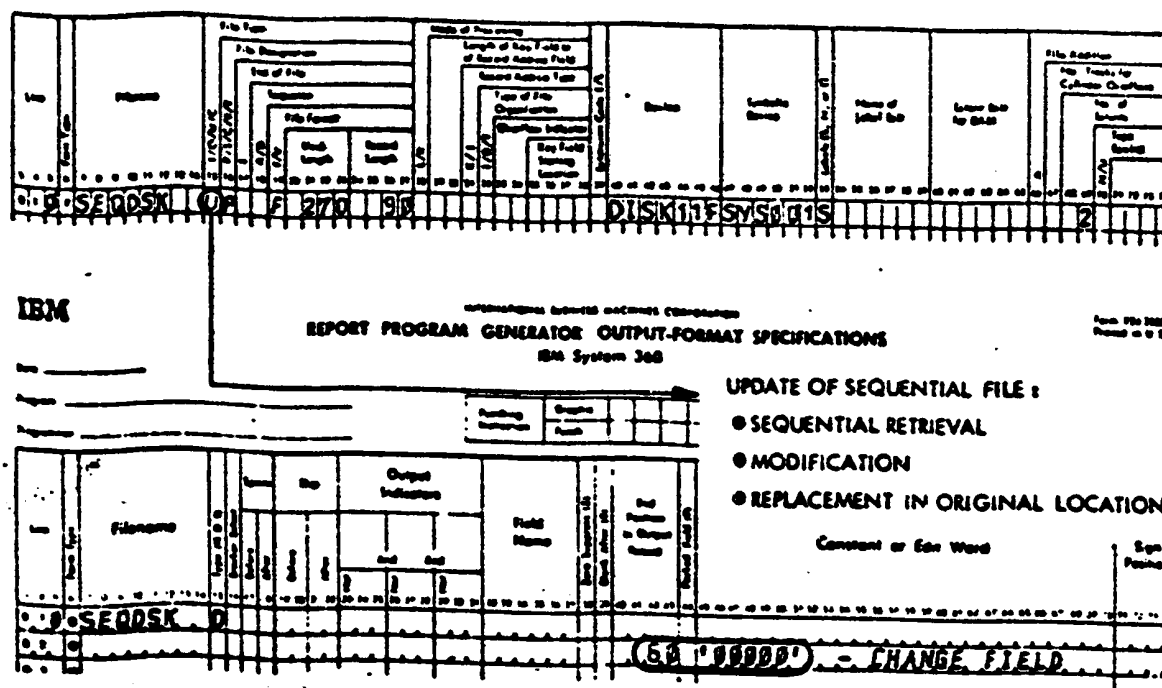


Figure S-5

Sequential File - Updating the File

It is not necessary to require the whole file to be updated if the record lengths are not modified, however if the length is to be modified, the entire file must be copied to an output-file.

SEQUENTIAL FILE - UPDATING AND ADDING RECORDS TO THE FILE (2 pages)

File Description Specifications

Line	Filename	Record Identification Codes	Field Location		Field Name	Field Indicators			Starting Page Position
			From	To		Pos	Len	Code	
01	SEQDSK AA	07 1 CS	1	90	RECORD				
02			2	6	CUST01L1M1				
03			10	14	UPDATE				
04			20	25	SALES				
05	CARDIN BB	02 1 C1							
06	OR	03 1 C2							
07			2	6	CUST02L1M1				
08			7	11	DATE				
09			12	17	ISSUE				
10			18	80	DATA				

Input Specifications

Line	Filename	File Type	File Organization	Length of Key Field or Record Address Field	Record Address Field	Start	Number of Records	Name of Data File	Length of Data File	Number of Columns
01	SEQDSK I	AF	270	90				DISK11FSYS01S		2
02	CARDIN IS	AF	80	80				READ01		
03	SEQDSK1 O	F	270	90				DISK11FSYS02S		2

THE FILE SEQDSK IS WRITTEN AS A NEW FILE SEQDSK 1.
RECORDS IN THE CARDIN FILE ARE OF 2 TYPES:
● 1 IN COL.1 - UPDATE RECORD
● 2 IN COL.1 - ADDITION
PROCESSING IS CONTROLLED BY MATCHING RECORDS

Figure S-6. Sequential File

Updating and Adding Records to the File, Part 1 of 2

Updating and Adding Records to the File (con't.)

Calculation Specifications

Line	Seq	In-File	In-File		Factor 1	Operation	Factor 2	Result Field	Field Length	Resulting Indicators			Comments
			Start	End						Rel	Cond	Code	
01		MR 02			SALES	ADD	ISSUE	SALES					
02		MR 02				MOVE	DATE	UPDATE					

Output-Format Specifications

Line	Seq	Filename	Seq	Seq	Output Indicators		Field Name	Field Length	Constant or Edit Word	Working Seq Position
					Start	End				
01		SEQDSK1			01	NMR				
02					03	RECORD	90			
03						CUST01	6	'S'		
04						DATE	11			
05						ISSUE	17			
06						DATA	80			
07							90	#####		
08					11	MR				
09						RECORD	90			
10						UPDATE	16			
11						SALES	25			

Figure S-7. RPG Coding. Sequential File

Updating and Adding Records to the File. Part 2 of 2

The File must be rewritten, but records can be updated during the same run. The files are specified as input and output, not as update, since update files must not include additions and deletions.

II INDEXED SEQUENTIAL FILE

File Description Specifications

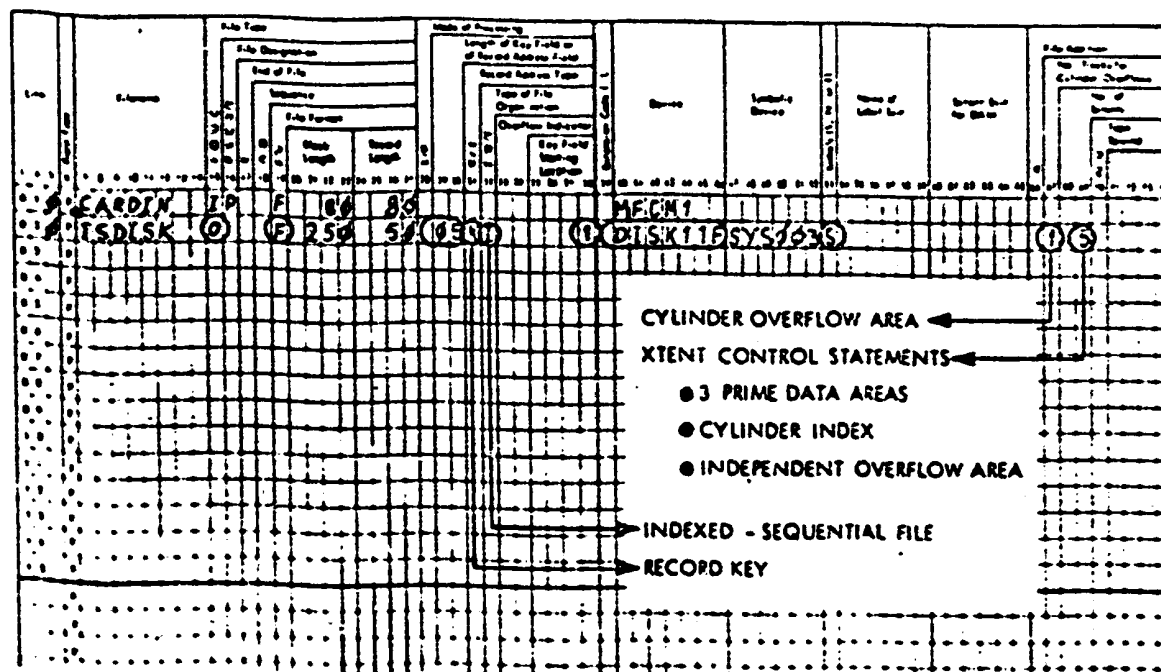


Figure S-8. RPG Coding

Indexed-Sequential File - Loading the File

Records are fixed length blocked or unblocked. As each file is written, indexes are created to provide a reference to records on each track or cylinder. The file is an output file designated by a 0 in column 15. There is a K in column 31 and an I in column 32. Columns 29-30 and 35-38 specify the length and location of the key field. There are two additional extents for the prime data area and one extent for the independent overflow area entered in columns 68-69. The user provides information about these with an extent control statement to the Job Control program when the load operation is enacted.

INDEXED SEQUENTIAL FILE - SEQUENTIAL RETRIEVAL

File Description Specifications

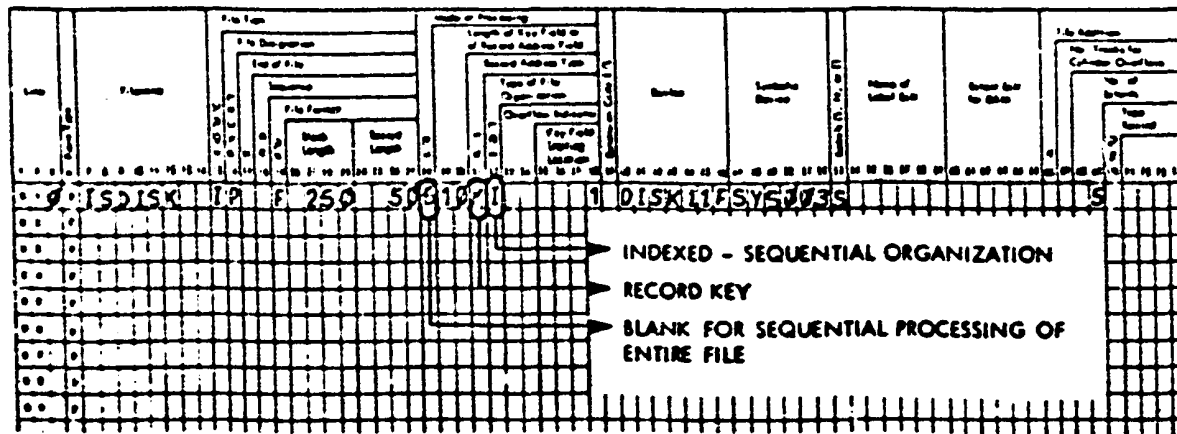


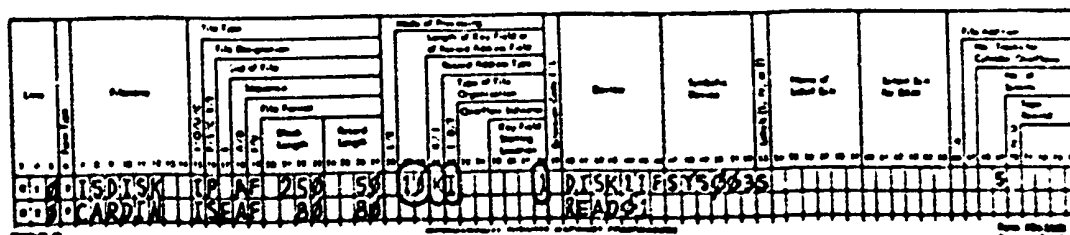
Figure S-9. RPG

Indexed-Sequential File - Sequential Retrieval

Retrieval is similar to retrieval of sequential files, and additions written in overflow areas are retrieved in sequential order. When multiple input files are specified, the order of processing is determined by the sequence they are entered on the Input Specifications. Coding follows the basic pattern for coding a disk file, plus a K in column 31 and an I in column 32. This indicates indexed-sequential organization. Column 28 is blank, and causes the entire file to be processed.

INDEXED SEQUENTIAL FILE - SEQUENTIAL RETRIEVAL
WITH MATCHING RECORD LOGIC

File Description Specifications



Input Specifications

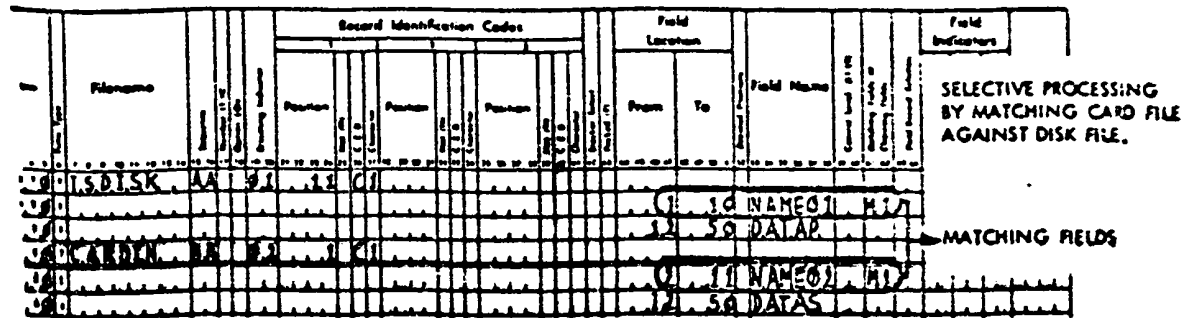


Figure S-10. RPG Coding. Indexed-Sequential File
Sequential Retrieval with Matching Record Logic

**INDEXED SEQUENTIAL FILE - SEQUENTIAL RETRIEVAL
WITH AN RA FILE OF LIMITS**

The RA (Record Address) is used to retrieve a segment of an indexed sequential file. The RA File defines the low and high limits of data to be processed. Two keys are contained in each record of the RA File (the lower limit and the upper limit). The record with a key equal to, or higher than, the lower limit is retrieved; all others above it follow in sequence up to, or equal to, the upper limit.

Only "keys" are contained in the RA file, and their purpose is to retrieve records. No reference is made to it on the Input Specifications so fields for update purposes can't be defined.

All information for the RA comes from the File Description and File Extension Specifications. An addition of a 1 in the "MODE OF PROCESSING" (column 28) of the File Extension Specifications tells that the RA File is related to the Indexed-Sequential File.

File Description Specifications

File	Name	File Type	File Description		Mode of Processing		Block	Volume	Length of File	Last Date	Last Time	File Extension	
			File Extension	File Name	Length of Key Field or of Record Address Field	Record Address Field							
000	RAFLMT	IR	F	89	80	10							
000	ISDISK	IP	F	253	56	10	1	DISK11FSYS003S					S
000	PRINTER	O						PRINTER					

File Extension Specifications

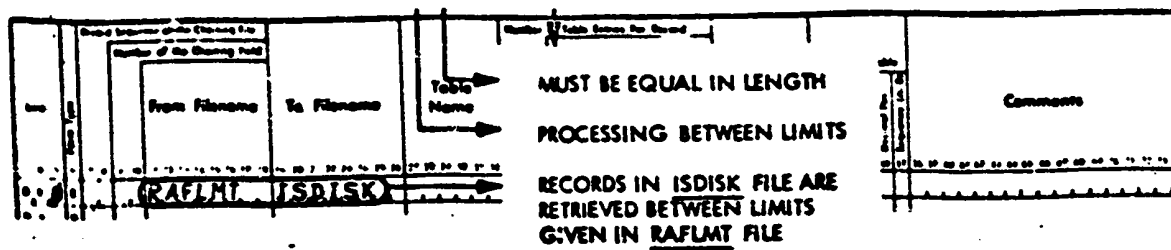


Figure S-11. Indexed-Sequential File Sequential Retrieval with RA File of Limits. Part 1

Input Specifications

File	Filename	Record Identification Codes			Field Location		Field Name	Field Indicators			Starting Sign Position
		Position	Position	Position	From	To		File	Block	Line	
000	ISDISK	AA	01	11	C1						
000						1	10	KEY			
000						12	50	DATA			

Figure S-11. Indexed-Sequential File Sequential with RA File of Limits. Part 2 of 2

INDEXED-SEQUENTIAL FILE - RANDOM
RETRIEVAL WITH CHAINING FILE

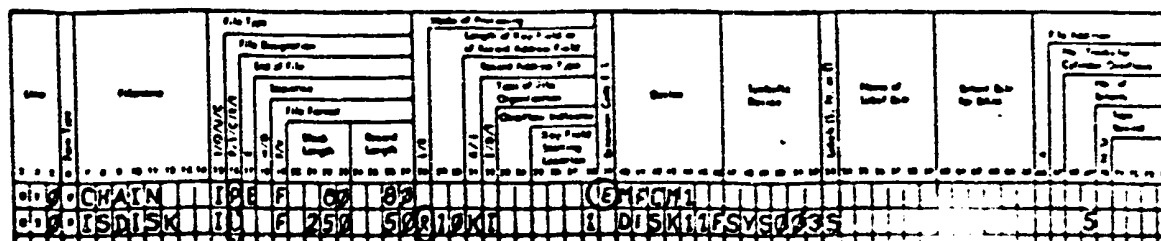
The "Chaining" technique is also used with the indexed sequential file. This method of random retrieval uses keys or chaining fields. They also serve as the link between two files. Each record in a chaining file contains the key of a record to be randomly retrieved from the chained file, which must be organized as an indexed-sequential file.

See Figure S-12. Chaining is encountered, in addition to the entries for the indexed sequential file:

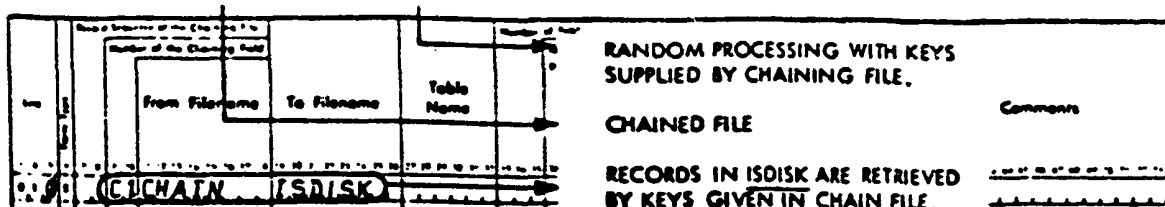
1. AC is included in the File Description (column 16) of the File Description.
2. An R is placed in Mode of Processing (column 28).
3. An I is placed in Extension Code (column 39), to reference the File Extension Specifications. It relates the file as the chaining file (columns 11 to 18) to the disk file, which is the chained file (columns 19-26) by the chaining field of C1 (columns 9-10).

Indexed-Sequential File
Random Retrieval with Chaining File (con't.)

File Description Specifications



File Extension Specifications



Input Specifications

Key or Chaining Field

Line	Filename	Record Identification				Codes				Field Location		Field Name	Control Field in Chaining File	Field Name in Chaining File	Field Indicators			Starting Sign Position
		Position	Length	Position	Length	Position	Length	Position	Length	From	To				Key	Chain	Field	
0.0	CHAIN	AA	01	1	CC													
0.0										11	11	CHAIN1	C1					
0.0										12	50	DATA						
0.0	ISDISK	AB	02	11	C1					1	10	KEY						
0.0										12	50	DATA						

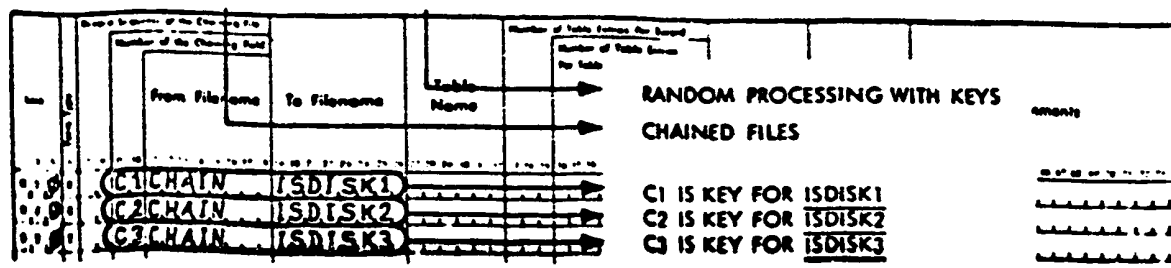
Figure S-12. Indexed-Sequential File Random Retrieval with Chaining File

INDEXED-SEQUENTIAL FILE - RANDOM RETRIEVAL WITH ONE CHAINING FILE FROM THREE CHAINED FILES

File Description Specifications

Line	Filename	File Type				Mode of Access				Name	Library	Name of Label	Number of Labels	File Address			
		File Description	File Format	File Organization	File Access Method	Length of Key Field or of Record Address Field	Record Address Type	Type of File Organization	Number of Records					File Address	File Address	File Address	
0.0	CHAIN	IPE	F	83	80					CHAIN1							
0.0	ISDISK1	IC	F	256	56	2KI				DISK11FSYS03S							5
0.0	ISDISK2	IC	F	256	10	2KI				DISK11FSYS04S							3
0.0	ISDISK3	IC	F	256	256	2KI				DISK11FSYS04S							2

File Extension Specifications



Input Specifications

Filename	Record Identification Codes				Field Location		Field Name	Field Indicators	Starting Position
	Position	Length	Code	Code	From	To			
CHAIN	AA	01	1	C1	2	11	CHAIN1	C1	
					12	15	CHAIN2	C2	
					16	19	CHAIN3	C3	
ISDISK1	AB	02	11	C1	18	80	DATA1		
					1	10	KEY1		
ISDISK2	AC	03	5	C2	12	50	DATA2		
					1	40	KEY2		
ISDISK3	AD	04	3	C3	6	100	DATA3		
					1	20	KEY3		
					4	25	DATA3		

Figure S-13. RPG Coding Indexed-Sequential File - Random Retrieval with One Chaining File from Three Chained Files

As depicted by Figure S-13, another method of using the chaining technique is to retrieve records from a chained file. Then use a field in those records to link to another file. The chained file then becomes a chaining file to the second chained file.

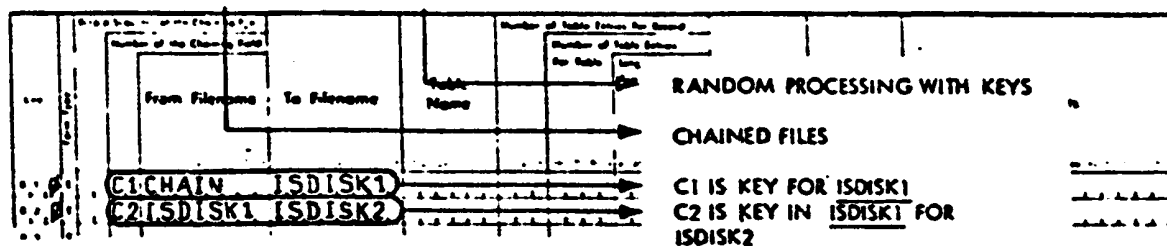
The chained file is unique type of RPG file. Records in this file are not retrieved until total time in the RPG processing cycle, while other files are read before total time. As a result, a chained file cannot contain control fields or matching fields. It is the only file which can be updated at total time.

INDEXED-SEQUENTIAL FILE - RANDOM RETRIEVAL
WITH CHAINED FILE AS A CHAINING FILE

File Description Specifications

File	File Name	File Type	File Organization		Record Length		Number of Records		File Status	File Class	File Code
			Block Length	Record Length	Number of Blocks	Number of Records					
0	CHAIN	TF	188	80							
0	ISDISK1	TC	250	50	10X1						S
0	ISDISK2	TC	500	100	4X1						3

File Extension Specifications



Input Specifications

File	File Name	Record	Record Identification Codes			Field Location		Field Name	Field Indicators		
			Position	Length	Code	From	To		Indicator	Indicator	Starting Position
0	CHAIN	AA	01	1	CC						
						2	11	CHAIN1			C1
						12	86	DATA1			
0	ISDISK1	AB	02	11	C1						
						1	10	KEY1			
						12	15	CHAIN2			C2
						16	50	DATA1			
0	ISDISK2	AC	03	5	C2						
						1	4	KEY2			
						6	100	DATA2			

Figure S-14. RPG Coding. Indexed-Sequential File Random Retrieval with Chained File as a Chaining File

ACTIVITIES

The final activities section presents a culmination of the book. The material was presented by models and methods for accomplishing these models.

If the reader will review Section I and the Preface, he will notice that this book could be used for a complete year.

This (year) can be easily accomplished by going back to the Swanson Study, and writing programs for the functional areas of the corporation.

It is suggested that the reader program for Financial Control first. This includes:

1. General Ledger and Budget Accounting

This activity reports all income, expenses and budget information to Management Planning.

2. Accounts Receivable

This activity accounts for all customer receivable information resulting from sales. Here customer records are maintained, reflecting credit status, statement of customer accounts, and customer receipts. The income is reported to General Ledger Accounting.

3. Cost Accounting

This activity accounts for and reports costs of production which includes material costs and labor costs.

4. Labor and Payroll

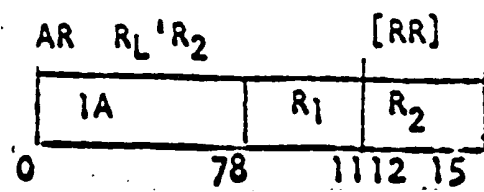
This activity accounts for payroll for all employees. A report is also made to General Ledger for payroll and

labor expenses. Cost Accounting also receives a report based on labor tickets from Production Planning and Control.

5. Accounts Payable

This activity accounts for, and reports payments for, goods and services received. Expenses are reported to General Ledger Accounting.

ADD REGISTER



The second operand is added to the first operand, and the sum is placed in the first operand.

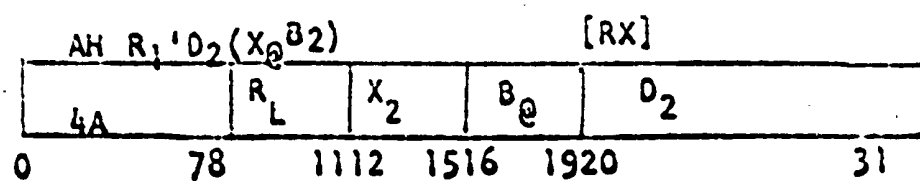
Both operands and the sum are 16-bit integers.

Condition Code:

- 0 Sum is zero
- 1 Sum is less than zero
- 2 Sum is greater than zero

AR10,11

ADD HALFWORD



The halfword second operand is added to the first operand, and the sum is placed in the first operand (register).

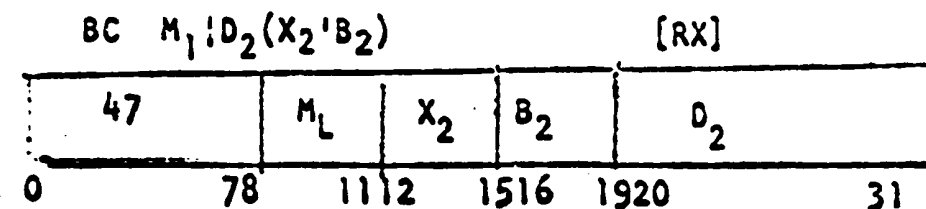
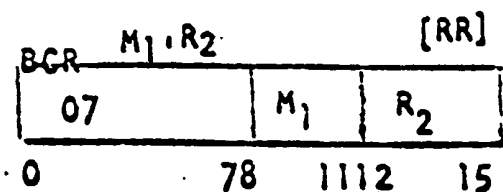
1. The second operand is a 16-bit signed integer.
2. The operand must be on a halfword integral boundary.
3. The first operand is a 16-bit signed integer.
4. The sum is a 16-bit signed integer.

Condition Code:

- 0 Sum is zero
- 1 Sum is less than zero
- 2 Sum is greater than zero

AH 8,HALF

BRANCH ON CONDITION



A branch to the address specified in the second operand is taken whenever the condition code matches a condition specified in the first operand (M).

To code this instruction:

1. Place the mask value corresponding to the desired condition 1.

Place the mask value corresponding to the desired condition code in the first operand.

Condition Code 0 1 2 3

Mask Value 8 4 2 1

Example:

A. Desired Condition Code is 1

B. Coding is BC4, TOP1

2. To test for more than one condition code, place the sum of the mask values corresponding to the desired condition codes in the first operand.

Example:

A. Desired condition codes are 0 and 2.

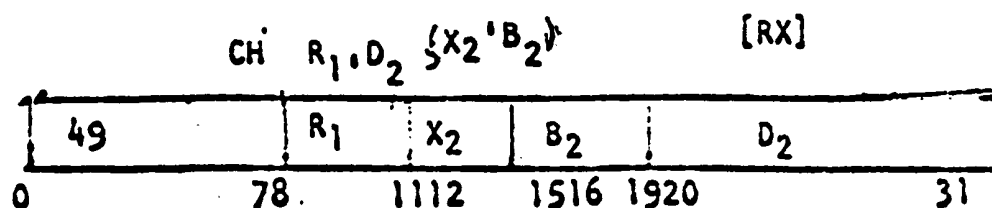
B. Coding is BC10, BRANCH

Note: Either condition code 0 or condition code 2 will cause a branch.

3. When the first operand is 15, the branch is always taken (unconditional branch).

4. When the first operand is 0, no branch is taken (a no-operation [no-op] equivalent).

COMPARE HALFWORD



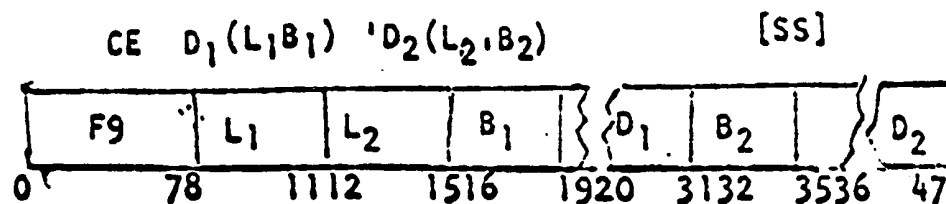
The first operand is compared algebraically with the halfword second operand, and the result determines the setting of the condition code.

1. Both operands are 16-bit signed integers.
2. The second operand must be on a halfword integral boundry.

Condition Code:

- 0 operands are equal
- 1 First operand is low
- 2 First operand is high
- 3 --

COMPARE DECIMAL



The first operand is compared algebraically with the second operand, and the result determines the setting of the condition code.

1. Both operands are in packed decimal format.
2. Comparison proceeds from right to left taking into account the sign as well as all the digits of each field.
3. The operand 2 field must not be longer than the operands 1 high-order zeros.
4. Plus zero and minus zero compare equal.

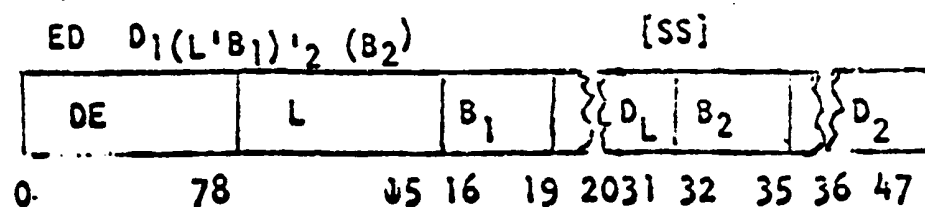
Condition Code:

- 0 Operands are equal
- 1 First operand is low
- 2 First operand is high
- 3 --

CP YTDF,AMT

BC 0,NOTX

EDIT



The format of the second (source) is changed from packed to zoned and is edited into the first operand.

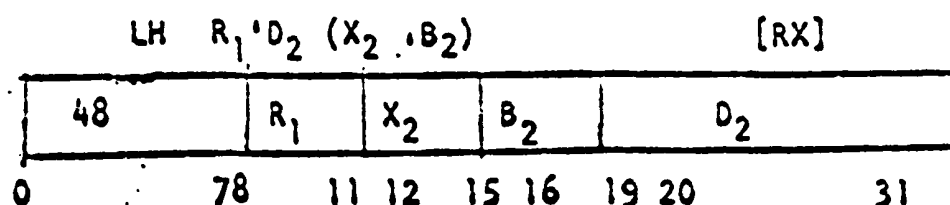
1. The second operand (source) must be in packed format.
2. Editing proceeds left to right, one character at a time.
3. The edited result replaces the pattern.

Condition Code:

- 0 Result is zero.
- 1 Result is less than zero.
- 2 Result is greater than zero.
- 3 --

ED OLPR-1(19), PRIN

LOAD HALFWORD



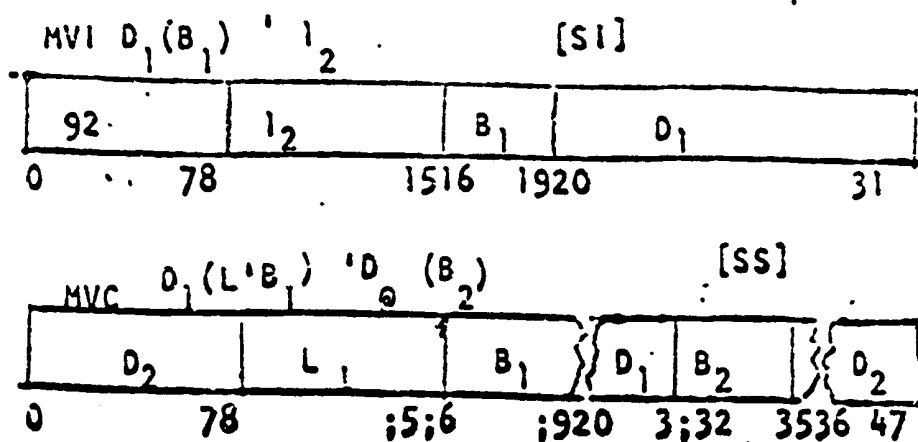
The halfword second operand must be on a halfword integral boundary.

Condition Code:

The code remains unchanged.

LH 11,HALF

MOVE



The second operand is placed in the first operand location.

MVC (Move Characters)

1. The bytes are moved one at a time in each field.
2. Movement is left to right.
3. The number of bytes is determined by the implicit or explicit length of the first operand.

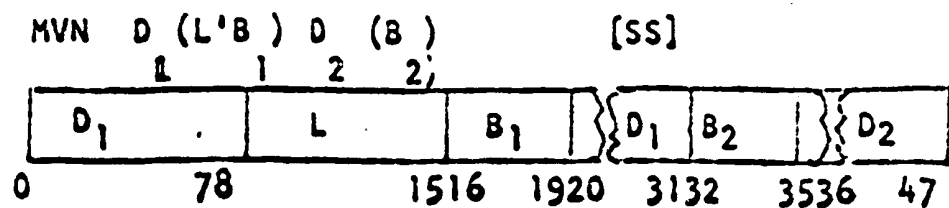
MVC (Move Immediate)

1. One byte of immediate data is stored in the first operand location. Immediate data supplied by the instruction itself, and in this case, is the second operand.

```
MVI SW,X"01"
```

```
MVI OLD,NEW
```

MOVE NUMERIC



The numeric portion (low-order four bits) of each byte in the second operand are placed in the numeric portion of the corresponding bytes in the first operand.

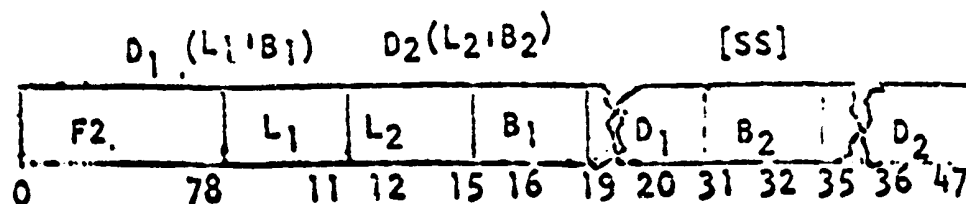
1. The number of bytes in the operation is determined by the implicit or explicit length of the first operand.
2. Movement is from left to right through each field.
3. Movement is one byte at a time.
4. Zones remain unchanged.

Condition Code:

The code remains unchanged.

MVN PINT+2(1),PINT+4

PACK



The signed or unsigned number in the zoned format, in the second location is changed to packed format and stored in the first operand location.

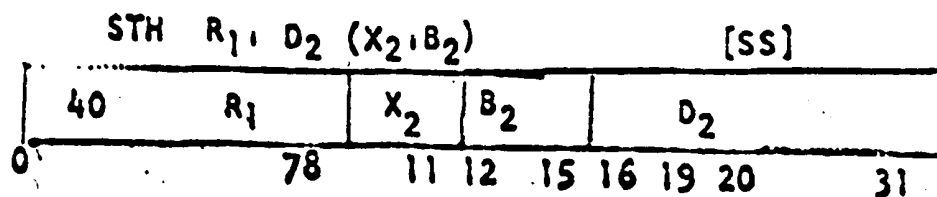
1. The fields are processed from right to left.
2. If the first operand field is too long, it will be filled with high-order zeros.
3. If the first operand field is too short, any remaining high-order digits in the second operand will be ignored.
4. The maximum size of the second operand (zoned field) is 16 bytes.

Condition Code:

The code remains unchanged.

PACK SGND,ZONE

STORE HALFWORD



The general register specified in the first operand is stored at the second operand halfword location.

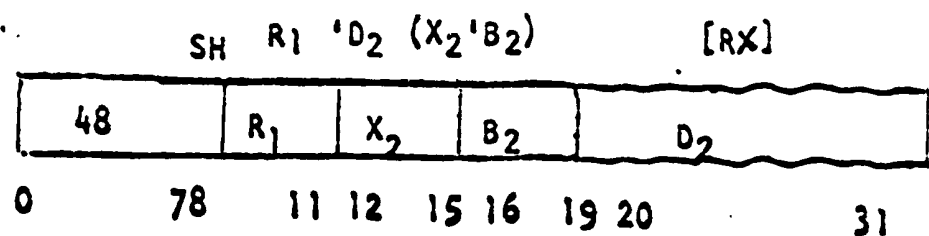
The second operand must be on a halfword integral boundry.

Condition Code:

The code remains unchanged.

STH 8,BUTE

SUBTRACT HALFWORD



The halfword second operand is subtracted from the register specified in the first operand, and the difference is placed in the register.

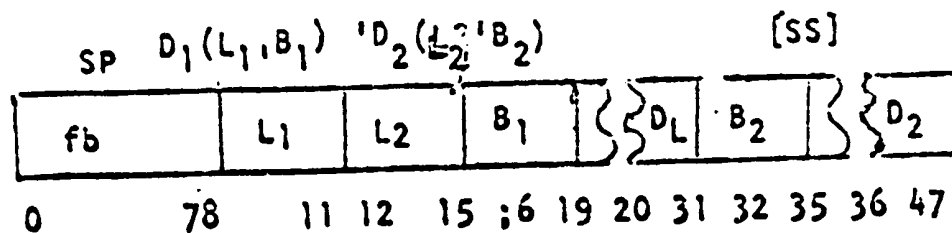
1. The second operand is a 16-bit signed integer.
2. The second operand must be on a halfword integral bound.
3. The first operand is a 16-bit signed integer.
4. The difference is a 16-bit signed integer.

Condition Code:

- 0 Difference is zero.
- 1 Difference is less than zero.
- 2 Difference is greater than zero.

SH 10,HALF

SUBTRACT DECIMAL



The second operand is subtracted from the first operand, and the difference is placed in the first operand.

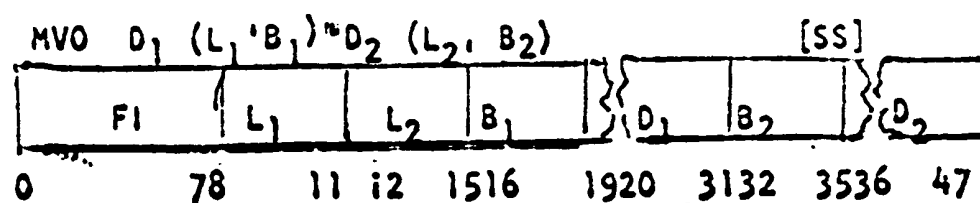
1. Both operands must be in packed format.
2. The difference is in packed format.
3. If the first operand is too short to contain the difference, overflow occurs, and the carry is lost.
4. If the second operand is shorter than the first, subtraction will take place normally.
5. A field may be subtracted from itself.
6. The second operand must not be longer than the first or a program error stop occurs.

Condition Code:

- 0 Difference is zero
- 1 Difference is less than zero.
- 2 Difference is greater than zero.
- 3 Overflow.

SP PAWT, PINT(3)

MOVE WITH OFFSET



The second operand is placed to the left of, and adjacent to, the low-order four bits of the first operand.

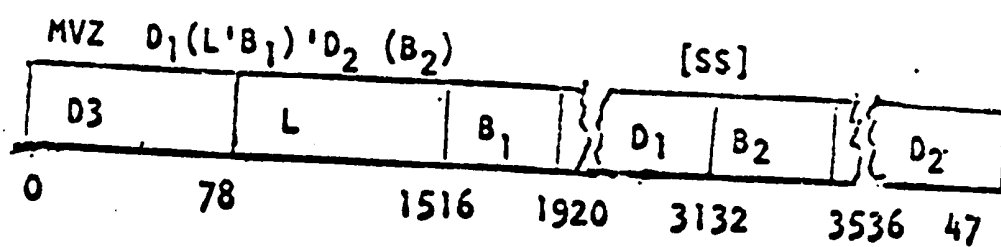
1. The fields are processed right to left.
2. If the second operand is shorter than the first operand, it is extended with high-order zeros.
3. If the first operand field is shorter than the second operand, the remaining information is ignored.

Condition Code:

The code remains unchanged.

MVO BDSN,GOSN(6)

MOVE ZONES



The high-order 4 bits (zone portion) of each byte of the second operand are placed in the first operand field.

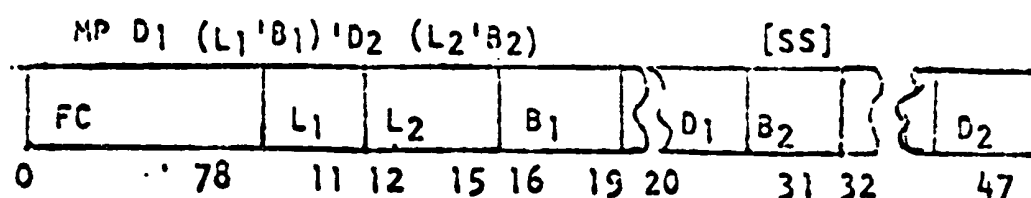
1. Movement is from left to right.
2. Movement is one byte at a time.
3. The number of zone portions moved is determined by the implicit or explicit length of the first operand.

Condition Code:

The code remains unchanged.

MVZ ZNUM+4(L),ZNUM+3

MULTIPLY DECIMAL



The first operand (multiplicand) is multiplied by the second operand (multiplier), and the product is placed in the first operand location.

1. Both the multiplicand and the multiplier must be packed.
2. The product is in packed format.
3. The length of the first operand in bytes must be equal to or greater than the number of bytes required to contain all of the multiplicand plus the total number of bytes in the multiplier (second operand) field.

Example:

Multiplier - XXXX 0X XX Xs (3 bytes)

Largest Multiplicand - XXXXXXXX XX XX XX XS (4 bytes)

The product field length must then be 7 bytes (3+4) or larger in length

XX XX XX XX XX XS

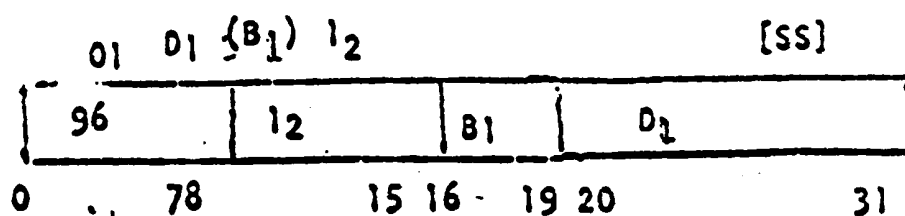
4. The multiplier may not exceed 15 digits and sign (8 bytes).
5. The maximum product size is 41 digits and sign (16 bytes).

Condition Code:

The code remains unchanged.

MP PINT, PRAT

OR IMMEDIATE



The first and second operands are examined on a corresponding bit by bit basis.

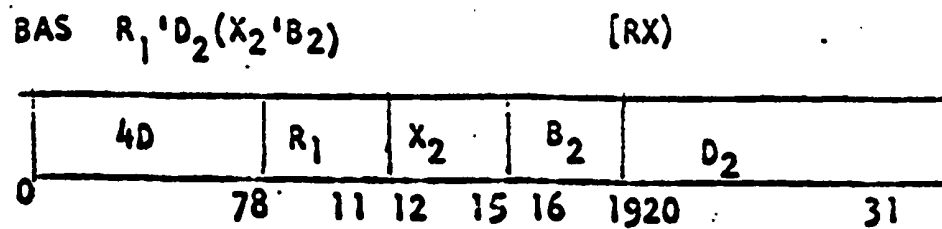
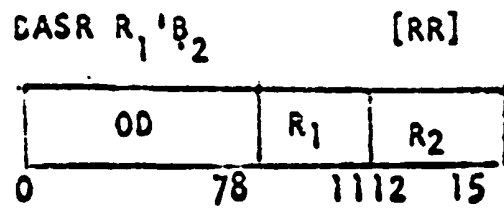
1. If either or both of the corresponding bits are ones, the result is a one and replaces the bit in the first operand.
2. If both bits are zeros, the result is zero and replaces the operand.
3. The second operand is one byte (8 bits) of immediate data which operates with one byte of data at the first operand storage location.

Condition Code:

- 0 Result is zero.
- 1 Result is not zero.

01 OPRD,X"08"

BRANCH AND STORE



The address of the next sequential instruction (nsi) is stored in the first operand, and a branch is taken to the address-stored in the second operand.

BASR (Branch and Store Registers)

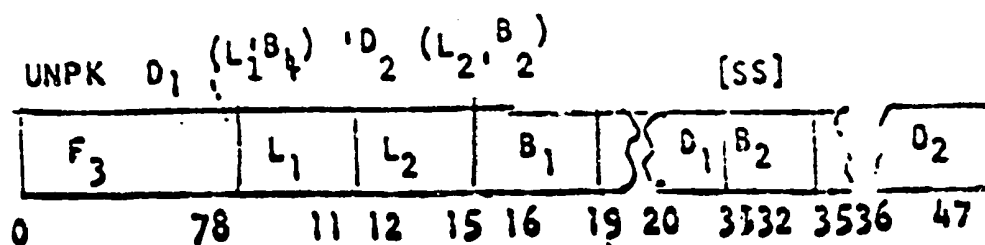
1. If the second operand is 0 (zero), no branch is taken.

Condition Code:

The code remains unchanged.

BEGIN BASR 11.0

UNPACK



The packed format second operand location is changed to signed zoned format and is placed in the first operand location.

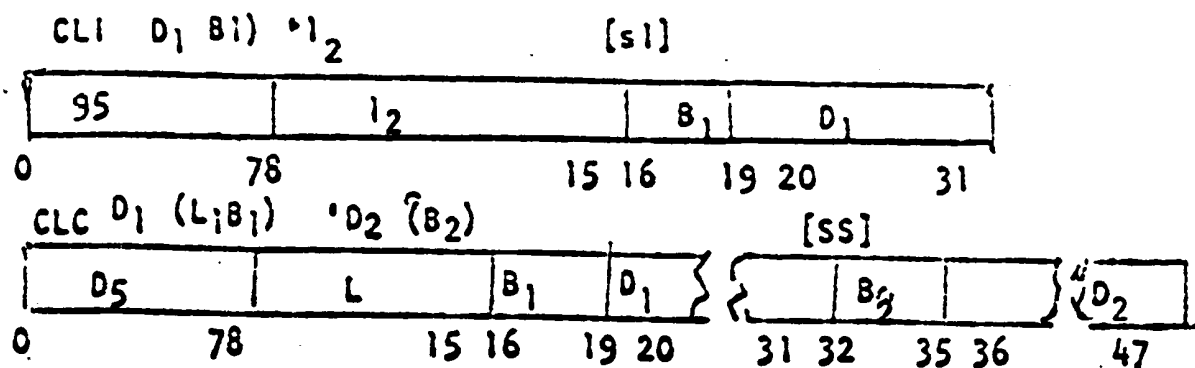
1. The fields are processed from right to left.
2. If the first operand field is too long, it will be filled with high-order zeros.
3. If the first operand field is too short, any remaining high-order digits will be ignored.
4. The maximum size of the first operand (zoned field) is 16 bytes.

Condition Code:

The code remains unchanged.

UNPK ZONE, PACK

COMPARE LOGICAL



The first operand is compared with the second operand, and the result is indicated in the condition code.

1. Comparison is binary (that is, bit by bit).
2. Comparison proceeds from left to right.
3. Comparison ends as soon as an inequality is found.

CLI only:

One byte in the storage location specified by the first operand is compared with one byte of immediate data.

CLC only:

The number of bytes to be compared is specified by the implicit or explicit length of the first operand.

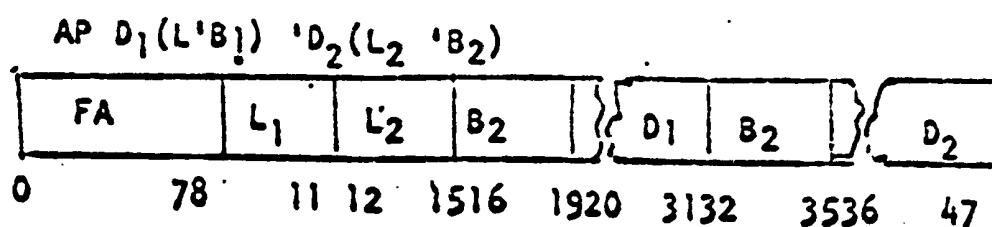
Condition Code:

- 0 Operands are equal.
- 1 First operand is low.
- 2 First operand is high.
- 3 —.

CLI CODE, C"E"

CLC NAME, C"SMITH"

ADD DECIMAL



The second operand is added to the first operand, and the sum is placed in the first operand storage location.

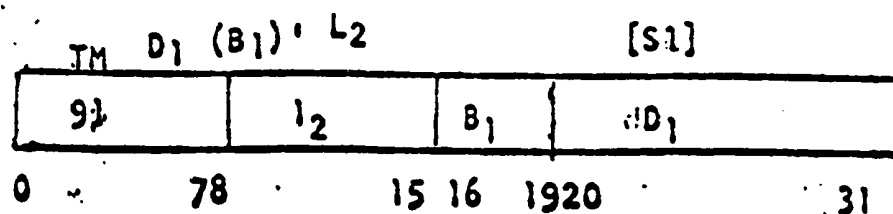
1. Both operands must be in packed format.
2. The sum is in packed format.
3. If the first operand is too short to contain the sum, overflow occurs and the carry is lost.
4. If the second operand is shorter than the first operand, addition will take place normally.
5. A field may be added to itself.
6. The second operand may not be longer than the first or an error stop occurs.

Condition Code:

- 0 Sum is zero.
- 1 Sum is less than zero.
- 2 Sum is greater than zero.
- 3 Overflow

AP PINT,INT

TEST UNDER MASK



The state of the first operand bits selected by a mask (second operand) is used to set the condition code.

1. 1-8 bits may be tested.
2. The mask is one byte (8 bits) of immediate data (second operand).
3. A mask bit of one indicates that the corresponding storage bit is to be tested.
4. A mask bit of zero indicates that the corresponding storage bit is to be ignored.

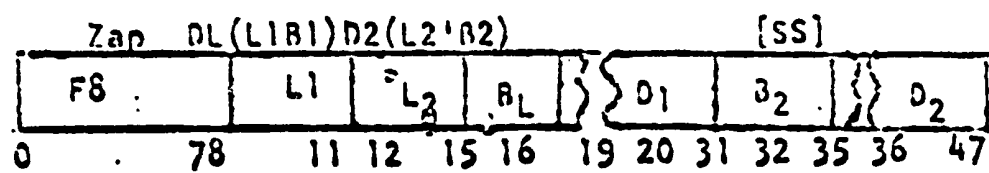
Condition Code:

- 0 Selected bits or mask, all-zero.
- 1 Selected bits mixed zero and one.
- 2 --
- 3 Selected bits all-one.

TM BUTE, X"FF"

BC 4, NWRT

ZERO AND ADD



The storage location specified by the first operand is cleared to zero, and then the second operand data (packed format) is added to the first operand.

1. The length of the second operand must not be greater than the first operand.

2. A negative zero balance will be made positive.

Condition Code:

0 Result is zero.

1 Result is less than zero.

2 Result is greater than zero.

GLOSSARY

Absolute Address: A pattern of characters that identifies a unique storage location or device without further modification. Synonymous with machine address.

Access Method: Any of the data management techniques available to the used for transferring data between main storage and an input/output device.

Accumulator: A register in which the result of an arithmetic or logic operator is formed.

Address: An identification, as represented by a name, or number for a register location in storage, or other data source or destination. Loosely, any part of an instruction which specifies the location of an operand for the instruction.

Address Calculation: A calculation performed on a key so that the result is an address.

Address Modification: The process of changing the address part of a machine instruction by means of coded instructions.

Address Register: A register that stores an address.

Alias: Another name by which a file is known.

Alphanumeric: A generic term for alphabetic letters, numerical digits and special characters.

Ascending Order: The word with the lowest key appears in the cell of the lowest number.

Assembler: A program that assembles. See assemble---A program which prepares an object language program by producing absolute or relocatable machine code from a source program of statements containing symbolic operation codes and symbolic operands.

Assemble: To prepare an object-language program from a symbolic language program by substituting machine operation codes for symbolic codes and absolute or relocatable addresses for symbolic addresses.

Auxiliary Function: A function that is either control or data oriented but does not cause data transmission.

Base Register: A register used for addressing purposes.

Basic Assembler Language: A symbolic language for the writing of source programs.

Basic Assembler Program: A program used to translate source programs written in basic assembler language into machine language.

Basic Monitor: The main control program of the DPS. Available in a card or disk version. Resident in core storage when control is required. Loads programs into core storage and causes their execution.

Bifurcating: No vortex is higher than three, for an absence only root-like, simple linking nodes, and one root are permitted.

Bifurcating Cell: Contains two pointers.

Binary: 1. Pertaining to a characteristic or property involving a selection, choice, or condition in which there are two possibilities.
2. Pertaining to the number representation system with a base of two.

Binary Code: A code that asks use of exactly two distinct characters, usually 0 and 1.

Binary-Coded Character: One element of a notation system for representing alphanumeric characters such as decimal digits, alphabetic letters, punctuation marks, etc., by a fixed number of consecutive binary digits.

Binary-Coded Decimal: Pertaining to a decimal notation in which the individual decimal digits are each represented by a binary code group: e.g., in the 8-4-2-1 coded decimal notation, the number twenty-three is represented by 0011, in binary notation, twenty-three is represented as 10111.

Binary Digit: A character used to represent one of the integers smaller than the Radix 2.

Bit: A binary digit.

Bit: The smallest unit of information in System/360. It can have either of the two binary values: zero or one.

Bit Extraction: Mapping by extracting specified bits from a key and then compressing.

Binary-to-Decimal Conversion: Conversion of a binary number to the equivalent decimal number; i.e., a base-two number to a base-ten number.

Blank: One of the characters in a character.

Blank Character: Any character or characters used to produce a character space on an output medium.

Blocks: 1. A physical set of records grouped for the purpose of conserving tape; or disk storage space; or increasing the efficiency of access or processing.

2. A group of bytes, transferred to or from a physical unit device in one operation, may contain one or more logical records.

Branch: To depart from the normal sequence of executing instructions in a computer. A machine instruction that can cause a departure as in (1). Synonymous with "transfer".

Buffer: Area in memory which holds data brought in from an input device.

Byte: A sequence of adjacent binary digits operated upon as a unit.

Byte: The basic unit of information in System/360. Every byte consists of eight bits each having a value of zero or one, (see bit).

Card Code: The combination of punched holes which represent characters (letters, digits, etc.) in a punched card.

Card Column: One of the vertical lines of punched positions on a punched card.

Card Field: A fixed number of consecutive card columns assigned to data of a specific nature.

Card Image: One-to-one representation of the contents on a punched card.

Card Punch: A device to record information in cards by punching holes in cards to represent letters, digits, and special characters.

Card Reader: A device which reads and translates into internal form the holes in punched cards.

Card-Resident-System: Consists of the card control programs, Basic Monitor, Job Control, and initial Program Loader. Used for the execution of object programs contained in punched cards.

Card Stacker: A mechanism which stacks cards in a pocket after they pass through a machine.

Cell: Continuous locations where a record is stored.

Character: One of a set of elementary symbols which may include decimal digits 0 through 9, the letters A through Z, punctuation marks, and any other symbols acceptable to a computer for reading, writing or storing.

Character Set: A list of characters acceptable for coding to a specific computer or input/output device.

Clear: To put a storage or memory device into a prescribed state, usually that denoting zero or blank.

Coded Decimal: A type of notation in which each decimal digit is identified by a group of binary ones and zeros.

Column Binary: Pertaining to the binary representation of data on punched cards in which adjacent positions in a column correspond to adjacent bits of data, e.g., each column in a 12 row card may be used to represent 12 consecutive bits of a 36 bit word.

Command: An instruction in machine language.

Communication: The process of transferring information from one point, person, or equipment to another.

Communication Region: An area of the Basic Monitor. Contains date, storage capacity specification, UPS1 byte, user areas 1 and 2, and program-name area. Provides for inter-program and intra-program communication.

Compiler: A compiler is a program which prepares an object language program.

Component: A basic part, an element.

Compound Key: Combining several fields of the record for either ordering or searching.

Computer: 1. A device capable of solving problems by accepting data, performing prescribed operations on the data, and supplying the results of these operations on the data. Various types of computers are calculators, digital computers and analog computers.

2. In information processing, usually an automatic stored-program computer.

Computer, Hexadecimal Number System: A number system using the equivalent of the decimal number sixteen as a base.

Computer Instructions: Same as machine instruction.

Constant: A fixed or invariable value or data item.

Control Field: A group of contiguous bytes that are within a data record. The sort, or merge of the records, is based on the collating sequence as applied to these bytes.

Control Statement: A punched card that contains information to tailor a program to the requirements of the user.

Convert: To change the representation of data from one form to another; e.g., to change numerical data from binary to decimal or from cards to tape.

- Core-Image Directory:** A table on the system disk pack containing the addresses and sizes of the program and/or program phases in the core-image library.
- Core Library:** A disk area containing the job control program, other IBM-supplied programs (except the Basic Monitor), and user's problem programs. Permits retrieval of programs and/or phases by the Basic Monitor.
- Core-Image Maintenance Program:** A service program. Updates the core-image library and directory. Is used to add and/or delete phases.
- Core Storage:** A form of high speed storage using magnetic cores.
- Counter:** A device such as a register or storage location used to represent the number of occurrences of an event.
- Cycle:** 1. An interval of space or time in which one set of events or phenomena is completed.
2. Any set of operations that is repeated regularly in the same sequence. The operations may be subject to variations of each repetition.
- Cylinder:** All the marks which can be accessed without moving the READ/WRITE Heads.
- Data:** Any representation, such as character quantities, to which meaning might be assigned.
- Data Conversion:** The process of changing data from one form of representation to another.
- Data File:** A collection of related records organized in a specific manner, for example, a payroll file (one record for each employee, showing his rate of pay, deductions, etc.) or an inventory file (one record for each inventory item, showing the cost, selling price, number in stock, etc.).
- Data Management System:** A data management system deals with records in a file. The records contain information called elementary data items, which are the object of processing.
- Data Processing:** A systematic sequence of operations performed on data.
- Data Processing System:** A network of matching components capable of accepting information, processing it according to a plan, and producing the desired results.
- Data Structure:** The arrangement and interrelation of records in a file form; a data structure.

Decimal: 1. Pertaining to a characteristic or property involving selection, choice or condition in which there are ten possibilities.

2. Pertaining to the number representation system with a radix of ten.

Deck: A collection of punched cards.

Decimal-to-Binary Conversion: The conversion of a decimal number to the equivalent binary number, i.e., a base-ten number to a base-two number.

Decision: A determination of future action.

Decision Instruction: An instruction that effects the selection of a branch of a program, e.g., a conditional branch instruction.

Decrement: The quantity by which a variable is decreased.

Decision Box: A flow-chart symbol whose interior contains the criterion for decision or branching.

Diagram: A schematic representation of a sequence of operations or routines.

Diagnostic: Pertaining to the detection and isolation of a malfunction or a mistake.

Digit: One of the symbols 0, 1.....9..... used to designate a quantity smaller than n for a base-n number system.

Directory: An auxiliary list for which each entry corresponds to one sublist in the object list; also an auxiliary list for searching a lower level directory and for finding a subdirectory there.

Direct Access: Retrieval or storage of data by providing a reference to its physical location on a volume.

Displacement: The difference (in bytes) between the contents of a base register (or the address represented by a symbol) and a referenced storage location.

Dummy: Pertaining to the characteristic of having the appearance of a specified thing but not having the capacity to function as such.

Edit: To modify the form or format of data; e.g., to insert or delete characters such as page numbers or decimal points.

Effective Address: The absolute address of the current operand. This may differ from that of the instruction in storage.

Error: 1. A general term to indicate that a data value is not correct or that a machine component is malfunctioning.

2. A specific term for the amount of loss in precision.

Execute: To carry out an instruction or perform a routine.

Explicit Addressing: Specification of an address by a base register and a displacement in the form D(B).

Expression: A symbol or self-defining term used in the operand of a statement.

Extent: Area of a disk file specified by an upper limit and a lower limit.

Fetch (program): 1. To obtain requested load modules and load them into main storage, relocating them as necessary.

2. A control routine that accomplishes: 1 holds to represent data and a card as in 1 before being punched.

File: A collection of related records treated as a unit, e.g., in inventory control, one line of an invoice forms an item, a complete invoice forms a record, and the complete set of such records forms a file.

File Label: Label containing information applicable to a given data file or portion of a data file stored on a particular volume.

File Reorganization: A term used to describe the process of writing a new file from an indexed sequential file, purging records that are tagged for deletion, and their sequential positions in the prime data area.

Fixed-Length Record: A record having the same length as all other records with which it is logically or physically associated.

Flow Chart: A graphical representation for the definition, analysis, or solution of a problem in which symbols are used to represent operations, data, flow and equipment.

Garbage Collection: Collection of cells deleted by linked lists operation to make available again for use.

Hardware: The mechanical magnetic, electrical and electronic devices or components of a resident system from card input.

Hierarchy: An information source which can be represented by an absence.

Hopper: A device that holds cards and makes them available to a card feed mechanism. Synonymous with input magazine. Contrast with card stacker.

Identification: A code number or code name which uniquely identifies a record, block, file or other unit of information.

Image: An exact logical duplicate stored in a different medium.

Immediate Address: The designation of an instruction address which is used as data by the instruction of which it is a part.

Implied Address: The address assigned to a symbol by the basic assembler program.

Index Register: A register whose content is added to or subtracted from the operand address prior to or during the execution of an instruction.

Indexing: A technique of address modification often implemented by means of index registers.

Indirect Pointer: A pointer which depicts the place a file pointer can be obtained.

Information Structure: Inherent property of the information in a given set of data. This property may be by design or by the natural way the data appear in the given set of data.

Initialize: To set certain counters, switches and addresses at specified times in a computer routine.

Initial Program Loader: A Control program, available in a card and a disk version; loads basic monitor into core storage; used to assign actual input/output addresses to symbolic addresses SYSRDR and/or SYSRES; places name of Job Control program into communication region of basic Monitor; required for the initialization of the card-resident or disk-resident system.

Input:

1. The data to be processed.
2. The state of sequence of states occurring on a specified input channel.
3. The device or collective set of devices used for bringing data into another device.
4. A channel for impressing a state on a device or logic element.
5. The process of transferring data from an external storage to an internal storage.

6. Pertaining to any entities such as are quoted above.

Input/Output:

1. Common abbreviation I/O. A general term for the equipment used to communicate with a computer.
2. The data involved in such communication.
3. The media carrying the data for Input/Output.

Input Area: The area of internal storage into which data is transferred from external storage.

Installation: A general term for a particular computing system in the context of the overall function it serves and the individuals who manage it, operate it, apply it to problems, service it, and use the results it produces.

Instruction: A statement that specifies an operation and the values of locations of all operands. In this context, the term instruction is preferable to the terms command or order which are sometimes used as synonyms. Command should be reserved for electronic signals. Order should be reserved for sequence, interpolation and related usage.

Instruction Format: The allocation of bits or characters of a machine instruction to specific functions.

Interpreter: Translator.

Interrupt: A break in the normal flow of a system or routine such that the flow can be reused from that point at a later time.

Inverted and File: A file that consists of one sub-file for each value a key may have and one inverted file corresponding to each key field in the record for which a search may be necessary.

Job Control Program: A control program, resides in core storage between jobs and provides for automatic job-to-job transition. Performs I/O device assignment. It causes the Basic Monitor to load the next program.

Job Control Statement: Any one of the control statements in the input stream that identifies a job or defines its requirements and options.

Label:

1. A physical identification record on magnetic tape located either preceding or following a data file, or both. If a data file extends beyond a single reel of tape, a label can be placed preceding and following the data on each reel.

2. A physical identification record disk which identifies the volume or file.

Language: A defined set of characters which are used to form symbols, words, etc., and the rules for combining these into meaningful communications i.e., English, French, Algol, Fortran, COBOL, etc.

Language Translator: A general term for any assembler, compiler, or other routine statements in one language which produces equivalent statements in another language.

Library: A group of files.

Library Management Programs: Collective term for four service programs; core-image maintenance, macro maintenance, directory service, and allocation organization programs.

Link: The pointer filled for a linked list.

Linkage: The interconnections between a main routine and a closed routine from the main routine.

Linkage Editor: A service program. It relocates programs or phrases and links separately assembled programs or phrases.

Load:

1. To fetch, i.e., to read a load module into main storage preparatory to executing it.
2. To place data into internal storage.

Load Program: A service program that, unlike the other service programs, is contained in punched cards. It creates a disk.

Load Module: The output of the linkage editor; a program in a format suitable for loading into main storage for execution.

Location: A position in storage that is usually identified by an address.

Logical Record: A record identified from the standpoint of its content, function, and use rather than its physical attributes. It is meaningful with respect to the information it contains. (Contrasted with physical record.)

Logical Unit Table: A feature of the Basic Monitor. It has twenty-six logical unit blocks, each of which refers to one specific symbolic I/O address. These symbolic addresses are related to physical I/O device addresses by means of ASSIGN control cards.

Loop: A sequence of instructions that is repeated until a terminal condition occurs.

Machine Address: Same as absolute address.

Machine Code: Same as operation code.

Machine Instructions: An instruction that the particular machine can recognize and execute.

Machine Language: A language that is used directly by a given machine.

Macro Instruction: An instruction that is replaced in a routine by a predetermined sequence of machine instructions.

Macro Library: A disk area containing the macro definitions required by the macro instructions issued in user-written programs. Contains source statements to generate commonly used routines.

Macro Maintenance Program: A service program. It updates the macro library and directory. It is used to add and/or delete macro definitions.

Magnetic Tape: Ink containing particles of magnetic substance which can be detected or read by automatic devices e.g., the ink used for printing on some bank checks for magnetic ink character recognition.

A tape with a magnetic surface on which data can be stored. A tape of magnetic surface used as the constituent in forms of magnetic codes.

Mask: An alphanumeric character string consisting of one or more digits, used to test or alter the contents of storage positions.

Mnemonic Code: A technique to assist the human memory. A mnemonic code enables the original word and is usually easy to remember, e.g., MPY for multiply and ACC for accumulator.

Module: (programming): The input to, or output from, a single execution of an assembler, compiler, or linkage editor; a source, object, or load module; hence, a program unit that is discrete and identifiable with respect to compiling, combining with other units, and loading.

Multi-Extent Disk File: File stored on a disk pack in several areas or defined by more than one extent.

Multilist: A linked list which may contain shared sublists.

Multi-Pack Disk File: File stored on more than one disk pack.

Multi-Volume Disk File: The same as Multi-Pack Disk File.

Name: An alphameric character string, normally used to identify a program.

Object Module: The output of a single execution of an assembler or compiler, which constitutes input to linkage editor. An object module consists of one or more control sections in relocatable, though not executable, form and an associated control dictionary.

Operation:

1. The act specified by a single computer instruction.
2. A program step undertaken or executed by a computer, e.g., addition, multiplication, extraction, comparison, shift or transfer. The operation is usually specified by the operation part of an instruction.

Operation Code: The code that represents the specific operations of a computer.

Output:

1. Data that has been processed.
2. The state or sequence of states occurring on a specified output of a device.
3. The device or collective set of devices used for taking data out of a device.
4. A channel for expressing a state on a device or logic element.
5. The process of transferring data from an internal storage to an external storage.
6. Pertaining to any entities such as are quoted above.

Output Area: The area of internal storage from which data is transferred to external storage.

Overflow: That portion of data that exceeds the capacity or the allocated unit of storage, pertaining to the generation of overflow as in Overflow Area.

Overflow Area: The area a load module or a segment of a load module used for addition of records to a list or sublist which has no room left in it.

Overlay: To place a load module or segment of a load module into main storage locations occupied by another load module or segment which has already been processed.

Pack: To combine two or more units of information into a single physical unit to conserve storage.

Packed Decimal: Storage technique whereby two digits or one digit and sign are stored per byte.

Padding: A technique used to fill out a block of information with dummy records, words or characters.

Phases: The smallest addressable unit in core-image library of a disk-resident system.

Physical and Logical Unit Tables Service Program: A service program. This program (PSERV) is used to display, and/or change the permanent device assignments, and/or to change the configuration byte of the Basic Monitor on the system disk pack.

Physical Record: A record identified from the standpoint of the manner or form in which it is stored and retrieved, that is, one that is meaningful with respect to access. (Contrasted with Logical Record.)

Physical Unit Table: A feature of the Basic Monitor. It has eight physical blocks, each of which contains a physical device address. Pointers to these entries are inserted into the logical unit table by means of ASSGN control cards.

Printer: A device which expresses coded characters as hard copy.

Problem Program: Any of the class of routines that perform processing of the type for which a computing system is intended, and including routines that solve problems, monitor and control industrial processes, sort and merge records, perform computations, process transactions against stored records, etc.

Processing Program: A general term for any program that is not a control program.

Program:

1. The plan for the solution of a problem including data gathering, processing and reporting.

2. A group of related routines which solve a given problem.

Programming Language: A language used to prepare computer programs.

Pseudo-Register: A register with fixed contents used in conjunction with an IBM System/360.

Punched Card: A card punched with a pattern.

Read: To transfer information from an input device to internal or auxiliary storage.

Reader: A device which converts information in one form of storage to information in another form of storage.

Record: A general term for any unit of data that is distinct from all others when considered in a particular context.

Register: A device capable of storing a specified amount of data such as one half word.

Relative Address: An address expressed by a previously defined symbol and a displacement. (e.g., FID-10)

Relocatable Area: An area on the system disk pack to temporarily hold an object module, thus permitting the assembly or compilation and the execution of a program or program phase in one job.

Relocate: In programming, to move a routine from one portion of internal storage to another and to automatically adjust the necessary address references so that the routine, in its new location, can be executed.

Report Generator and Report Program Generator (RPG): A program which constructs reports or report-writing programs in accordance with input specifications of the data file and of the desired report.

Reset:

1. To restore a storage device to prescribed initial state, not necessarily that denoting zeros.

2. To place a binary cell into the zero state.

Restart: To return to a previous point in a program and resume operation from that point.

Search Key: A key used to find a record which has a presearched identity within a file.

Seek: To position the access mechanism of a direct-access device at a specified location.

Self-Defining Term: A term with an implied value (e.g., 300, X"2A", C"F").

Semantics: The meanings which govern subsequent interpretation.

Service Programs: A collective term used to refer to the Library Management Program.

Single Extent Disk File: A file stored on disk where the file has exactly one extent.

Sort: The process by which a data set of logical records is sequenced according to the collating-sequence value of the control field of the records. Also a program that performs the process.

Source Language: A language that is an input to a given translation process.

Source Program: A program written in a source language.

Special Characters: In a character set, a character that is neither a numeral nor a letter, e.g., *, -, \$ and blank.

Statement: In computer programming, a meaningful expression or generalized instruction in a source language.

Step:

1. One instruction in a computer routine.
2. To cause a computer to execute one instruction.

Store:

1. To enter data into a storage device.
2. To retain data in a storage device.

Storage:

1. Pertaining to a device into which data can be entered and from which it can be retrieved at a later time.
2. Loosely, any device that can store data.

Storage Allocation: The assignment of blocks of data to specified blocks of storage.

Storage Capacity: The amount of data (in bytes) that can be contained in a storage device.

Subroutine: A routine that can be part of another routine.

Switch:

1. A symbol used to indicate a branching point, or a set of instructions to condition a branch.
2. A physical device which can alter flow.

Symbol Table: A mapping for a set of symbols to another set of symbols or numbers.

Symbolic Address: An address expressed in symbols convenient to the programmer.

Symbolic I/O Device (e.g., SYRES, SYSIPP, SYS005): This address is related to an actual address by means of the logical unit table.

Symbolic Language: The discipline that treats formal logic by means of a formalized language or symbolic calculus whose purpose is to avoid the ambiguities and logical inadequacies of natural language. Advantages of the symbolic method are greater exactness of formulation and greater exactness of formulation and power to deal with complex material.

Syntax: The rules for constructing admissible combinations of the characters in the basic alphabet.

System:

1. A collection of consecutive operations and procedures required to accomplish a specific objective.
2. An assembly of objects united to form a functional unit.

System Directory: A table on the system disk pack listing the addresses and sizes of the core-image library directory, the macro library and the directory, and the locatable area.

System Disk Pack: The disk pack on which the user's disk-resident system is located.

Table: A collection of data, each item being uniquely identified either by some label or by its relative position.

Table Look-Up: A procedure for obtaining the function value corresponding to an argument from a table of function values.

Throughput: A measure of system efficiency, the rate at which work can be handled by the computing system.

Truncate: To cut off a specified spot (as contrasted with round or pad).

Unpack: To recover the original data from packed data.

User: Anyone who requires the services of a computing system.

Volume: That portion of a single unit of storage media that is accessible to a single read/write mechanism. For example, a reel of magnetic tape for a 2415 magnetic tape drive, or one 1316 Disk Pack for an IBM 2311 Disk Storage Drive.

Volume Label: Label which uniquely identifies the volume.

Volume Table of Contents (VOTC): A table associated with a disk volume, which describes each data set on the volume.

Zero: The elimination of non-significant zeros in a number.

Zone: The 12, 11 or 0 punches in IBM card code.

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